

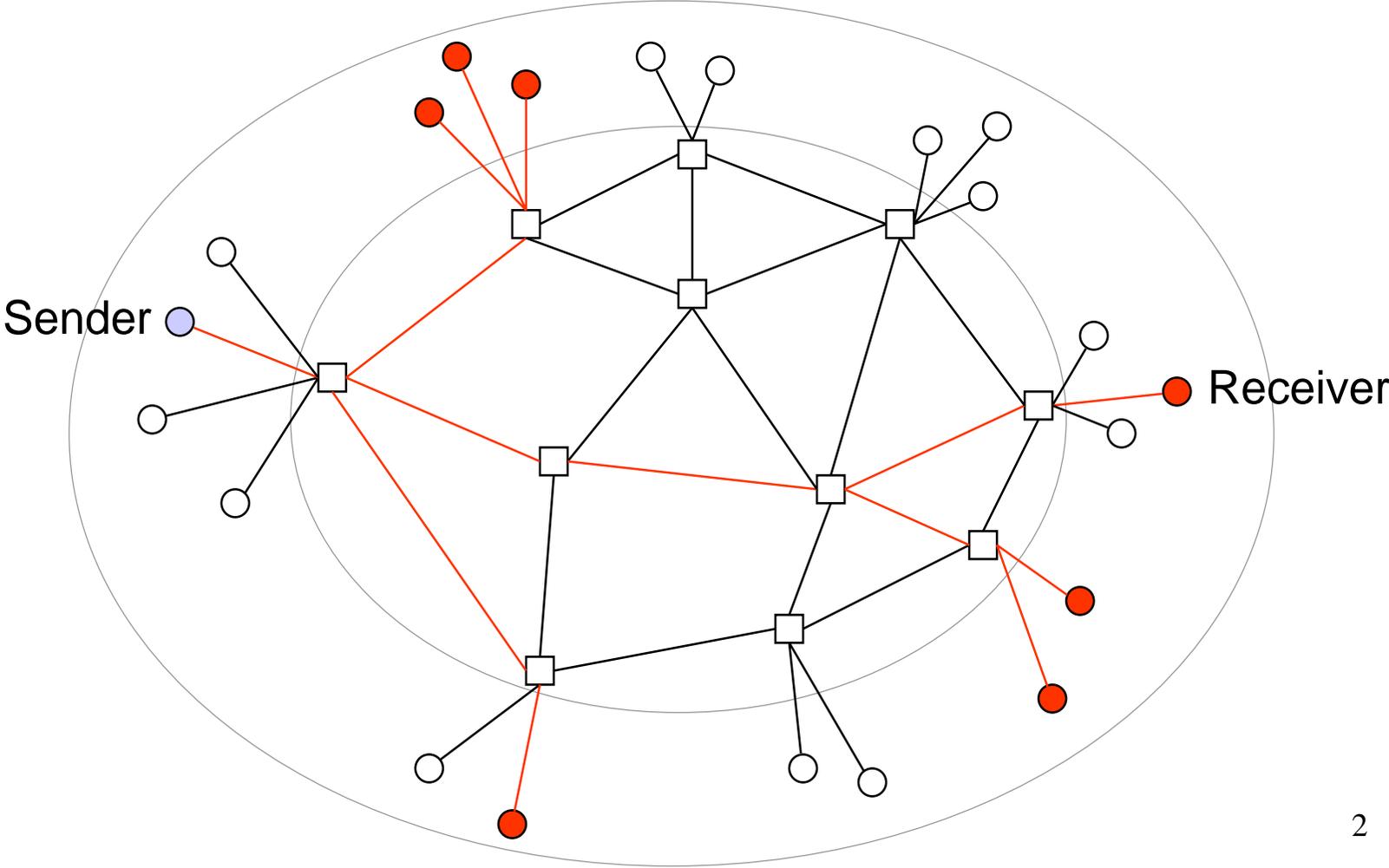


New Algorithms for Calculating Blocking Probabilities in Multicast Networks

