

NAT Traversal in SIP

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Outline

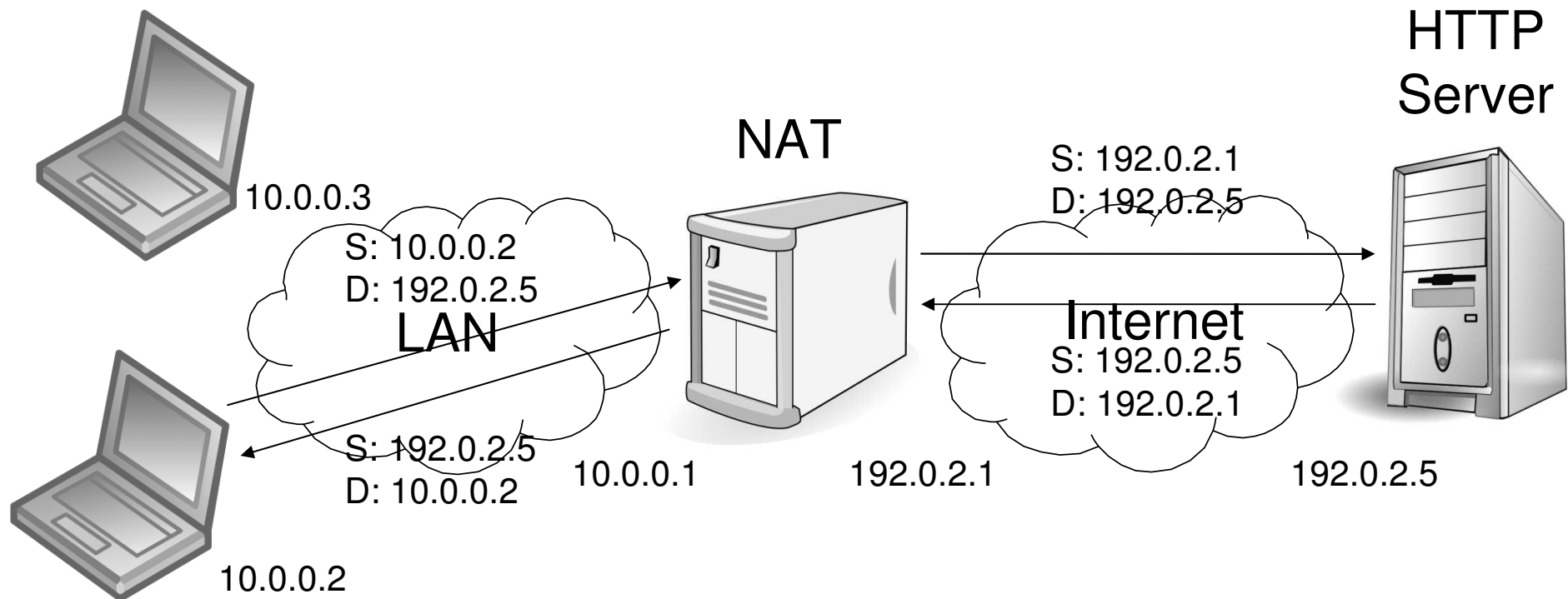
- § Introduction to NATs
- § NAT Behavior (actual and recommended)
 - UDP
 - TCP
- § IAB UNSAF Considerations
- § STUN
- § STUN Relay Usage
- § NAT Traversal in SIP
 - SIP Response Routing
 - User Agent Reachability
 - ICE

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Origin of NATs

- § Created to resolve the IPv4 address exhaustion problem
- § Designed with the web in mind
 - Client/server paradigm



Side-effects of NATs

- § Hosts behind NATs are not reachable from the public Internet
 - Sometimes used to implement security
 - Breaks peer-to-peer (as opposed to client/server) applications
- § NATs attempt to be transparent
 - Troubleshooting becomes more difficult
- § NATs are a single point of failure
- § NATs' behavior is not deterministic
 - Difficult to build applications that work through NATs

IETF BEHAVE WG

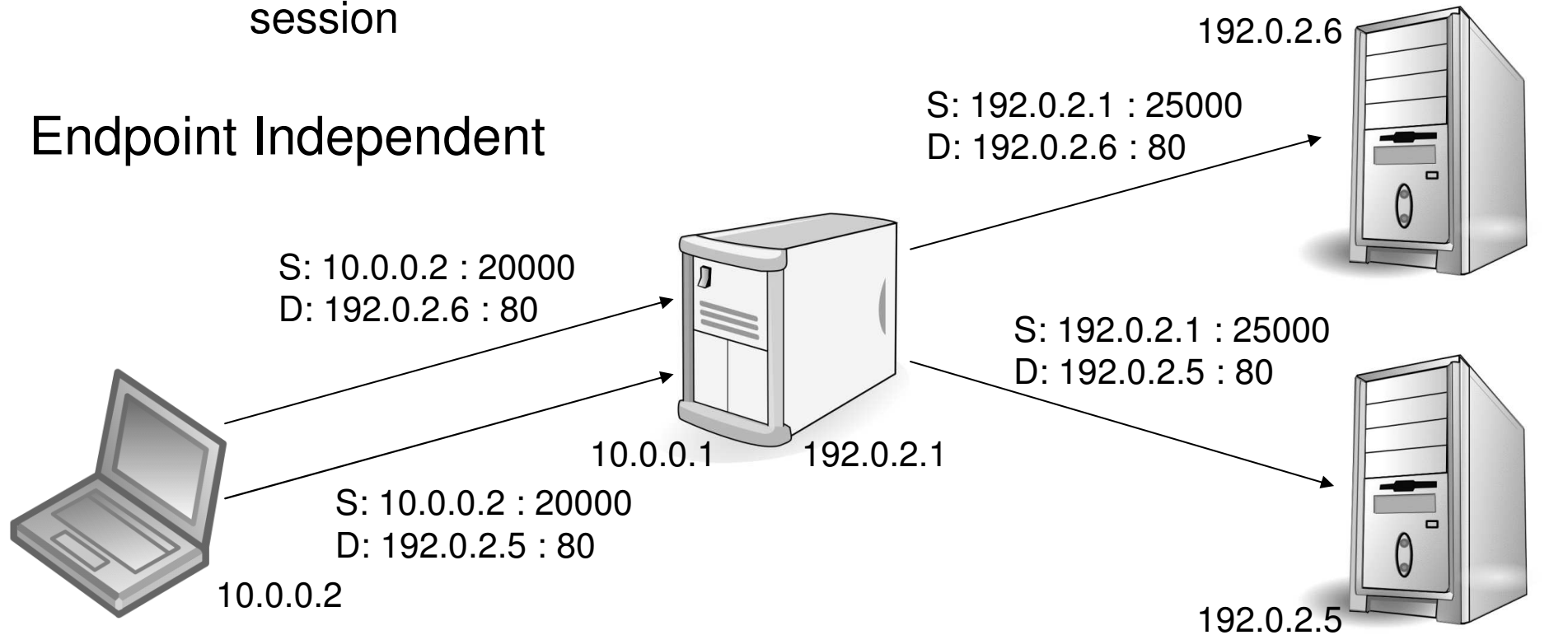
- § Classification of current NAT behaviors
 - Existing terminology was confusing
 - § Full cone, restricted cone, port restricted cone, and symmetric
 - New terminology needed
- § Recommendations for NAT vendors
 - BEHAVE-compliant NATs are deterministic
 - Easier to build applications

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Mapping Behavior

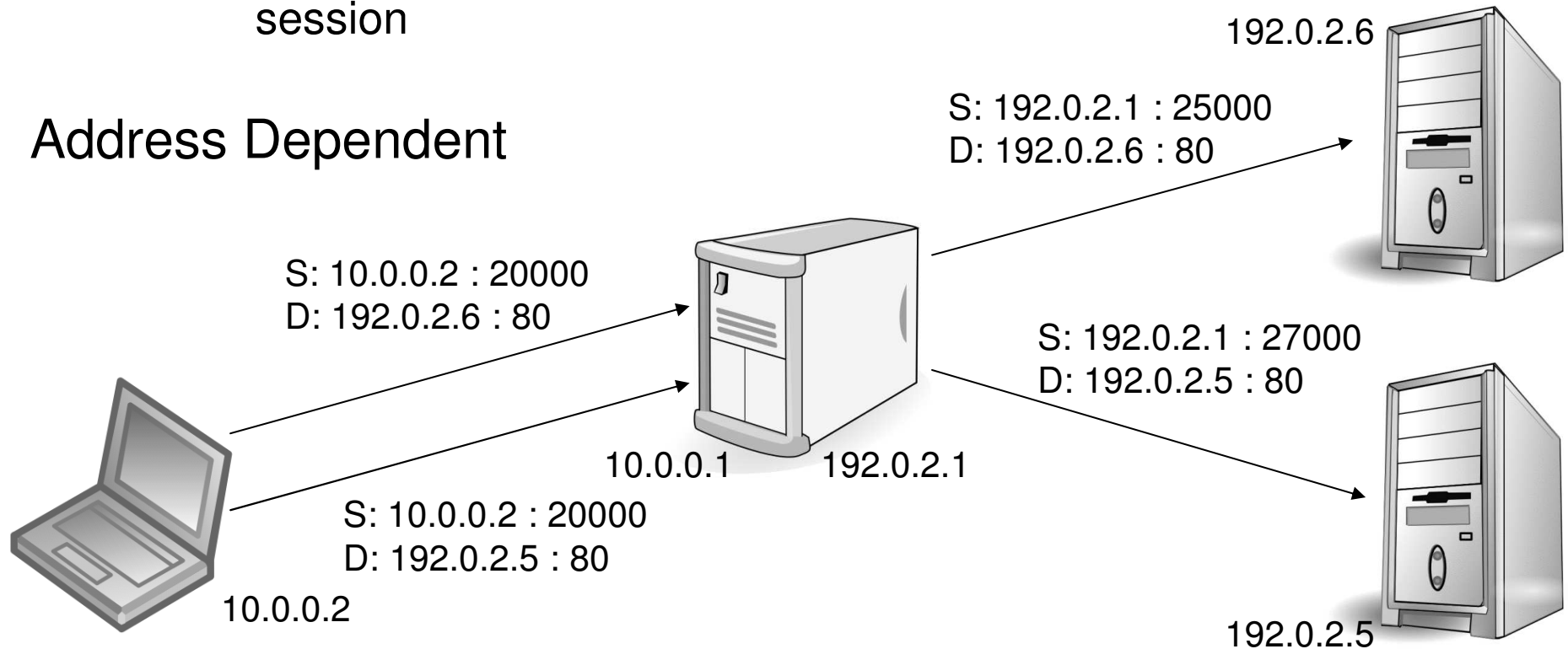
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- Endpoint independent: same mapping to different sessions
 - § MUST use it
 - Address dependent: same mapping to sessions to the same host
 - Address and port dependent: a mapping only applies to one session



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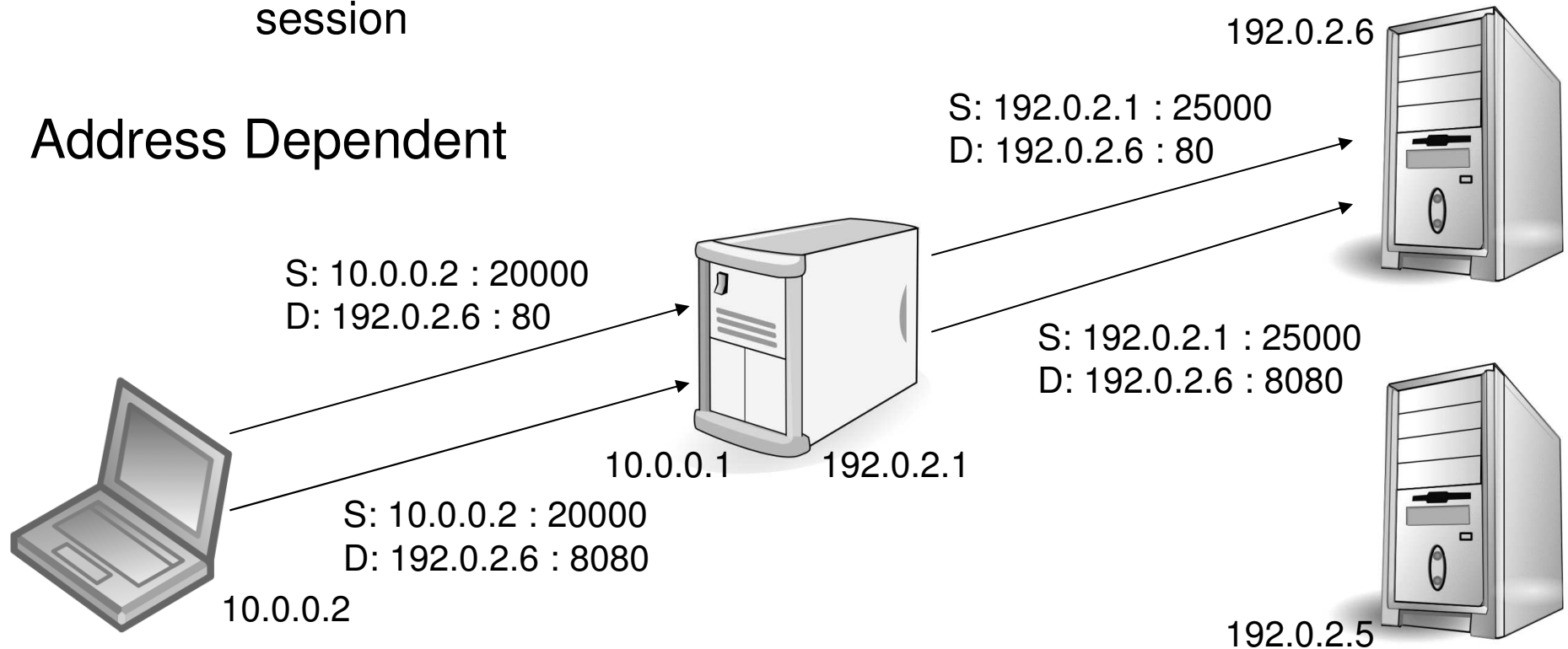
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Address Dependent



