



Some Findings from Assignment 1

Wide variety of specifications:

- ▶ From 1 to 10 pages
- ▶ More protocol spec vs. more implementation spec
- ▶ More or less complete (at a first glance)



Encodings

- ▶ JSON (text-based)
 - Base64-encoded data
- ▶ Unknown
- ▶ 3 x Box notation (binary)
 - Single bit to distinguish between data and control packet
 - Packet type field
- ▶ Text-based for control, binary for packets?
- ▶ RFC 822-style with 8-bit encoding for the data part
- ▶ Binary 64-bit words + scrambling to avoid deterministic bit errors
- ▶ HTTP-style + RFC 822



Protocol Operation (1)

- ▶ Connection setup
 - Explicit establishment via some handshake mechanism
 - Two-way, three-way, four-way, cookies against DoS
- ▶ File transmission
 - Various forms of checksums (e.g., MD5, HMACs)
 - Sequence numbers
 - Data + ACK (cumulative, selective ACKs)
 - Data + NACK + a final ACK
 - Dynamic RTO calculation
- ▶ Flow control
 - Explicit window size indication
 - Fixed window (negotiated at session setup)
 - Window size derived from delay x bandwidth product



Protocol Operation (2)

- ▶ Completion
 - Explicit end signaling + confirmation
 - Implicit server-side detection leads to confirmation
 - Plain shutdown (and hope)?
- ▶ Parallel upload
 - Transparent to the protocol



Some Observations on Possible Constraints

- ▶ Limited sizes (filename < 255 characters)
- ▶ Manual mapping: media type -> binary constant
 - Need to keep up to date
- ▶ Sometimes many options
 - Is there a common baseline?
- ▶ Did you think about sequence number wrap around?



Protocol Design

Assignment 2:

1. Solution analysis
2. Stress tests
3. adaptive fip



Reminder: Group Info Needed

- ▶ Send one email per group in exactly the following format (one line per group member)
“Last name:First name:ID:email address”
Mustermann:Erika:12345Z:erikam@example.com
- ▶ Just about two groups (out of nine!) got this right!



1. Solution Analysis

- ▶ Take a look at someone else's protocol from assignment 1
- ▶ Write down your observations (high level perspective)
 - Is the design spec sufficient to create interoperable implementations?
 - Where is it not? What is missing?
 - Is the protocol spec robust?
 - Do you find errors? (concepts rather than details)
 - What else do you observe?
- ▶ Practical matters:
 - We will pair two groups (in one case: three groups)
 - We will send out the design documents to the respective groups
 - You may update each other later on (but CC our course assistants)



2. Attack your implementations

- ▶ Analyze your peer group's and your own protocol specifications
 - Which are angles that an attacker could use?
 - To kill the server
 - To launch a DoS attack against a competing journalist
 - ...
- ▶ Go for it!
 - Challenge your own implementation
 - Challenge your peer group's implementation
- ▶ Important: both 1) and 2) are to learn
 - Grading of another solution will not depend on what you say about it
 - You may perform analysis and testing jointly
 - But we want independent submissions



Attacking fip... (1)

- a) Write a small program that can generate arbitrary UDP packets
 - Use it to generate data and control packets to send to your uft client and/or server
 - Packets should be somewhat close to real ones, yet random
 - Some suggestions: right total size but arbitrary contents, inconsistent field values (e.g., mismatch of packet length and length field), undefined values for selected fields, strange file names, ...
 - Observe and document what happens
 - Suggest reasonable fixes
 - To your protocol specification
 - To your implementation
 - Implement selected ones that can be done with reasonable effort



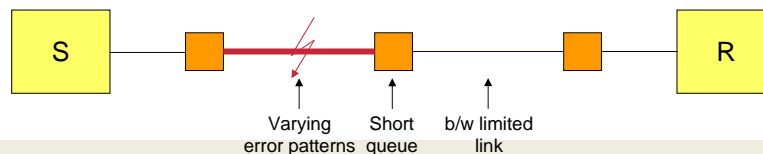
Attacking fip... (2)

- b) Exploit knowledge about your / the other protocol
 - Construct malicious packets to subvert protocol operation
 - Of the server
 - Of the communication relationship between the server and another client
 - ▶ What do you learn?
 - Document your observations
 - What type of protocol refinements (if any) would be needed to fix this?
 - If you like, try this out with your small random packet generator
 - Sniff an existing session (e.g., one transmitting a large file at low rate)
 - Add a second sender sending a different file and check the result



3. Adaptive FIP

- ▶ Optimize your file transfer protocol
 - Imagine a cost function like $C = C_d + C_v$
 - $C_d = \text{EUR } 0.01 / \text{ms delay (reception time - transmission timestamp)}$
 - $C_v = \text{EUR } 0.01 / \text{byte sent}$
 - Minimize the cost for a file transfer
 - Calculate the cost for delay on the receiver side
 - Running both sender and receiver on the same machine provides clock sync
 - Return the cost in the last confirmation
 - Calculate the cost for the volume on the sender side and compute the total cost
 - How would your algorithm change if the cost function changes?
- ▶ Target environment
 - Error-prone link concatenated with low speed link, high latency



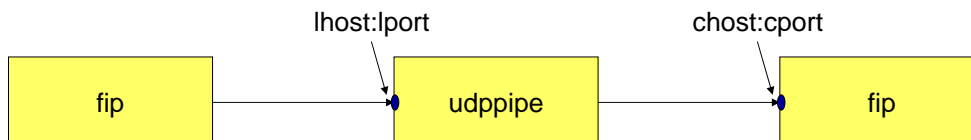
3. Adaptive FIP (2)

- ▶ The error pattern will vary in the mid-term
 - Adapt!
- ▶ Enhance your protocol for the optimization
 - Retransmissions incur cost in terms of delay (but minimize overhead)
 - Doing only retransmissions will be too costly
 - Proactive repair (e.g., XOR-based) FEC incurs cost in terms of overhead
 - Sending every packet ten times (or 200% FEC) will also be too costly
 - Provide feedback about the observed errors from the receiver to the sender
 - Examples: loss rate, loss patterns
 - Make the sender adapt for the optimization
 - Reduce FEC, increase FEC
- ▶ Remember: FEC requires packets of equal length
 - You may need to do some padding
- ▶ Implement and test!

Testing: udppipe

```
udppipe -l [lhost:]lport -c [chost:]cport -b <bitrate> [-d <delay>]
```

- l: transport address to receive UDP packets on from first uft peer; in the opposite direction, packets are sent to the address they were received from
- c: transport address to send UDP packets to (the other uft peer needs to transmit its responses to the address taken from recvfrom ())
- b: bit rate specified in kbit/s
- d: delay specified in milliseconds (default: none)



Only the forward path will be impacted by udppipe modifications.