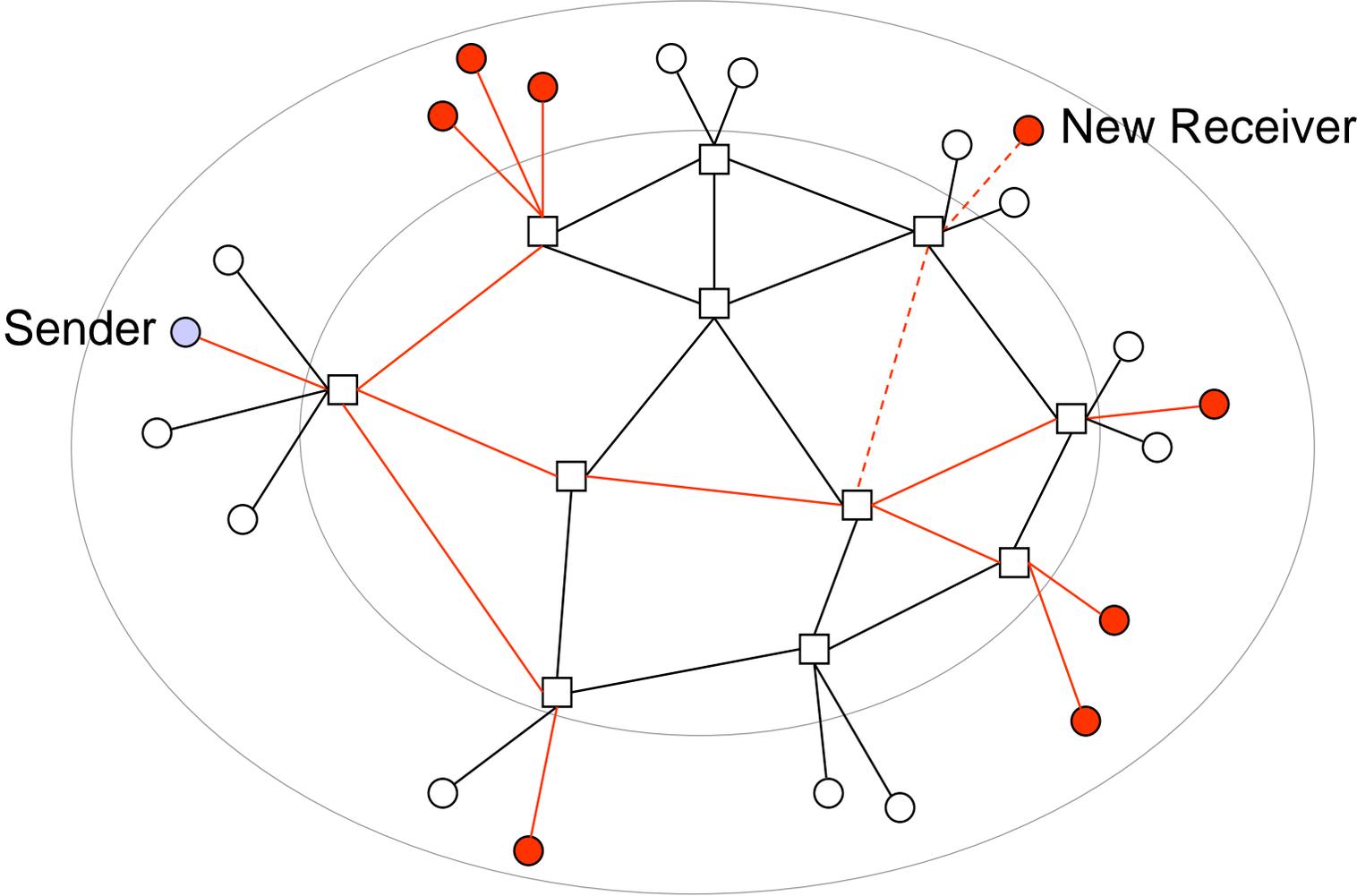
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Multicast connection with dynamic membership (1)