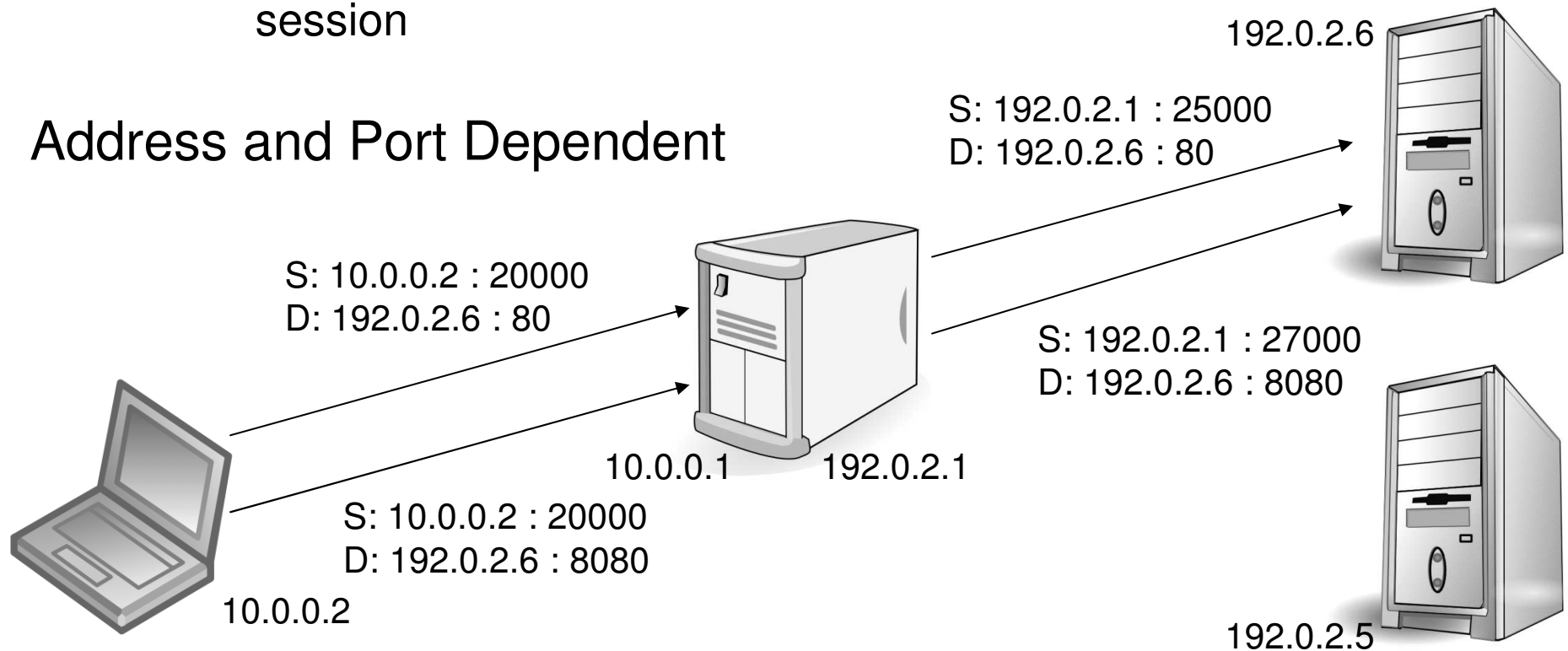
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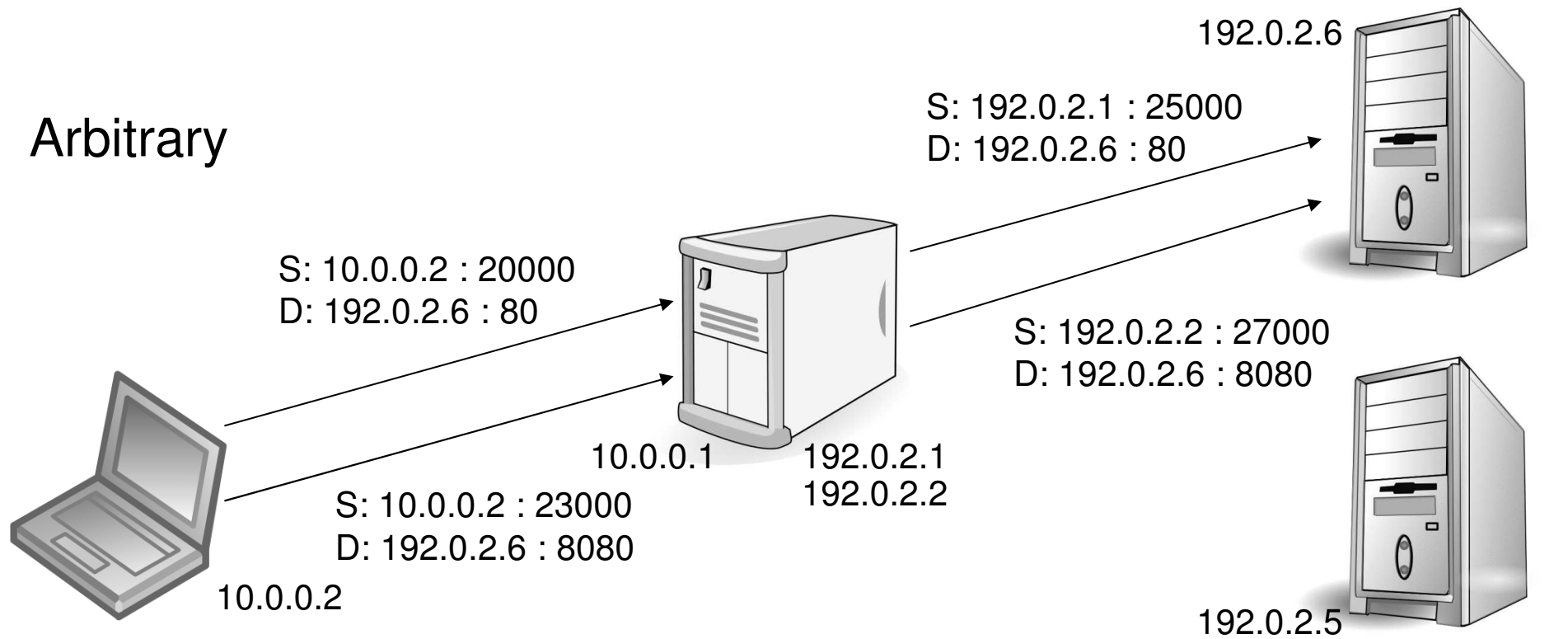


IP Address Pooling Behavior

§ NATs with a pool of external IP addresses

- Arbitrary: an endpoint may have simultaneous mappings corresponding to different external IP addresses of the NAT
- Paired: same external IP address of the NAT

§ RECOMMENDED

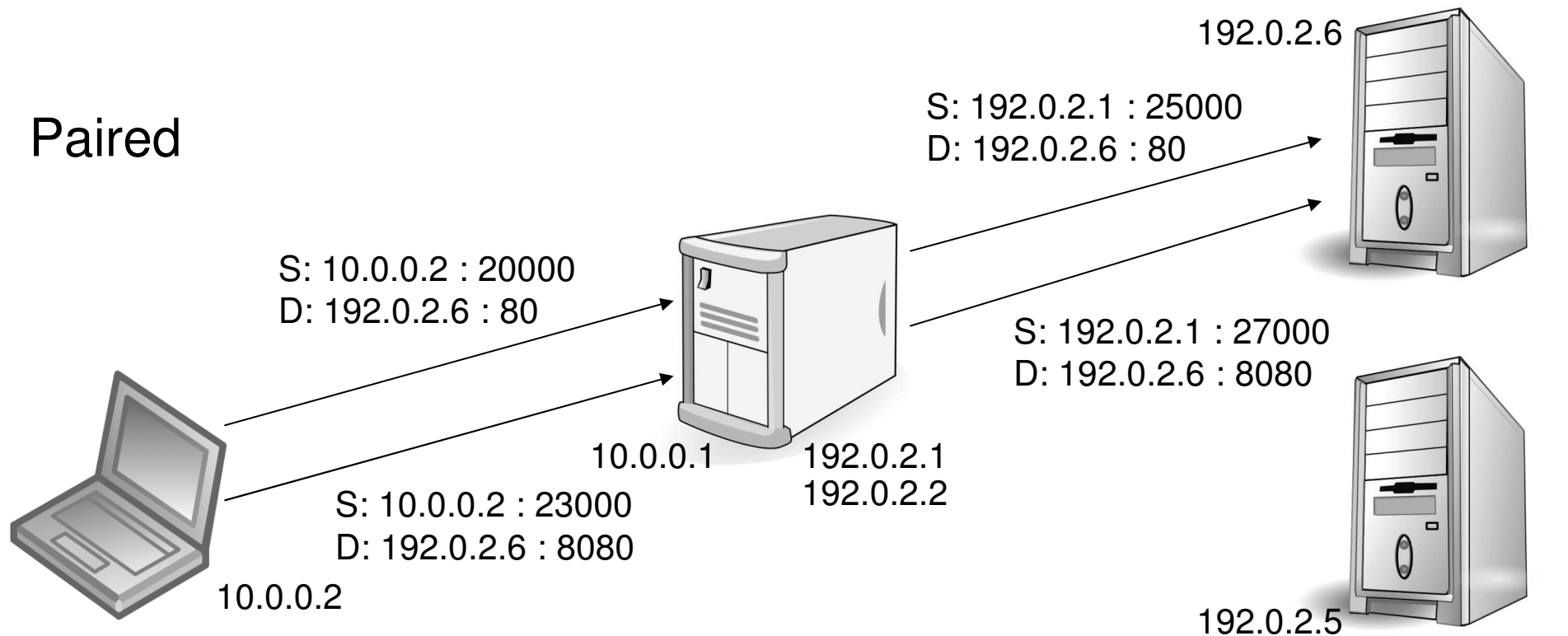


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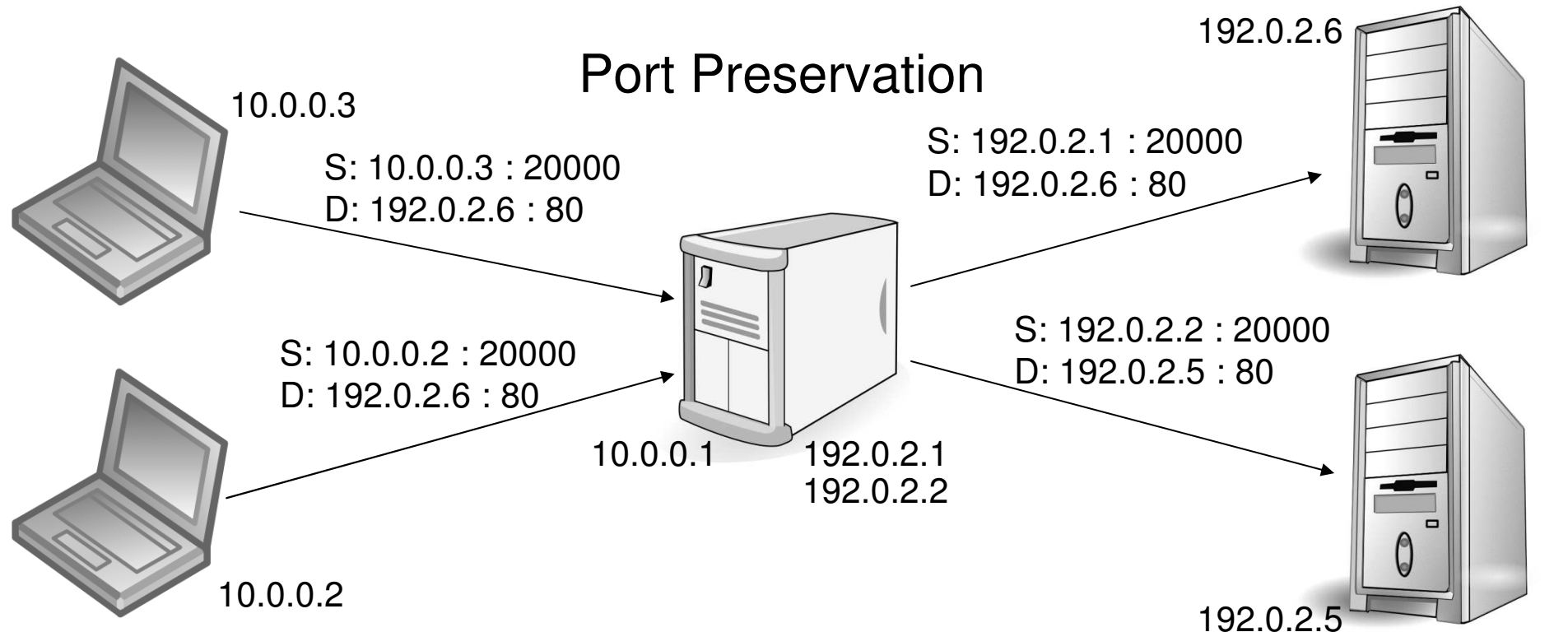
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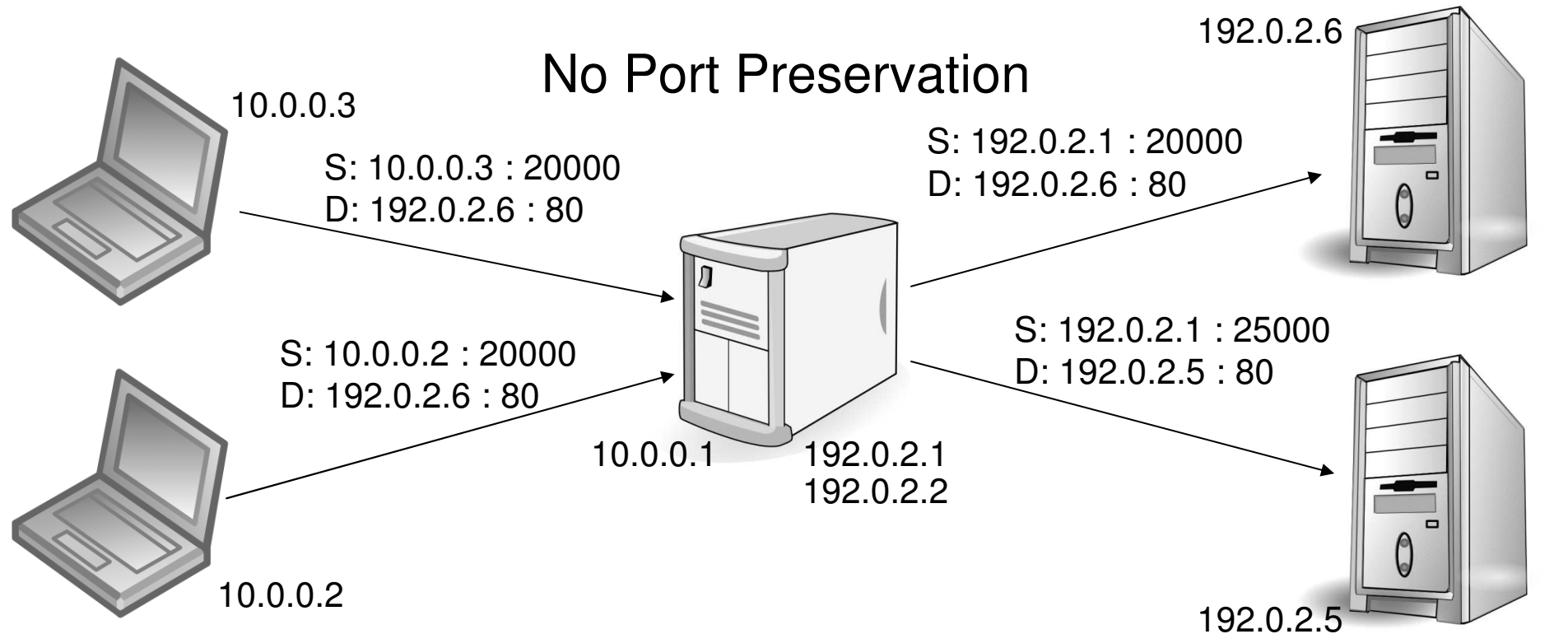
Port Assignment

- § Port preservation: preserves the port as long as there are available IP addresses in the NAT's pool
- § No port preservation
- § Port overloading: the port is preserved always, even without available IP addresses in the NAT's pool
 - The NAT relays on the source of the response