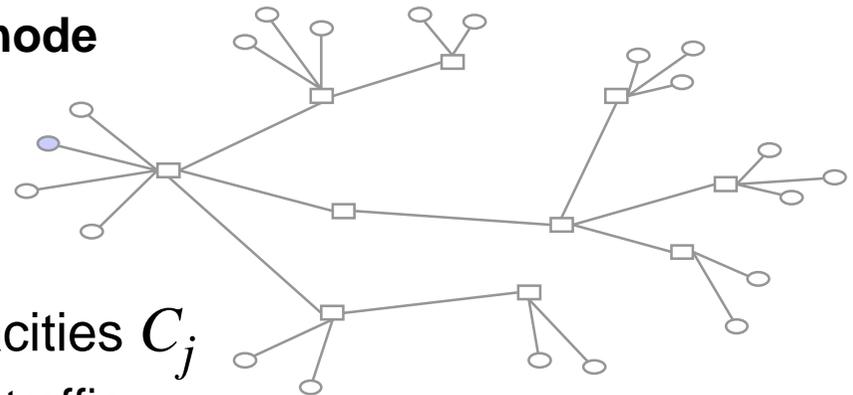


Multicast connection with dynamic membership (2)



Setup

- A (unique) **service center** offers a variety of **channels** $i \in I$
 - e.g. audio or video streams
 - required capacities d_i
- Each channel is delivered to users $u \in U$ by a **multicast connection** with **dynamic membership**
- Each multicast connection uses the same **routing tree**
 - service center located at the **root node**
 - users located at the **leaf nodes**
- Physical links $j \in J$ with finite capacities C_j
 - (possibly) shared with bg (unicast) traffic



Teletraffic study

- Calculation of call blocking (of a user requesting a channel)
 - call blocking = probability that the user fails to join the requested multicast connection
- **1st step:**
Find out how to calculate time blocking (for that user)
 - time blocking = probability of such network states that do not allow the user to join the requested multicast connection
- **2nd step:**
Express call blocking by means of time blocking (possibly in a slightly modified network)
 - exact expression depends always on the chosen user model

Time blocking (1)

- \mathbf{X} = network state (**without** capacity constraints)
- $\tilde{\mathbf{X}}$ = network state (**with** capacity constraints)
- $\tilde{\Omega}$ = network state space (**with** capacity constraints)
- $\tilde{\Omega}_{ui}$ = nonblocking states for user u and channel i
- B_{ui}^t = time blocking probability for user u and channel i

$$B_{ui}^t := 1 - P\{\tilde{\mathbf{X}} \in \tilde{\Omega}_{ui}\} = 1 - \frac{P\{\mathbf{X} \in \tilde{\Omega}_{ui}\}}{P\{\mathbf{X} \in \tilde{\Omega}\}}$$

Due to the Truncation Principle!

numerator

denominator

Time blocking (2)

- In principle, there is a closed form analytical expression both for the numerator and the denominator
- **Problem:** computationally extremely complex
 - exponential in U
- **Solution:** use **recursive convolution-truncation algorithms** to calculate the numerator and the denominator
 - linear in U

The algorithm

- Denominator:

$$P\{\mathbf{X} \in \tilde{\Omega}\} = \sum_{\mathbf{y}} Q_J(\mathbf{y})$$

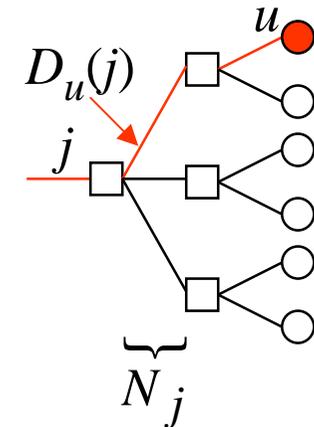
truncation operator

link state

root link

$$Q_j(\mathbf{y}) = \begin{cases} T_j[\pi_j](\mathbf{y}), & j \in U \\ T_j[\bigotimes_{k \in N_j} Q_k](\mathbf{y}), & j \notin U \end{cases}$$

“convolution” operator



- Numerator:

$$P\{\mathbf{X} \in \tilde{\Omega}_{ui}\} = \sum_{\mathbf{y}'} Q_J^{ui}(\mathbf{y}')$$

ui-mod. trunc. oper.