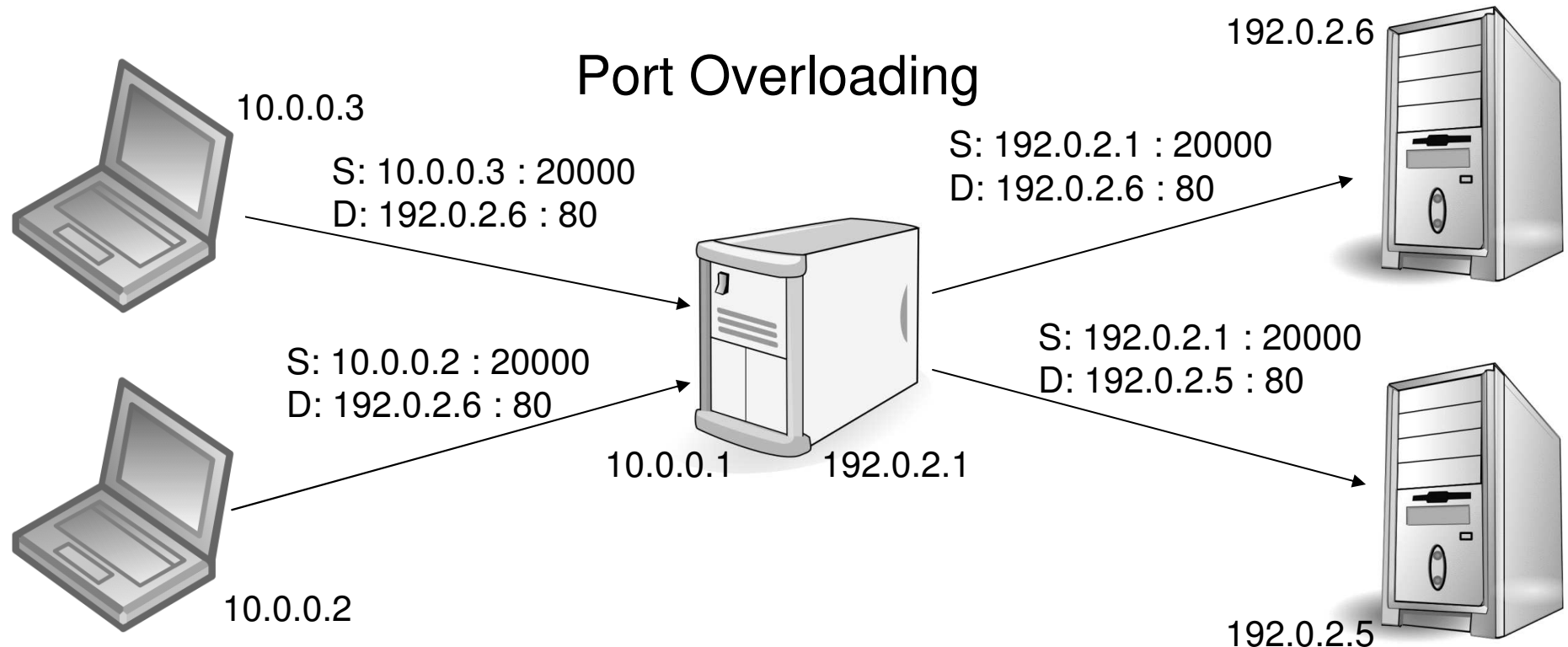
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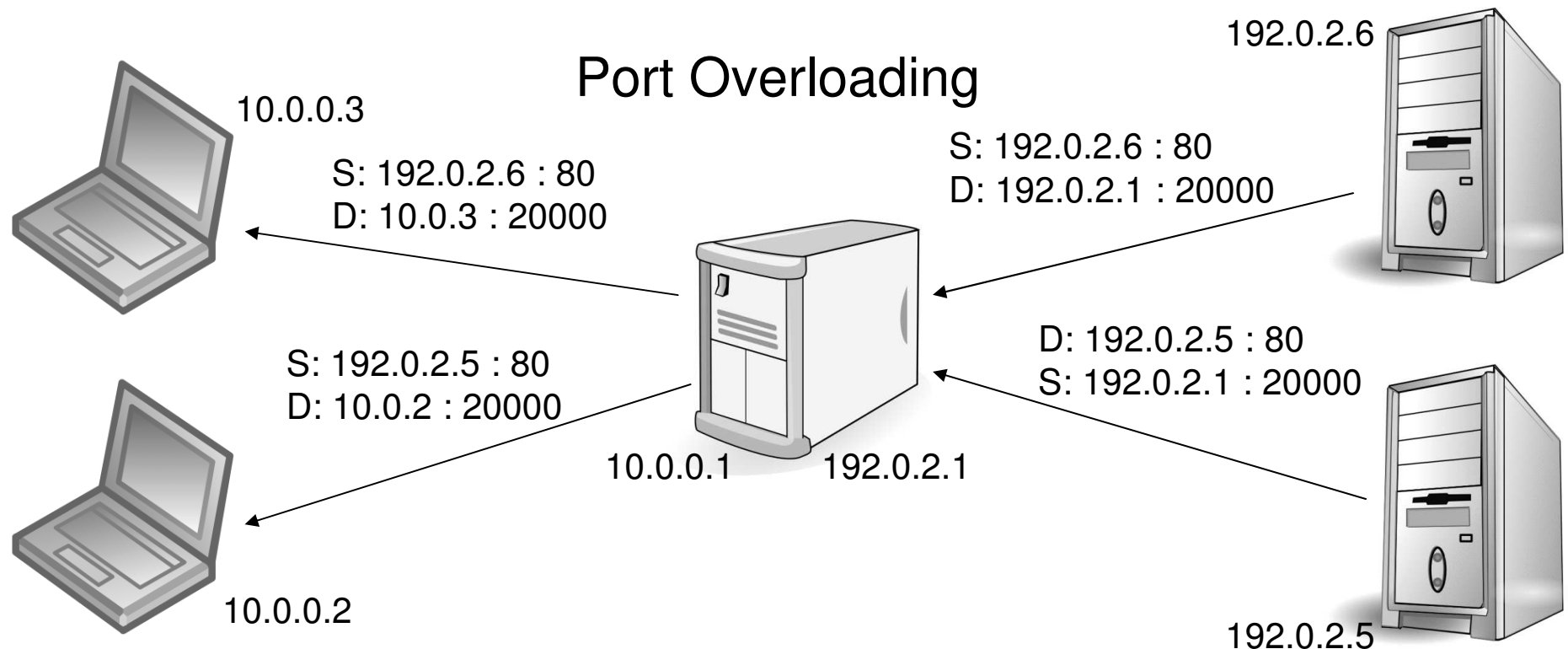
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Port Ranges

- § 1- 1023 Well known
- § 1024 – 49151 Registered
- § 49152 – 65535 Dynamic / Private

- § RECOMMENDED to preserve the following ranges
 - 1 – 1023
 - 1024 – 65535

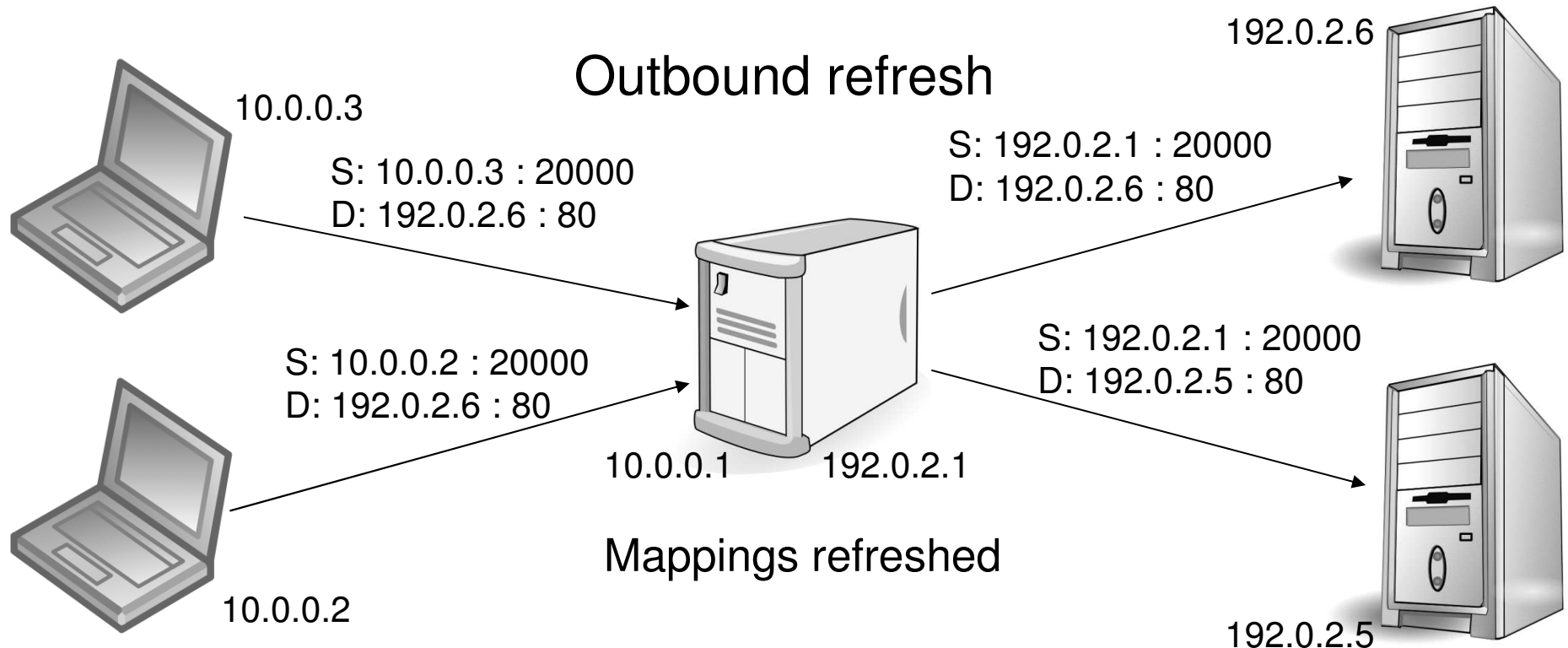
- § Port overloading **MUST NOT** be used
 - Problems when two internal hosts connect to the same external host
- § It is RECOMMENDED that NATs preserve port parity (even/odd)
- § No requirement for port contiguity

Mapping Timeout

- § A NAT UDP mapping **MUST NOT** expire in less than 2 minutes
- § NATs can have application-specific timers
 - Well-known ports
- § It is **RECOMMENDED** to use more than 5 minutes

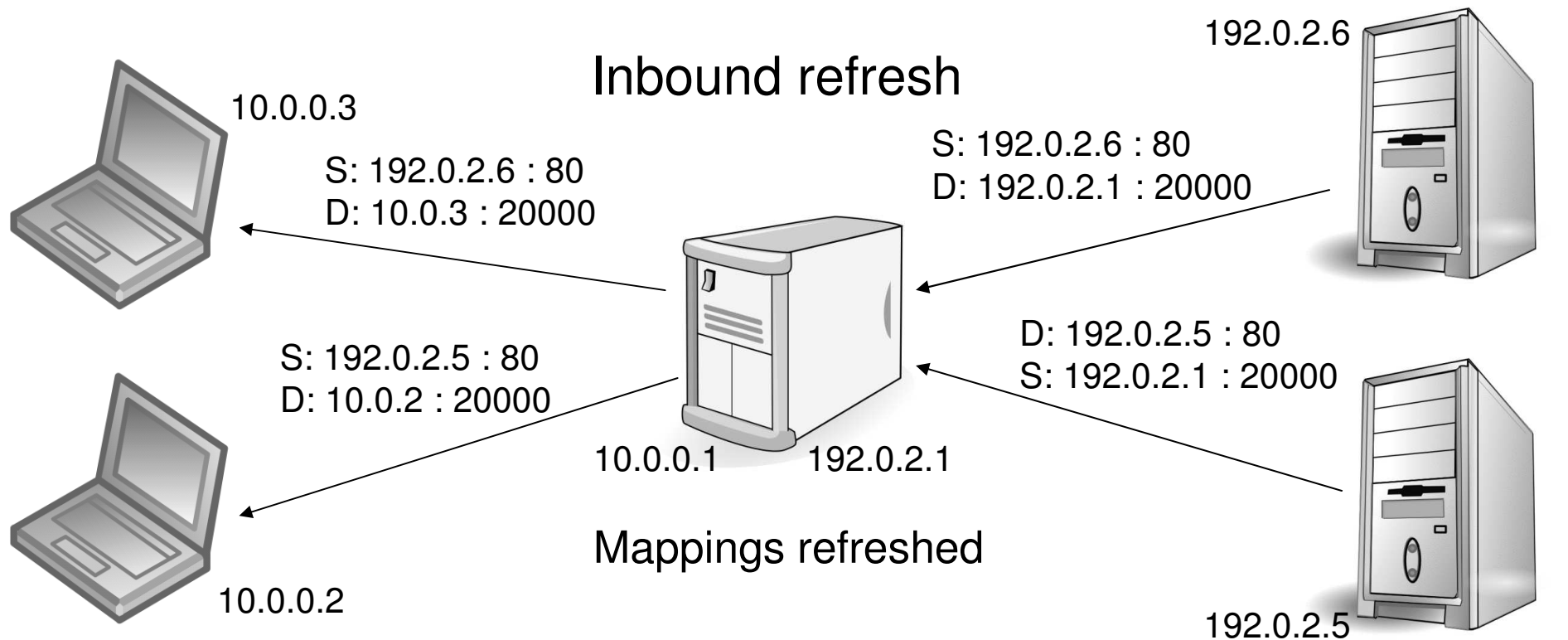
Mapping Refresh

- § NAT outbound refresh: packets from the internal to the external interface
 - MUST be used
- § NAT inbound refresh: packets from the external to the internal interface
 - Attackers may keep the mapping from expiring
 - MAY be used



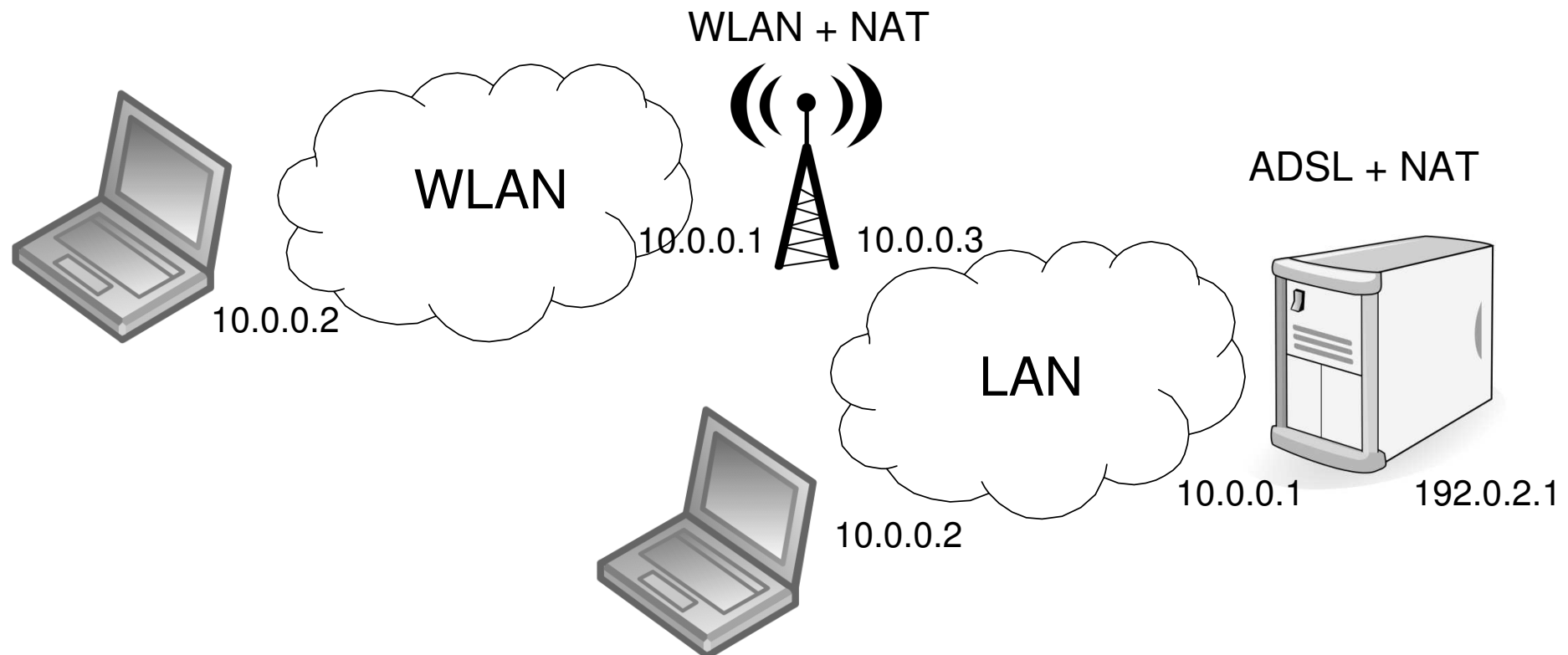
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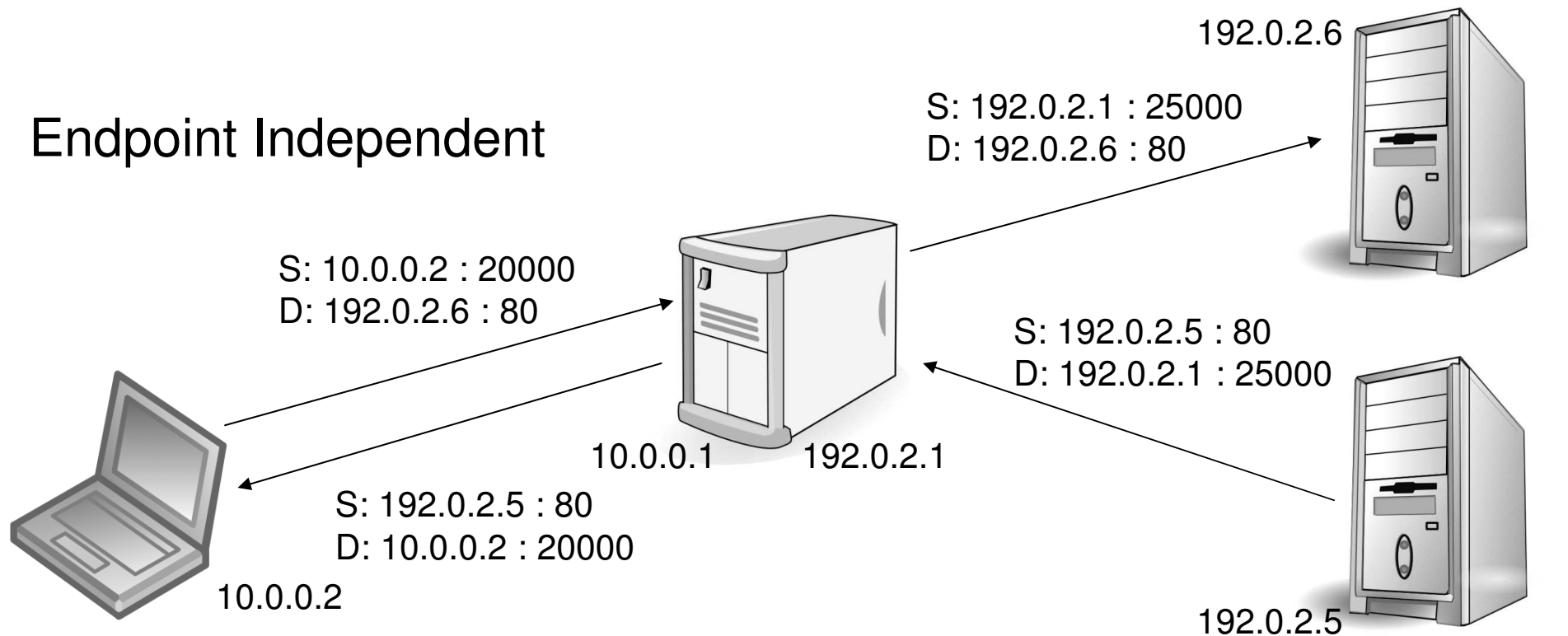
External Address Spaces

- § NATs MUST be able to handle external address spaces that overlap with the internal address space
- Internal nodes cannot communicate directly with external nodes that have the same address as another internal node
 - However, they can use STUN techniques



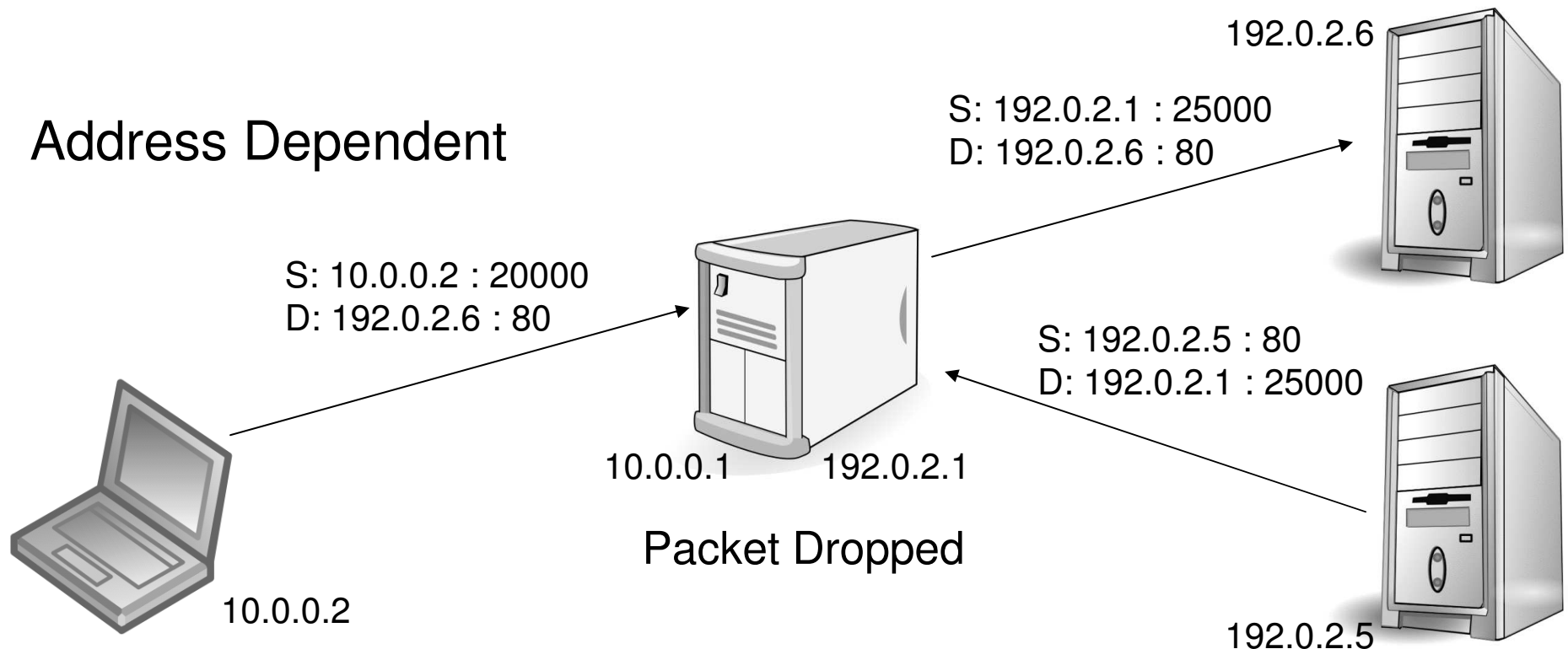
Filtering Behavior

- § Endpoint independent: any packets allowed back
- § Address dependent: external hosts can return packets
- § Address and port dependent
 - Packets sent to an address + port
 - Incoming packets allowed only from that address + port



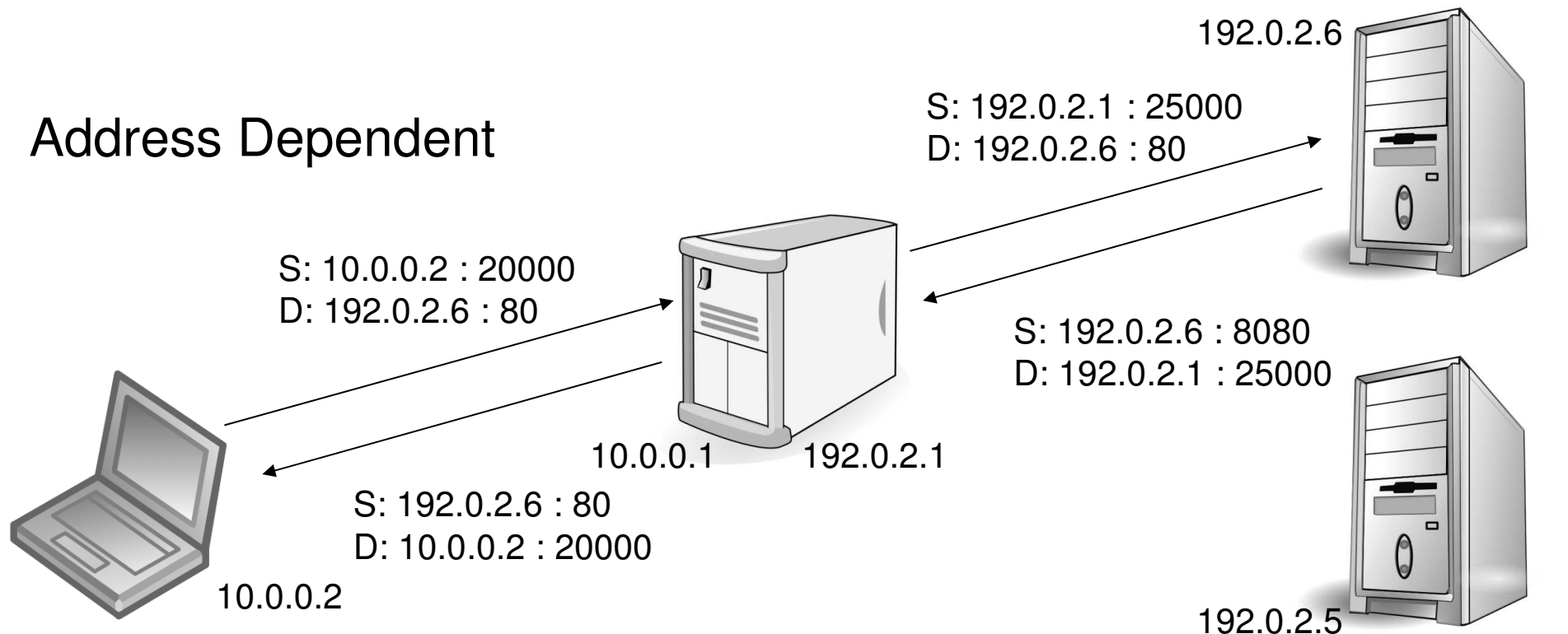
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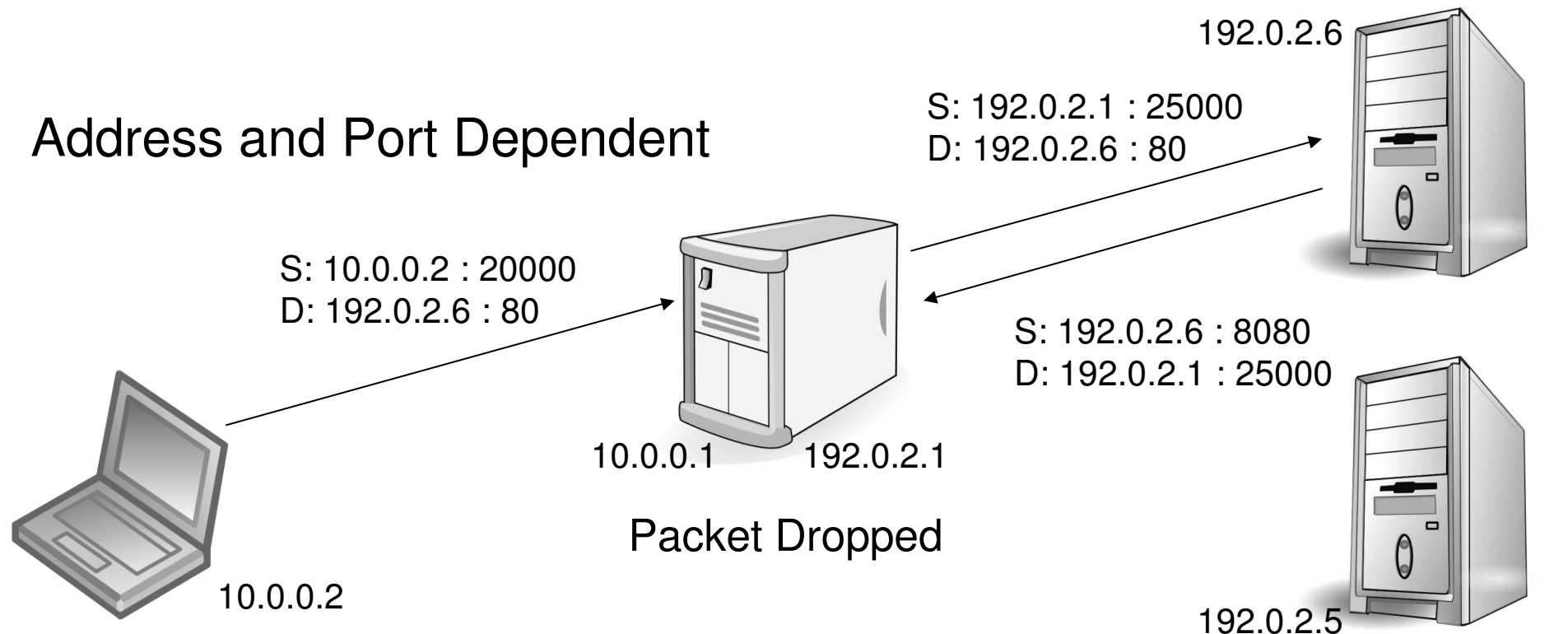
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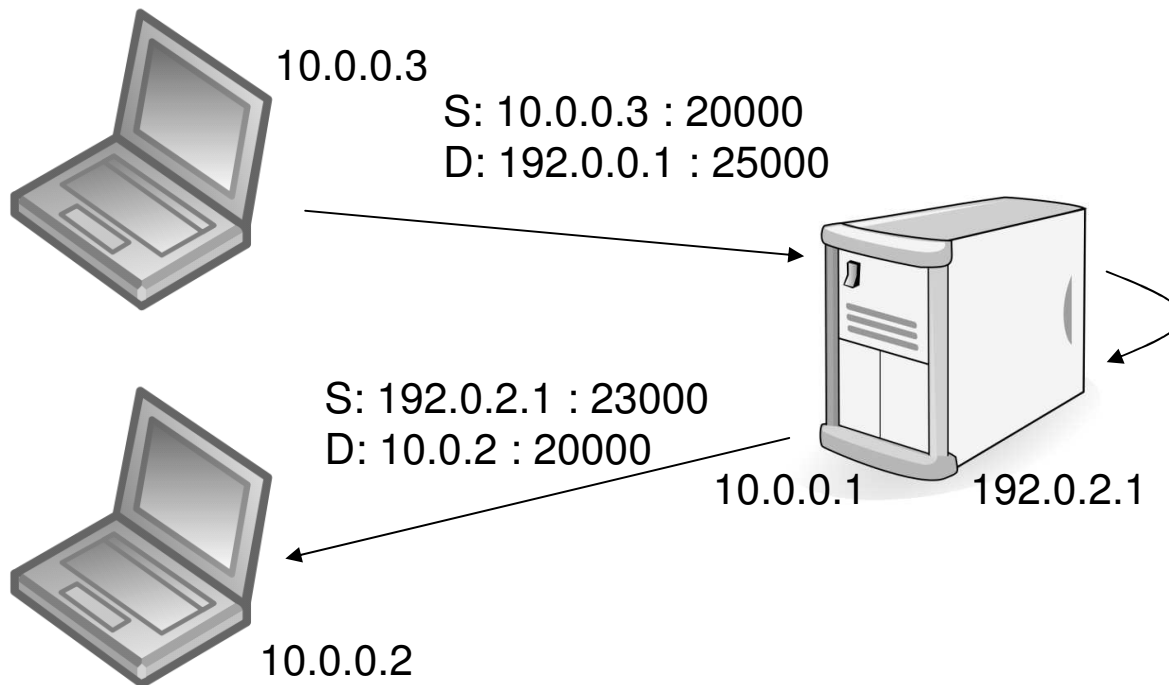


Filtering Behavior

- § Endpoint independent filtering is RECOMMENDED
 - Opens up ports for attackers
- § If a more stringent filtering is required
 - Address dependent filtering is RECOMMENDED

Hairpinning

- § Internal hosts communicate using external addresses
 - MUST be supported



Fragmented Packets

- § Receive fragments ordered
 - Only able to receive fragments in order
- § Receive fragments out of order
 - MUST be supported
 - As long as DoS attacks do not compromise the NAT's ability to process in order fragments and unfragmented packets
- § Receive fragments none

Various

- § ALGs SHOULD be turned off
 - MAY interfere with UNSAF methods

- § NATs MUST be deterministic
 - NATs MUST NOT change their mapping or filtering behavior

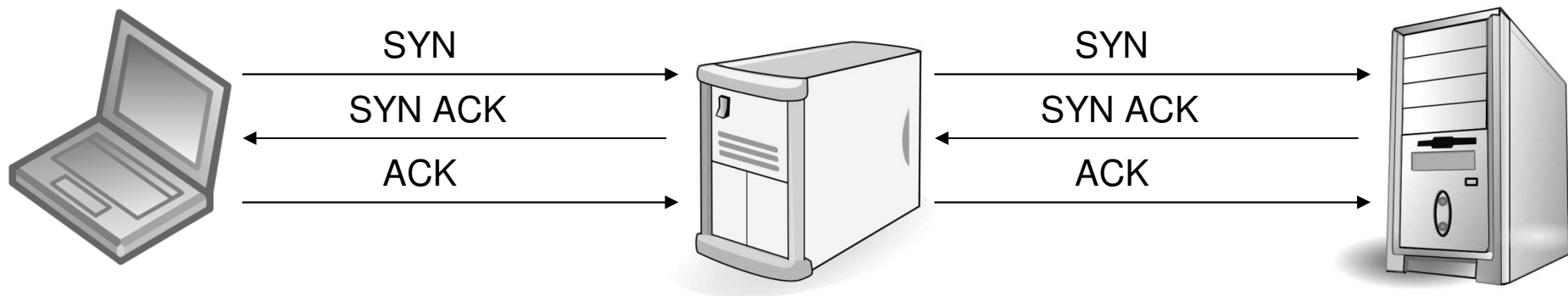
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Connection Establishment

- § Three-way handshake
 - MUST be supported
- § Simultaneous open
 - MUST be supported