i-modified link state

root link

$$Q_j^{ui}(\mathbf{y}') = \begin{cases} T_j^{ui}[\pi_j](\mathbf{y}'), & j = u \\ T_j^{ui}[Q_{D_u(j)}^{ui} \bigotimes_{k \in N_j \setminus R_u} Q_k](\mathbf{y}'), & j \in R_u \setminus \{u\} \end{cases}$$

route to user u

modified “convolution” operator

Cases studied

- Single link analysis (Karvo, Virtamo, Martikainen & Aalto, 1997-1998)
 - starting point
- Network wide analysis (Nyberg, Virtamo & Aalto, 1999)
 - all channels handled individually
 - first convolution-truncation algorithm

- Multi-class case (Aalto, Karvo & Virtamo, 2000)
 - class = group of statistically indistinguishable channels
 - combinatorial convolution needed
- Multi-layer case (Karvo, Aalto & Virtamo, 2000-2001)
 - layered coding of audio/video streams
 - individual channels / single class / multiple classes

New algorithms!

Individually handled channels

- Denominator:

- **Link state** = $\mathbf{y} = (y_1, \dots, y_I)$
 - $y_i = 1$ (0) if channel i is active (idle)

- **Q-function:**

$$Q_j(\mathbf{y}) = P\{\mathbf{Y}_j = \mathbf{y}; r(\mathbf{Y}_{j'}) \leq C_{j'}, j' \in M_j\}$$

total capacity req.

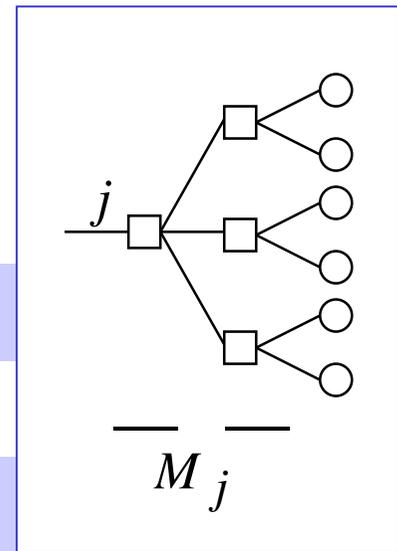
- **Truncation:**

$$T_j f(\mathbf{y}) = f(\mathbf{y}) \mathbb{1}\{r(\mathbf{y}) \leq C_j\}$$

- **OR-convolution:**

$$[f \otimes g](\mathbf{y}) = \sum_{\mathbf{u} \oplus \mathbf{v} = \mathbf{y}} f(\mathbf{u}) g(\mathbf{v})$$

componentwise OR



- Numerator: slightly different!

Single class

- Denominator:
 - **Link state** = n = number of active channels
 - **Q -function:**

$$Q_j(n) = P\{N_j = n; r(N_{j'}) \leq C_{j'}, j' \in M_j\}$$

total capacity req.

- **Truncation:**

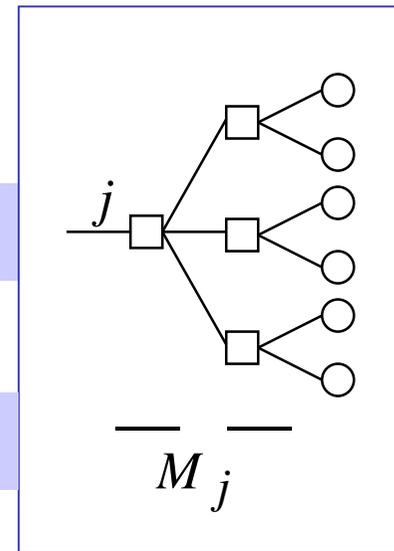
$$T_j f(n) = f(n) \mathbb{1}\{r(n) \leq C_j\}$$