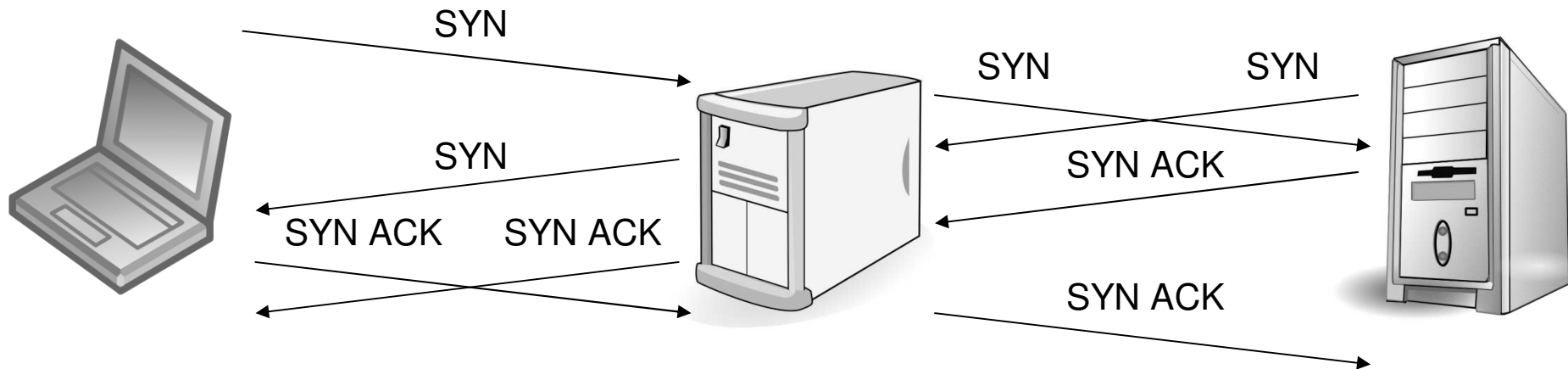
Three-way Handshake



Connection Establishment

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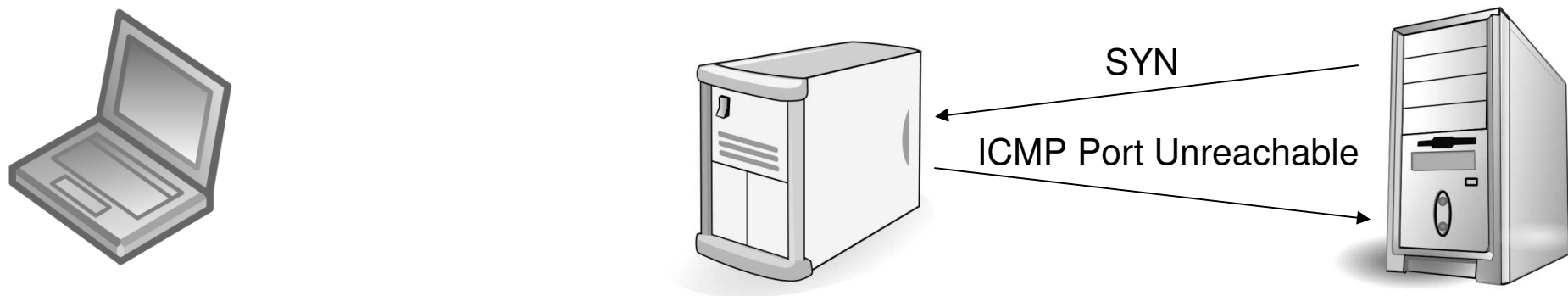
Simultaneous Open



Filtering External SYNs

- § Incoming SYN on the external interface.
 - No existing mapping for it
- § If the SYN is an error, the NAT should generate an ICMP error message
- § If it is a simultaneous open, the NAT should silently drop the packet

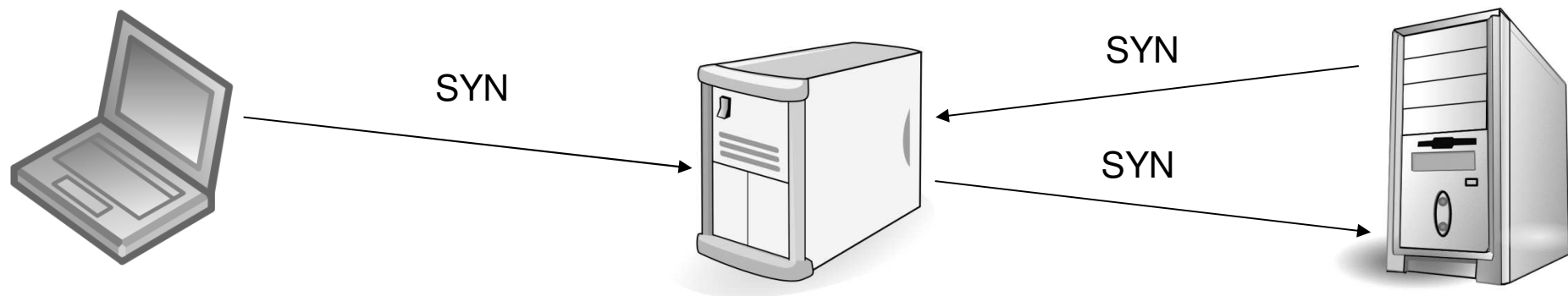
Error Situation



Filtering External SYNs

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Simultaneous Open



Filtering External SYNs

- § RECOMMENDED to respond to unsolicited SYNs
 - With an ICMP error with a delay of at least 6 seconds
 - If a matching outbound SYN is received, cancel the sending of the ICMP error

NAT Session Timeout

- § Established connections
 - MUST NOT be less than 2 hours and 4 minutes
 - TCP sends keep alives every 4 hours
- § Partially opened or partially closed connections
 - MUST NOT be less than 4 minutes
- § TIME_WAIT timeout not specified

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UNSAF

- § Unilateral self-address fixing
- § An application needs to refer to itself with a valid IP address
- § Problems
 - Reflectors may be in a different domain than the destination
 - Solutions may circumvent security
 - NAT behavior is not deterministic
 - Midcom has not been successful
- § Solutions need to have an exit strategy and a limited scope
- § Final goal is to get rid of NATs

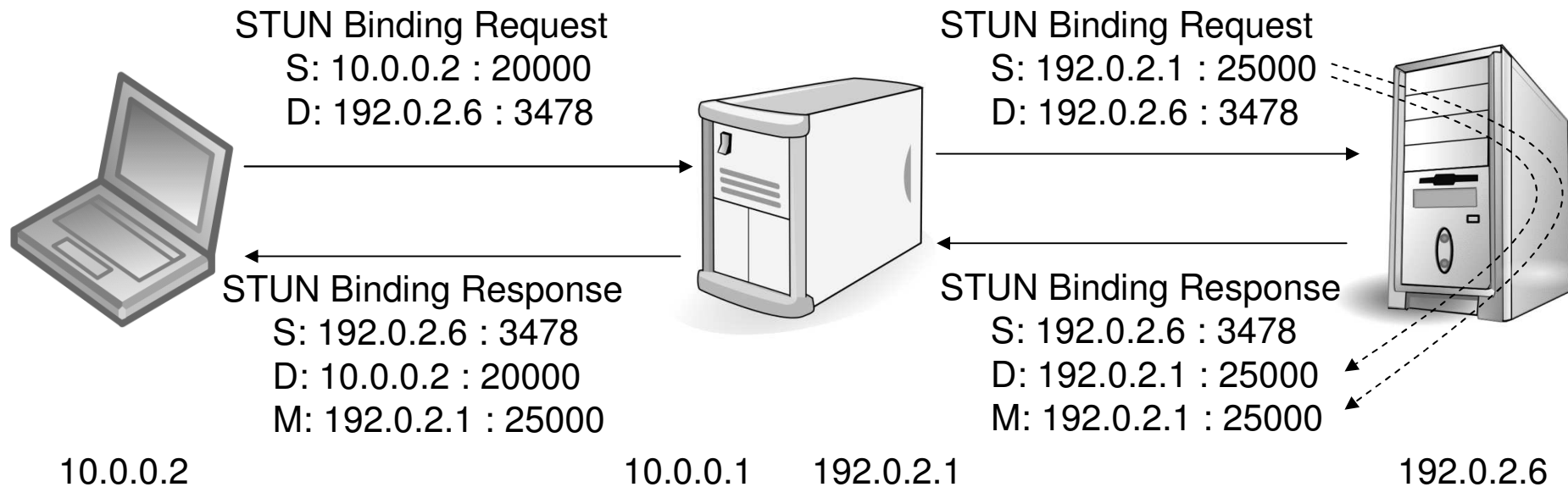
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Introduction to STUN

- § Originally a protocol between endpoints and reflectors
- § Revised specification defines usages
 - Binding discovery
 - NAT keep-alives
 - Short-term password
 - Relay (previously known as TURN)
- § TLV encoded
- § Can run on UDP, TCP, or TLS/TCP
- § STUN server located using DNS SRV
- § Transactions
 - Request/response
 - Indications (not delivered reliably)
- § Can be multiplexed with other protocols
 - Two first bits are zeros
 - Magic cookie
 - FINGERPRINT attribute

Binding Discovery



XOR-MAPPED-ADDRESS

- § In addition to the MAPPED-ADDRESS attribute
- § Some NATs inspect packets and translate IP addresses known to them

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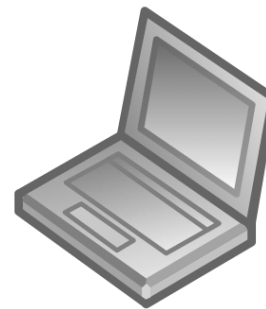
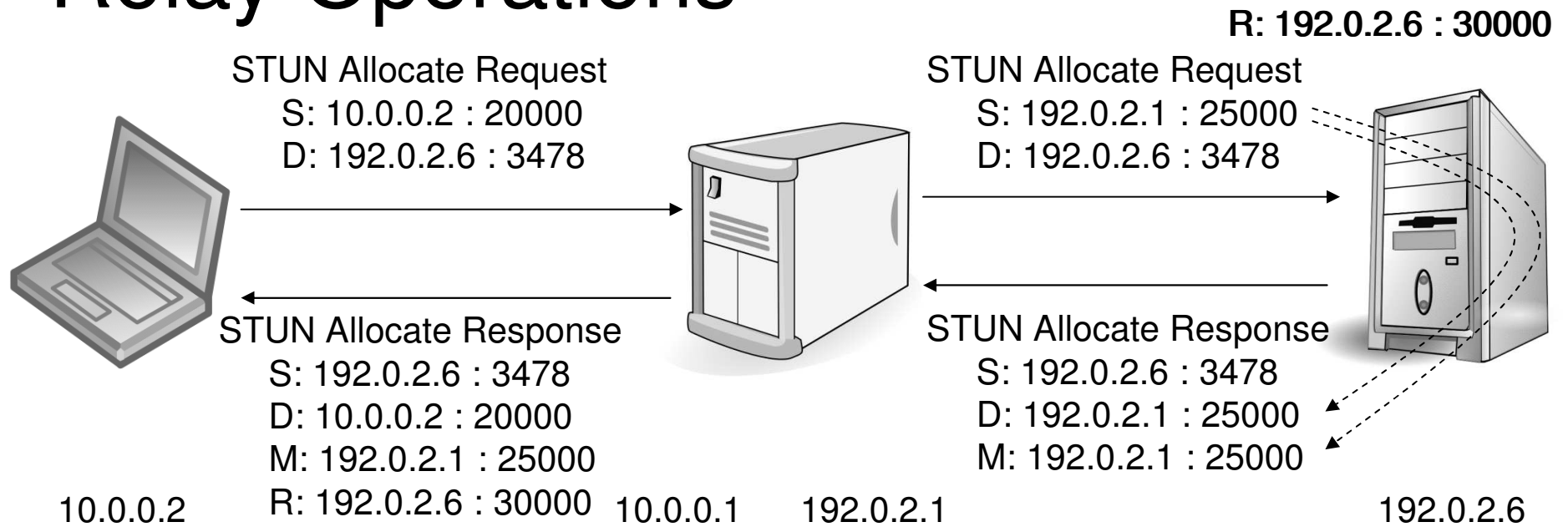
Introduction to STUN Relay Usage

- § Previously known as TURN
- § Allocate request / response
 - Allocate an external address at the relay
 - Responses carry a MAPPED-ADDRESS
- § Send indication
 - Send data to a remote endpoint through the relay
- § Data indication
 - Data received from remote endpoints through the relay
- § Set Active Destination request / response
 - Send and receive data to and from a single remote endpoint without using Send and Data wrappers
- § Connect request / response
 - Requests the relay to establish a TCP connection with the remote endpoint
- § Connection Status Indication
 - The relay informs the endpoint about the status of a TCP connection with the remote endpoint
 - LISTEN, ESTABLISHED, CLOSED

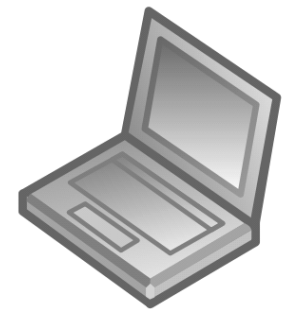
Transport Protocols

- § STUN relay clients can use UDP, TCP, or TLS/TCP
 - UDP to TCP is not supported

Relay Operations

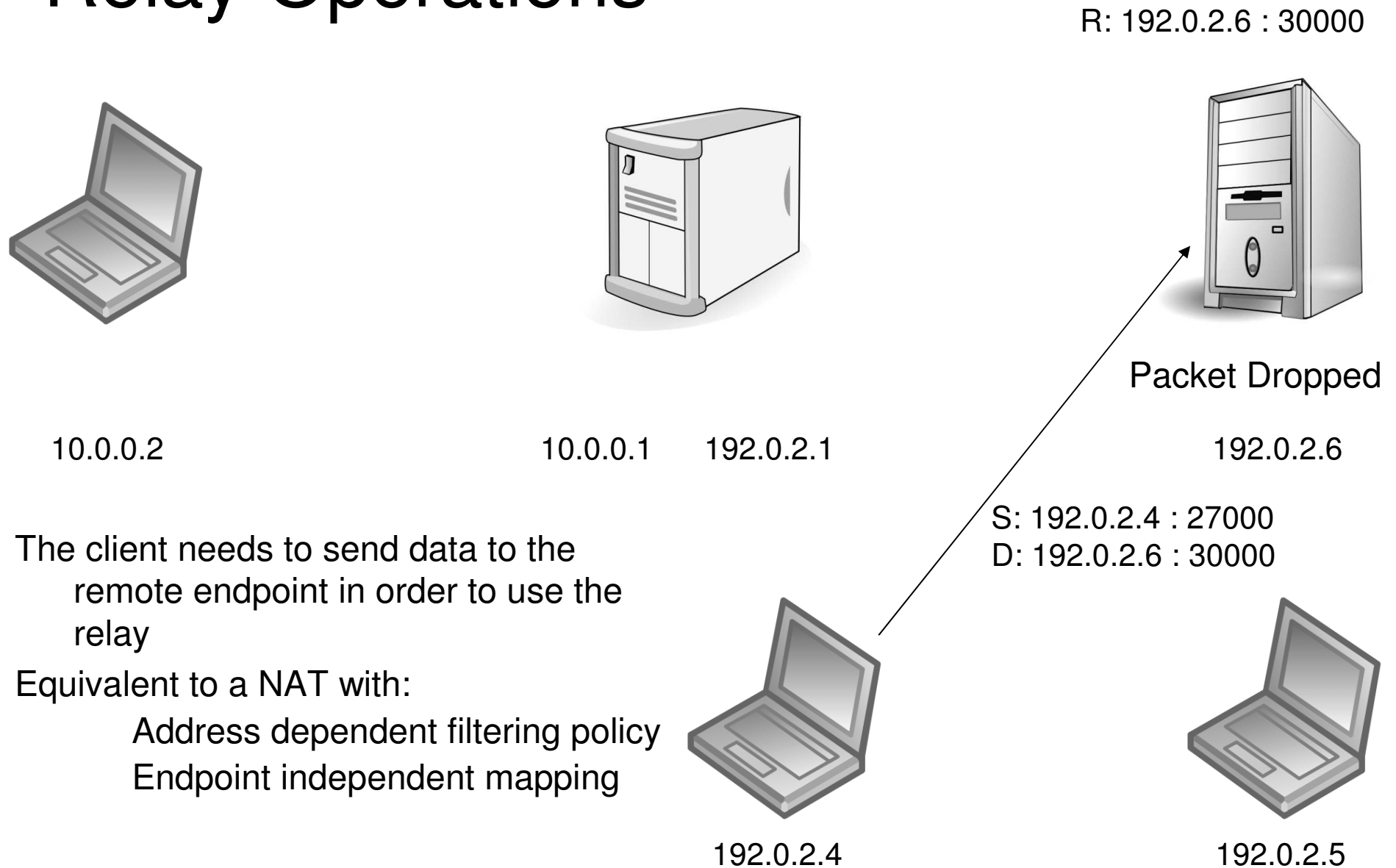


192.0.2.4



192.0.2.5

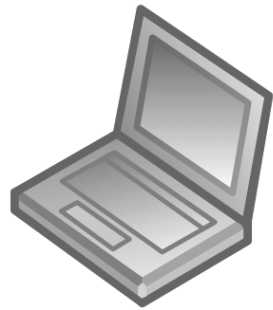
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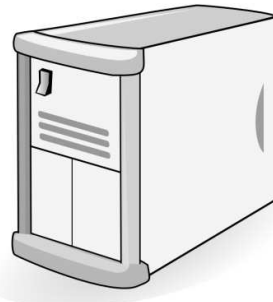
The client needs to send data to the remote endpoint in order to use the relay

Equivalent to a NAT with:
Address dependent filtering policy
Endpoint independent mapping

Relay Operations



10.0.0.2



10.0.0.1

192.0.2.1

R: 192.0.2.6 : 30000



Packet Dropped

192.0.2.6

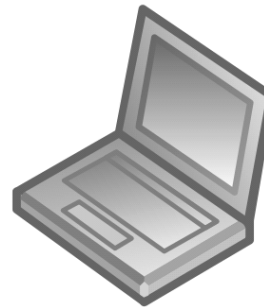
S: 192.0.2.5 : 27000

D: 192.0.2.6 : 30000

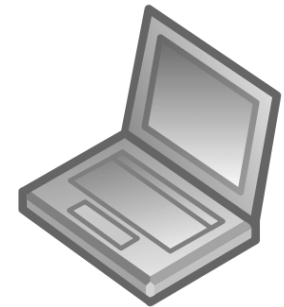
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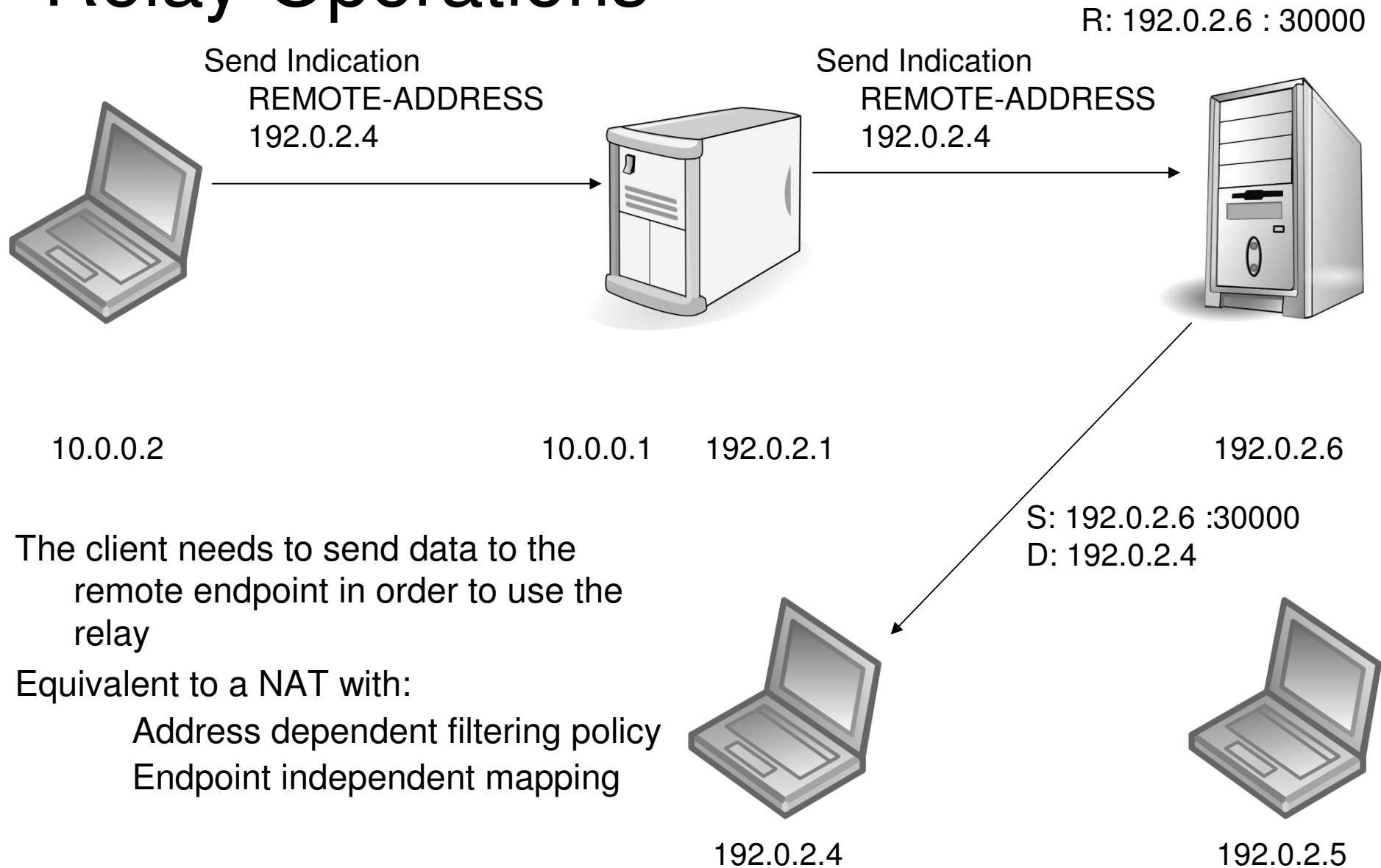


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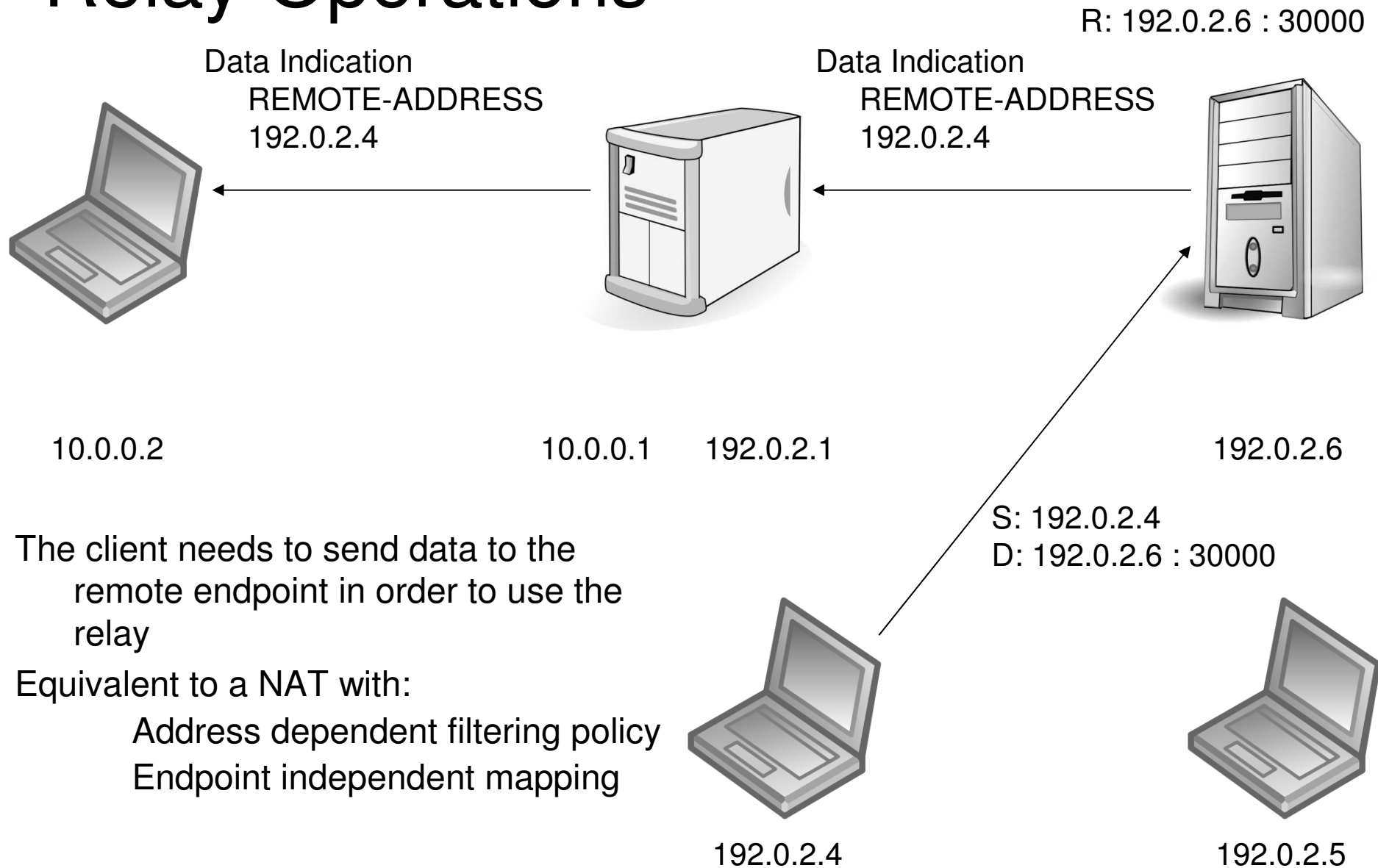


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Relay Operations



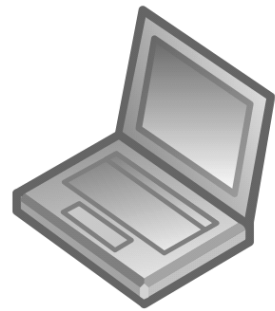
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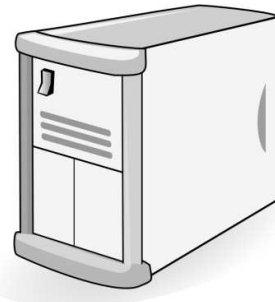
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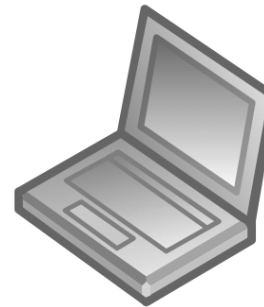
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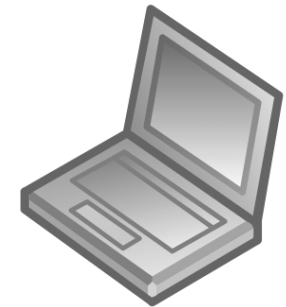
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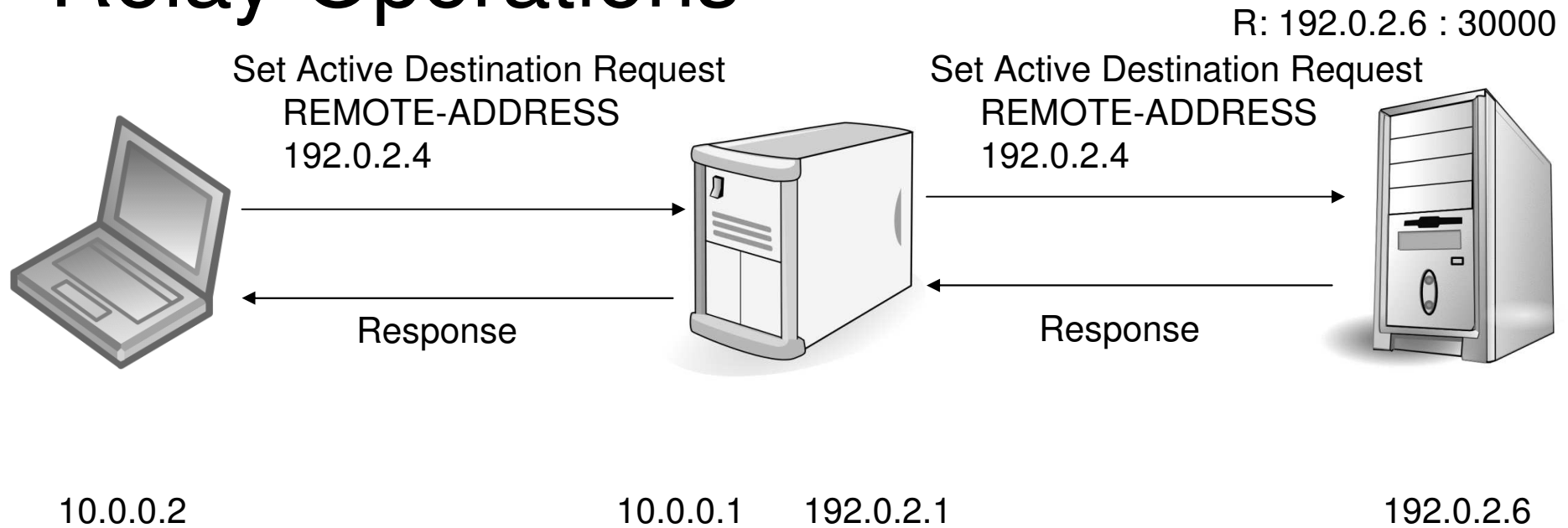


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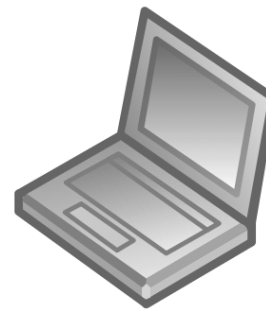


The client needs to set a destination as active in order to receive and send data to it without using indications.

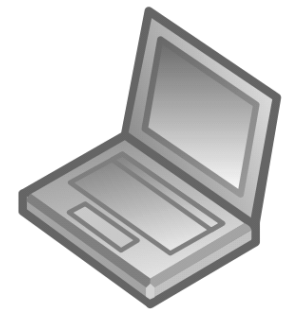
Data is framed:

Type: end-to-end or STUN

Length

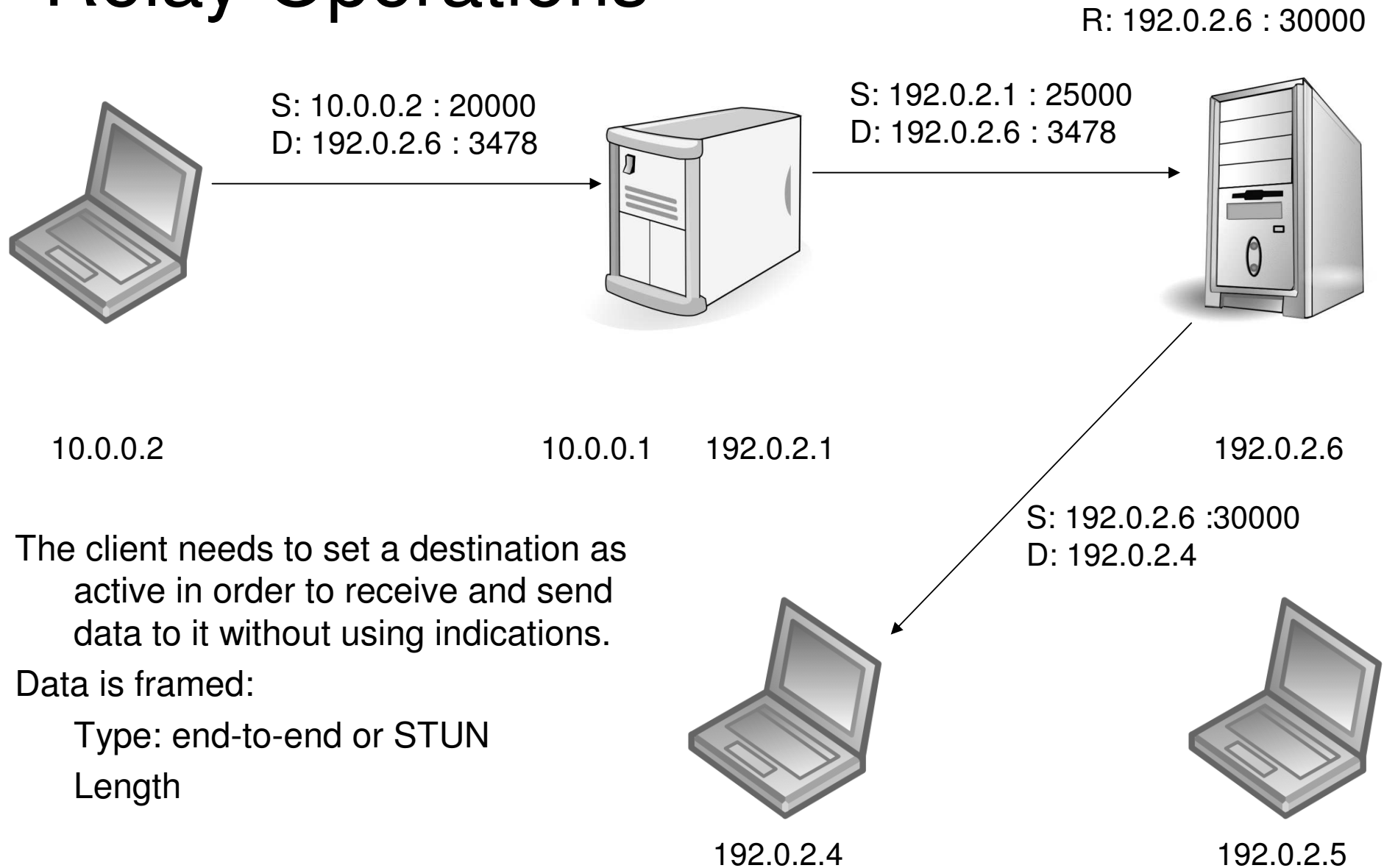


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Relay Operations



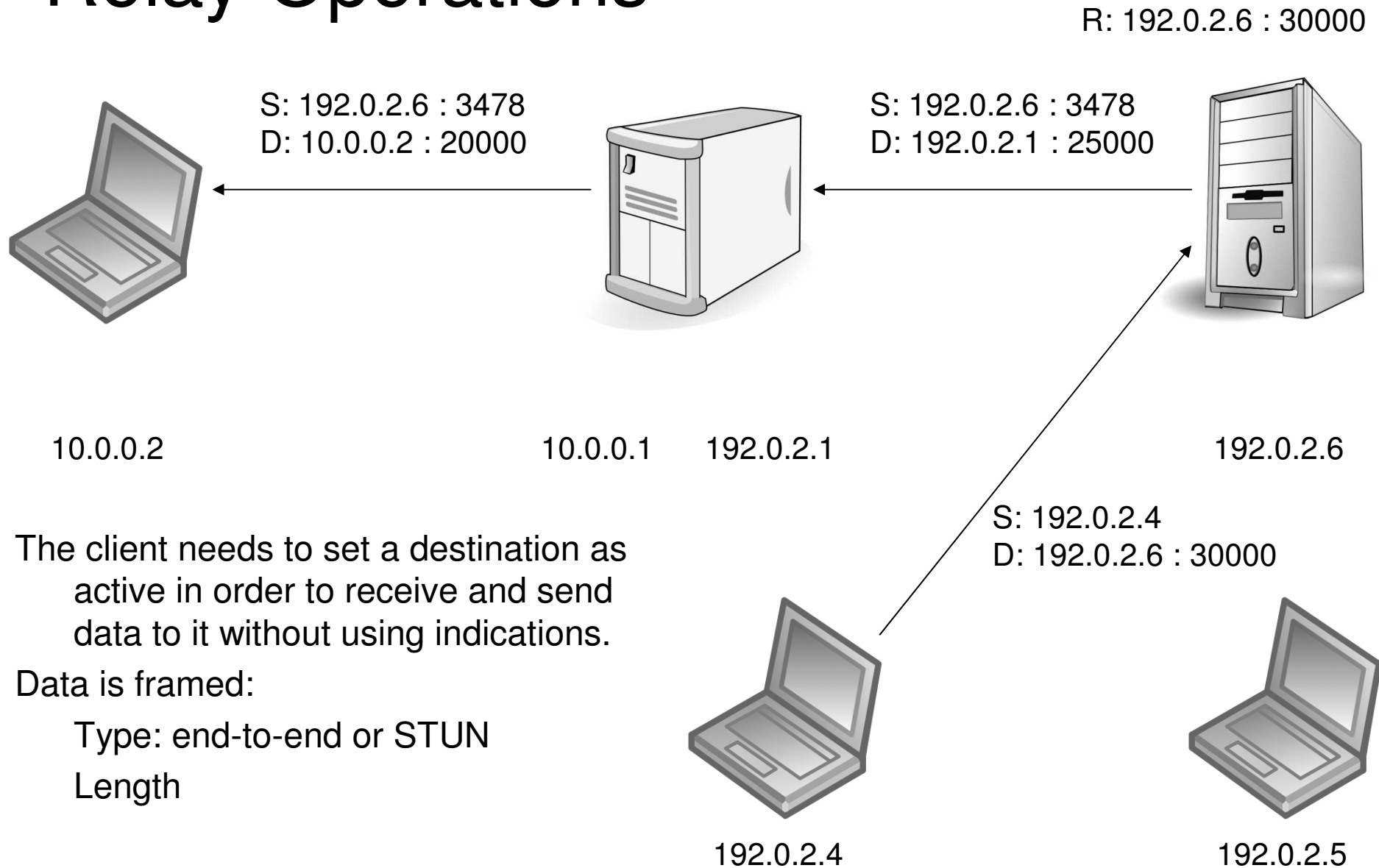
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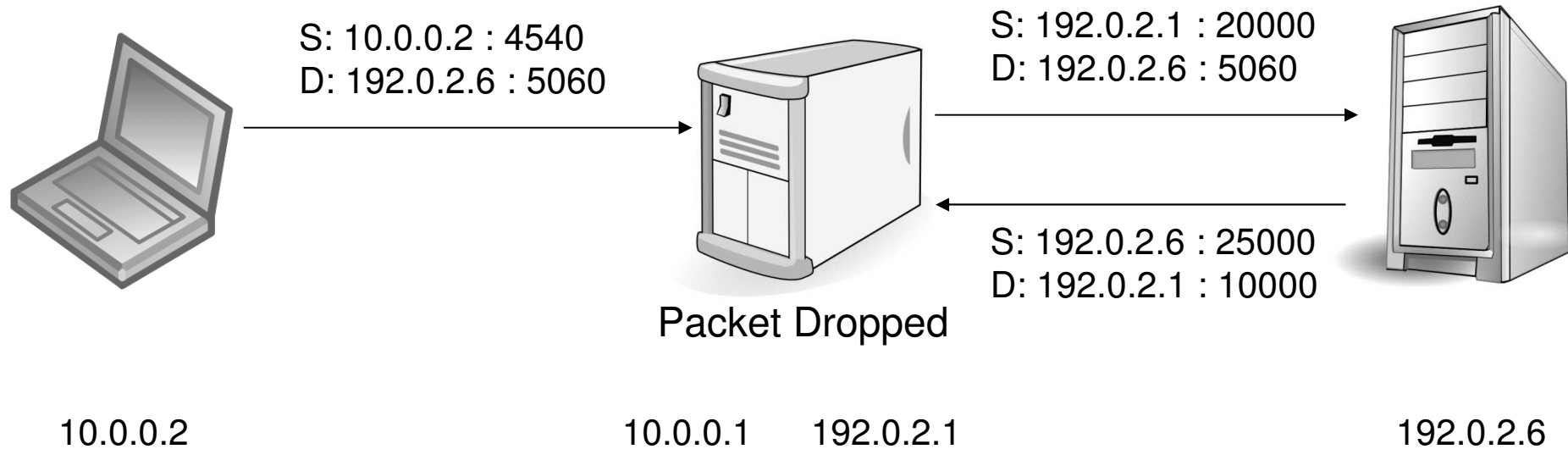
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Regular SIP Response Routing

- § Responses are sent to
 - Request's source IP address
 - Port in the Via entry
- § Responses are sent from
 - Typically the same IP address as the request was sent to
 - Any port

Regular SIP Response Routing



§ Request

- Via: SIP/2.0/UDP 10.0.0.2:10000;branch=z9hG4bKkjshdyff

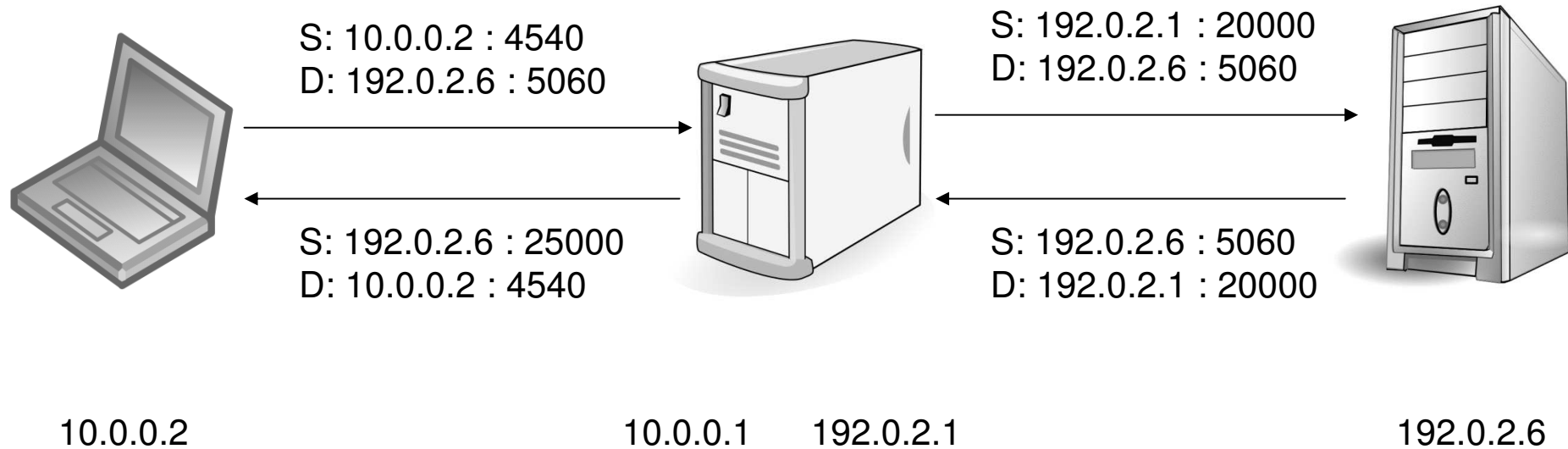
§ Response

- Via: SIP/2.0/UDP 10.0.0.2:10000;
received=192.0.2.1;branch=z9hG4bKkjshdyff

Symmetric SIP Response Routing

- § Responses are sent to
 - Request's source IP address
 - Request's source port
- § Responses are sent from
 - IP address the request was sent to
 - The port the request was sent to

Symmetric SIP Response Routing



§ Request

- Via: SIP/2.0/UDP 10.0.0.2:4540;
rport;branch=z9hG4bKkjshdyff

§ Response

- Via: SIP/2.0/UDP 10.0.0.2:4540;
received=192.0.2.1;rport=20000;branch=z9hG4bKkjshdyff

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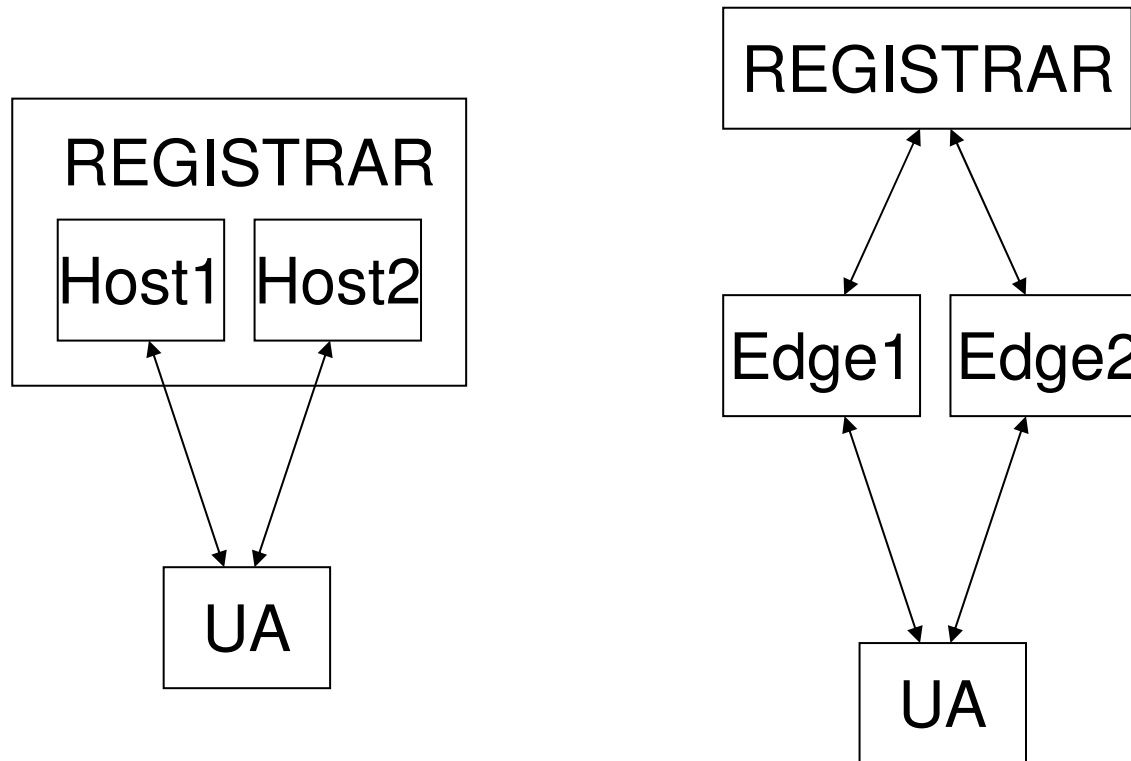
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 - User Agent Reachability
 - ICE

User Agent Reachability

- § Proxy to user agent connections
 - Problems with NATs and firewalls
- § User agent always starts the connection
- § Keep alives to keep NAT bindings up
 - Different alternatives analyzed
 - STUN
 - § Multiplexed with SIP
 - § New registration if a new binding is discovered
- § User agent identification (e.g., for service profiles)
 - Instance identifier

Reliability and Scalability

§ Different flows identified by their registration IDs



Outline

- § Introduction to NATs
- § NAT Behavior (actual and recommended)
 - UDP
 - TCP
- § IAB UNSAF Considerations
- § STUN
- § STUN Relay Usage
- § NAT Traversal in SIP
 - SIP Response Routing
 - User Agent Reachability
 - ICE

Introduction to ICE

- § Endpoints gather all the addresses they can
- § They run connectivity checks between them
- § They choose the highest priority pair that works

Gathering Addresses

§ Address types

- Host candidates
- Server-reflexive candidates
- Relayed candidates
- Peer-reflexive candidates

§ A candidate's base: the address used to send data

- The base for a reflexive candidate is a host candidate

§ Duplicated addresses are removed

- Candidates with the same transport address but different base are considered different

§ Foundation: used to freeze addresses (related to connectivity checks)

- Same type
- Bases with the same IP address
- Same STUN server

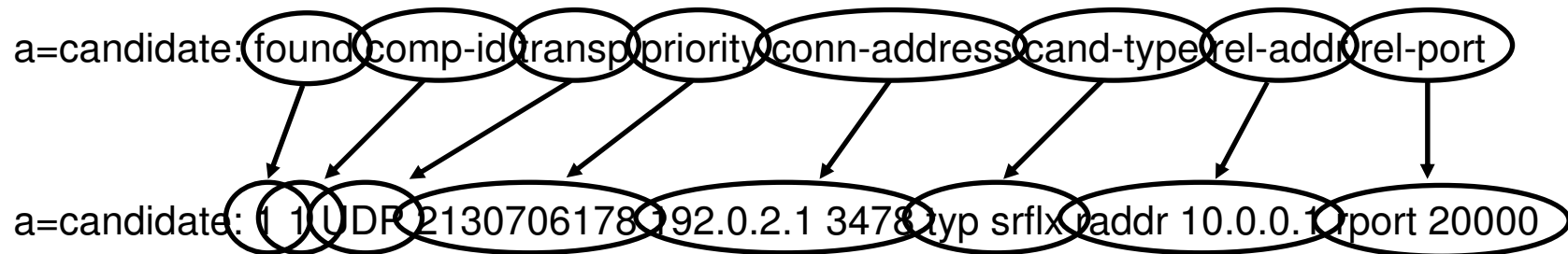
Prioritizing Addresses

$$\text{Priority} = 2^{24} (\text{type preference}) + 2^8 (\text{local preference}) + 2 (256 - \text{component ID})$$

- § Type preference [0-126]: preference for the type of candidate (e.g., server reflexive)
- § Local preference [0-65535]: preference for the interface the candidate was obtained from (e.g., multihomed hosts)
- § Component ID [1-256]: 1 for RTP and 2 for RTCP

Generation of an offer

- § The candidate with the highest probability to work goes into the m and c lines
 - A relayed address initially
- § The rest of the candidates go into 'candidate' attributes



- § User name and password for connectivity checks in 'ice-frag' and 'ice-pwd' attributes
 - tcp-so
 - tcp-act
 - tcp-pass
- Other transports are:
Related address: base

Prioritizing Pairs

§ After the offer/answer exchange

$$\text{Pair Priority} = 2^{32} \text{ MIN}(O-P, A-P) + 2 \text{ MIN}(O-P, A-P) + (O-P > A-P : 1 ? 0)$$

§ O-P: priority in the offer

§ A-P: priority in the answer

Connectivity Checks

- § Five states for a pair:
 - Waiting, in progress, succeeded, failed, frozen
- § Periodic checks and triggered checks
 - Periodic checks performed in priority order
- § Connectivity is checked with STUN Binding Requests
 - Carry a concatenation of user names and the remote password