- **Combinatorial convolution:**

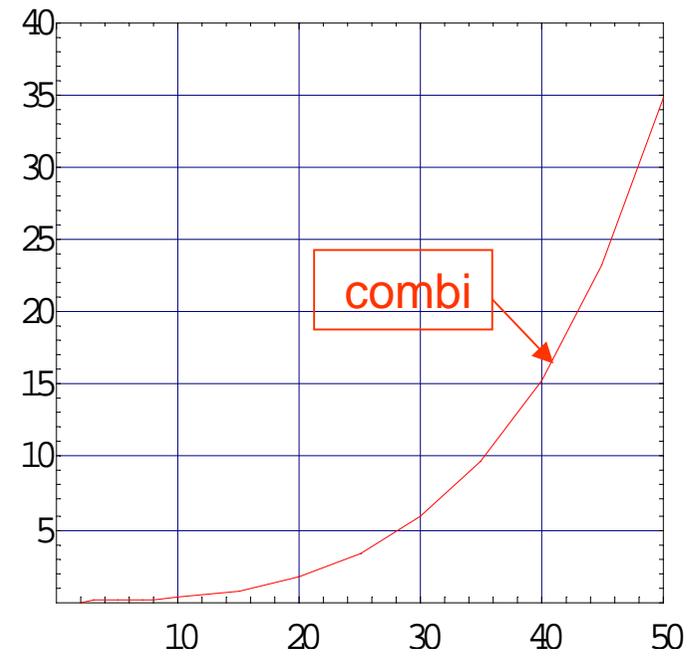
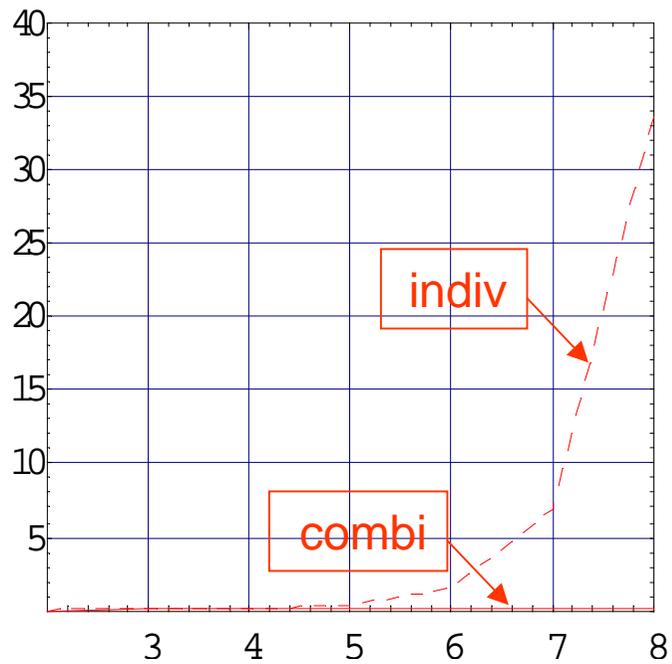
$$[f \otimes g](n) = \sum_{l,m} s(n | l, m, I) f(l) g(m)$$

- Numerator: slightly different!

simple combinatorial factor by a random sampling' -argument



Processing time (sec) vs. number of channels



Multi-class case

- Denominator:

- **Link state** = $\mathbf{n} = (n_1, \dots, n_K)$

- n_k = number of active channels of class k

- **Q-function:**

$$Q_j(\mathbf{n}) = P\{\mathbf{N}_j = \mathbf{n}; r(\mathbf{N}_{j'}) \leq C_{j'}, j' \in M_j\}$$

total capacity req.

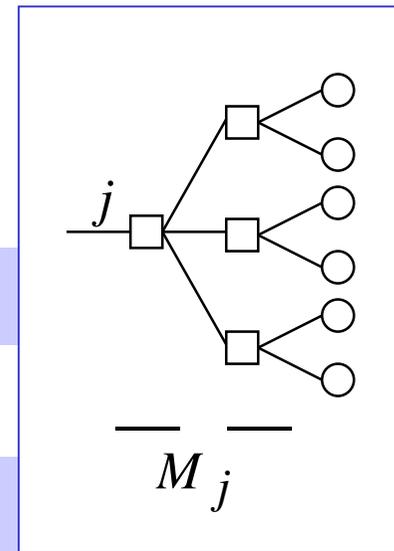
- **Truncation:**

$$T_j f(\mathbf{n}) = f(\mathbf{n}) \mathbb{1}\{r(\mathbf{n}) \leq C_j\}$$

- **Combinatorial convolution:**

$$[f \otimes g](\mathbf{n}) = \sum_{\mathbf{l}, \mathbf{m}} \prod_k s(n_k | l_k, m_k, I_k) f(\mathbf{l}) g(\mathbf{m})$$

the same combinatorial factor as before



- Numerator: slightly different!

Multi-layer case (individually handled channels)

- Denominator:

- **Link state** = $\mathbf{y} = (y_1, \dots, y_I)$

- y_i = layer on which channel i is active (0 when idle)

- **Q-function:**

$$Q_j(\mathbf{y}) = P\{\mathbf{Y}_j = \mathbf{y}; r(\mathbf{Y}_{j'}) \leq C_{j'}, j' \in M_j\}$$

total capacity req.

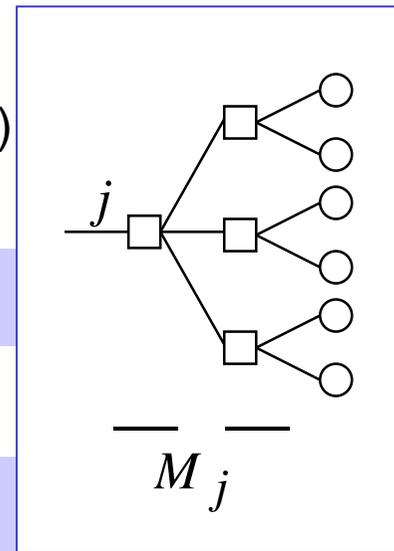
- **Truncation:**

$$T_j f(\mathbf{y}) = f(\mathbf{y}) \mathbb{1}\{r(\mathbf{y}) \leq C_j\}$$

- **Max-convolution:**

$$[f \otimes g](\mathbf{y}) = \sum_{\max\{\mathbf{u}, \mathbf{v}\} = \mathbf{y}} f(\mathbf{u})g(\mathbf{v})$$

- Numerator: slightly different! componentwise max



Multi-layer case (single class)

- Denominator:

- **Link state** = $\mathbf{n} = (n_1, \dots, n_L)$

- n_l = number of active channels on layer l

- **Q-function:**

$$Q_j(\mathbf{n}) = P\{\mathbf{N}_j = \mathbf{n}; r(\mathbf{N}_{j'}) \leq C_{j'}, j' \in M_j\}$$

total capacity req.

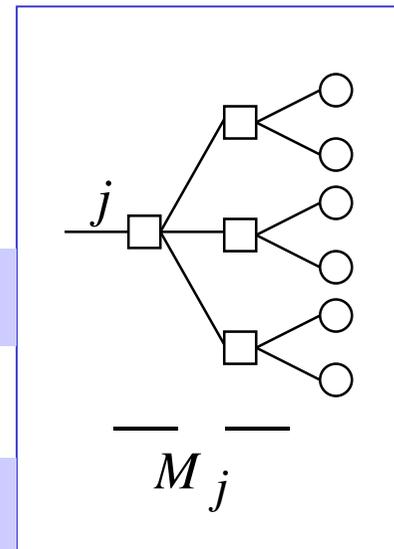
- **Truncation:**

$$T_j f(\mathbf{n}) = f(\mathbf{n}) \mathbb{1}\{r(\mathbf{n}) \leq C_j\}$$

- **Combinatorial convolution:**

$$[f \otimes g](\mathbf{n}) = \sum_{\mathbf{l}, \mathbf{m}} s(\mathbf{n} | \mathbf{l}, \mathbf{m}) f(\mathbf{l}) g(\mathbf{m})$$

much more complicated combinatorial factor



- Numerator: slightly different!

Multi-layer case (multiple classes)

To be done!

THE END