ICE Roles

§ Controlling agent

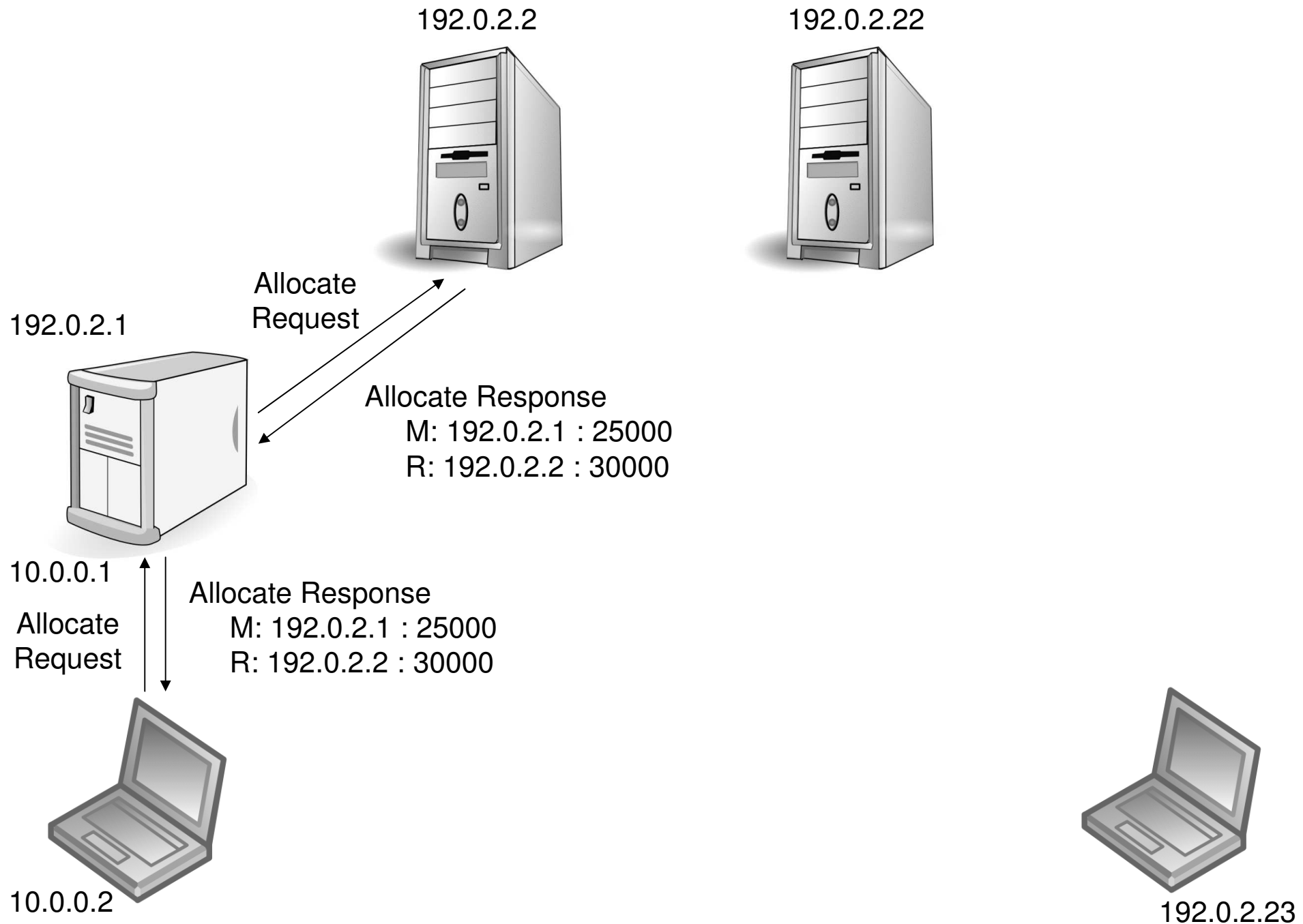
- Selects which pair to use
- USE-CADIDATE attribute
 - § Default algorithm: included in every check
- Agent that generates the initial offer

§ Passive-only agents

- They know they are not behind a NAT
 - § E.g., PSTN gateways, conferencing servers
- Include 'a=ice-passive' in their session descriptions
- Respond to checks
- Generate triggered checks
- Generate keepalives

ICE Example (1)

- § One endpoint is behind a NAT
- § One endpoint has a public IP address
- § Endpoints use STUN servers that support the relay usage



Host candidate:

10.0.0.2 : 20000

Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

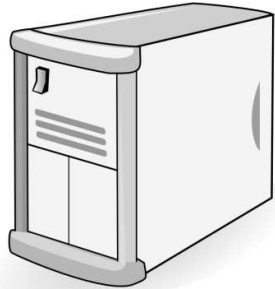
192.0.2.2



192.0.2.22



192.0.2.1

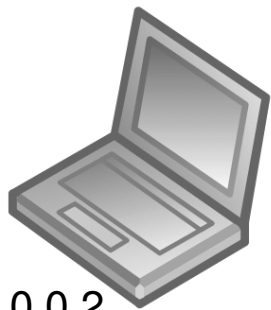


10.0.0.1

Allocate Response

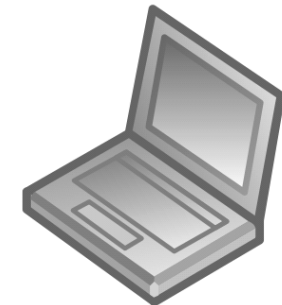
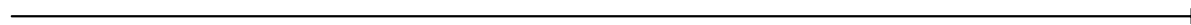
M: 192.0.2.1 : 25000

R: 192.0.2.2 : 30000



10.0.0.2

INVITE (offer)



192.0.2.23

Host candidate:

10.0.0.2 : 20000

Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

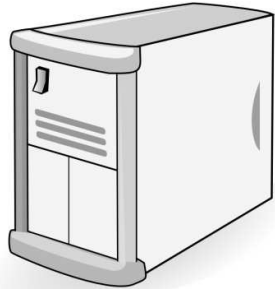
192.0.2.2



192.0.2.22



192.0.2.1



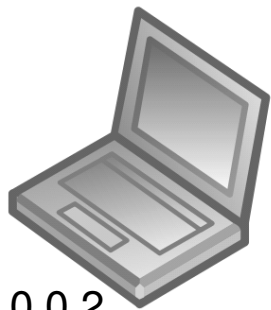
10.0.0.1

Allocate Request

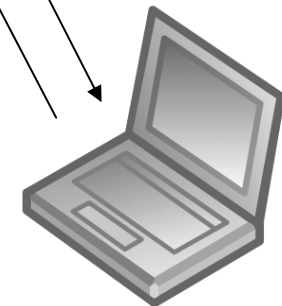
Allocate Response

M: 192.0.2.23 : 35000

R: 192.0.2.22 : 45000



10.0.0.2



192.0.2.23

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

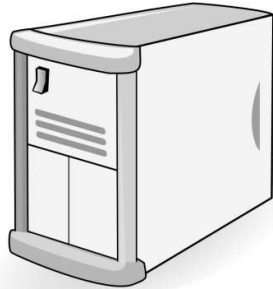


192.0.2.22



Host candidate:
~~192.0.2.23 : 35000~~
~~Server reflexive:~~
~~192.0.2.22 : 45000~~
Relayed:
192.0.2.22 : 45000

192.0.2.1

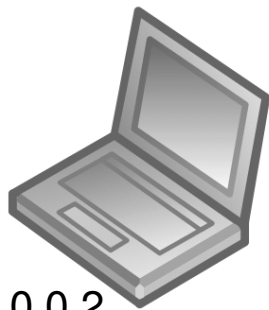


10.0.0.1

Allocate Response

M: 192.0.2.23 : 35000

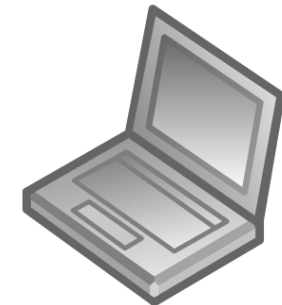
R: 192.0.2.22 : 45000



10.0.0.2

200 OK (answer)

ACK



192.0.2.23

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

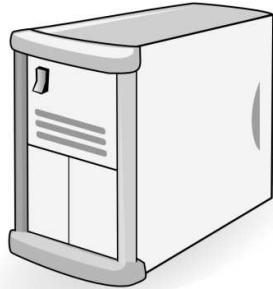


192.0.2.22



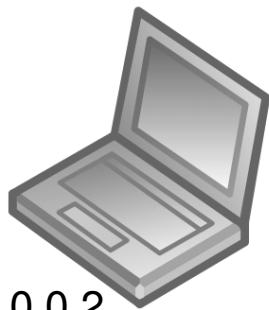
Host candidate:
192.0.2.23 : 35000
Relayed:
192.0.2.22 : 45000

192.0.2.1



10.0.0.1
Binding Request

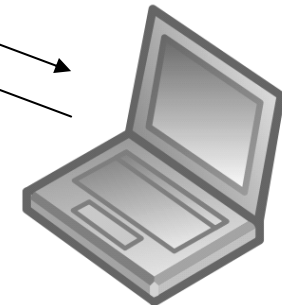
Binding Response
M: 192.0.2.1 : 25000



10.0.0.2

Binding Request

Binding Response
M: 192.0.2.1 : 25000



192.0.2.23

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

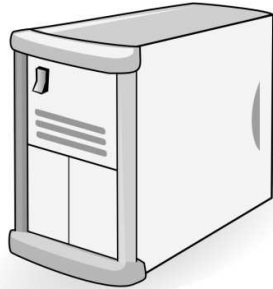


192.0.2.22



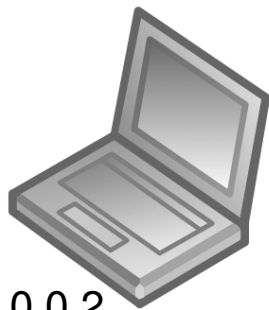
Host candidate:
192.0.2.23 : 35000
Relayed:
192.0.2.22 : 45000

192.0.2.1



10.0.0.1
Binding Request

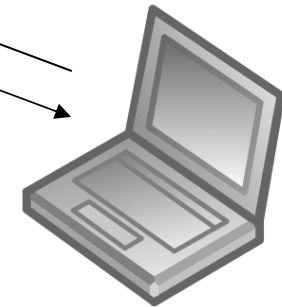
Binding Response
M: 192.0.2.23 : 35000



10.0.0.2

Binding Request

Binding Response
M: 192.0.2.23 : 35000



192.0.2.23

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

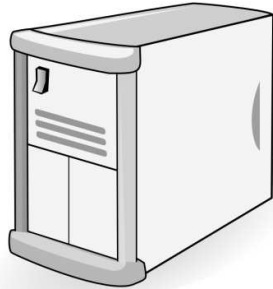


192.0.2.22

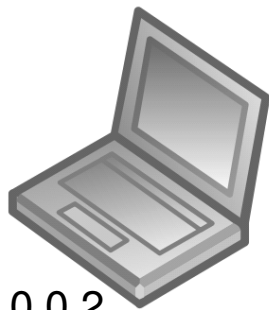


Host candidate:
192.0.2.23 : 35000
Relayed:
192.0.2.22 : 45000

192.0.2.1



10.0.0.1

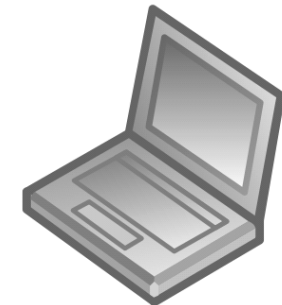


10.0.0.2

INVITE (offer)

200 OK (answer)

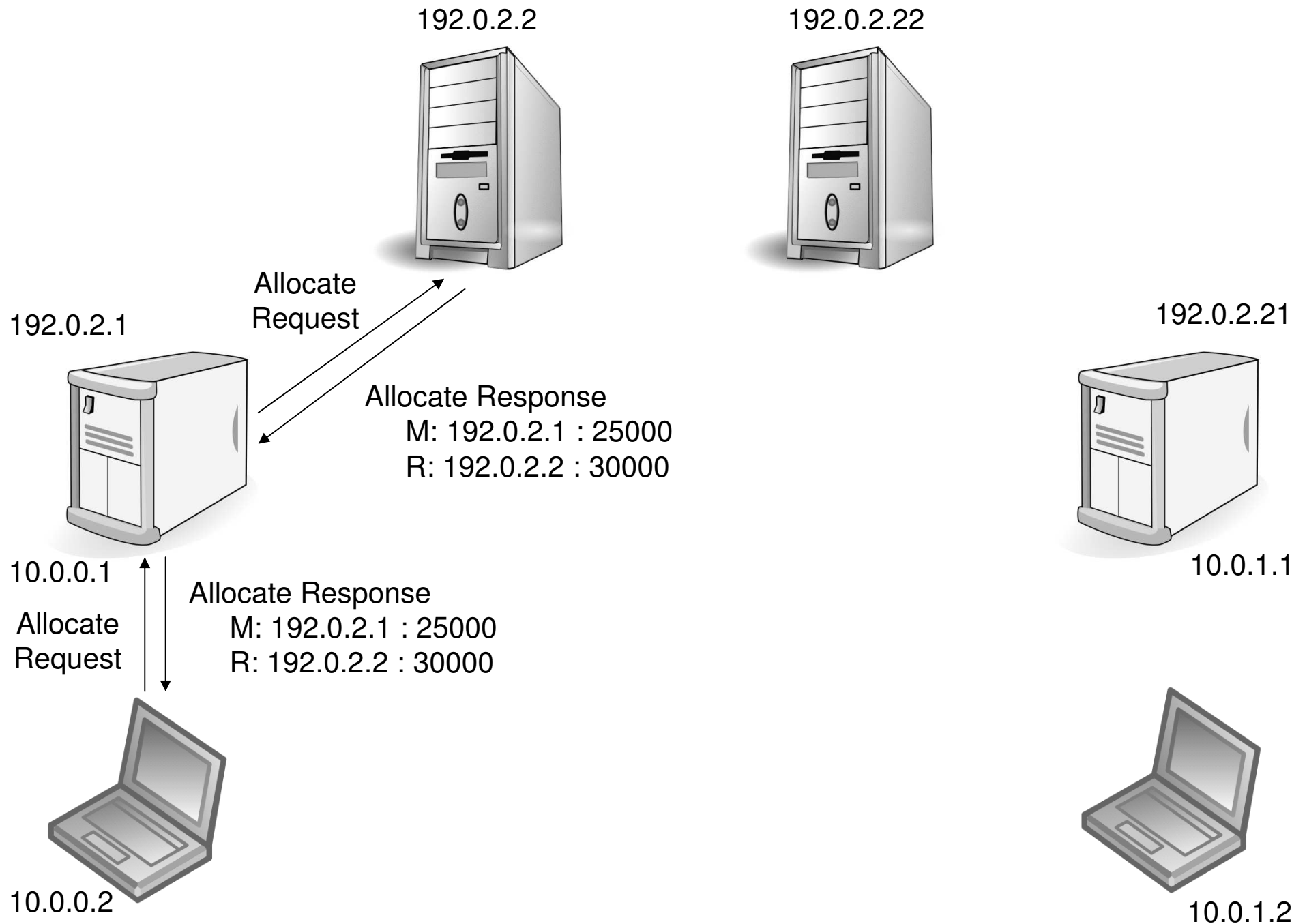
ACK



192.0.2.23

ICE Example (2)

- § Both endpoints are behind NATs
- § Endpoints use STUN servers that support the relay usage



Host candidate:

10.0.0.2 : 20000

Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

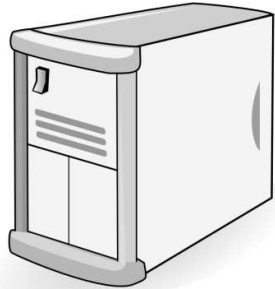
192.0.2.2



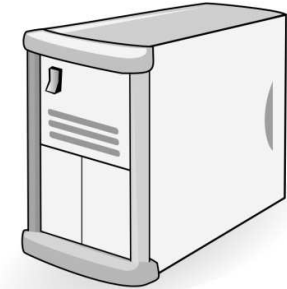
192.0.2.22



192.0.2.1



192.0.2.21



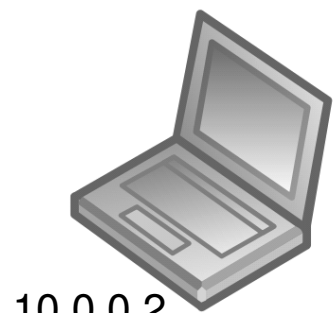
10.0.0.1

Allocate Response

M: 192.0.2.1 : 25000

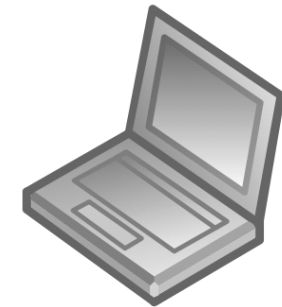
R: 192.0.2.2 : 30000

10.0.1.1



10.0.0.2

INVITE (offer)



10.0.1.2

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2



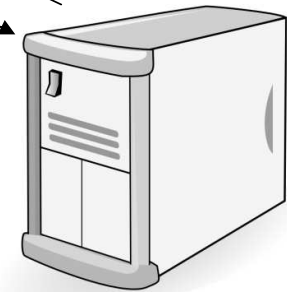
192.0.2.22



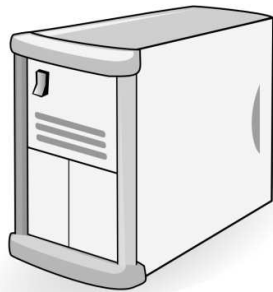
Allocate Request

Allocate Response
M: 192.0.2.21 : 25000
R: 192.0.2.22 : 30000

192.0.2.21



192.0.2.1

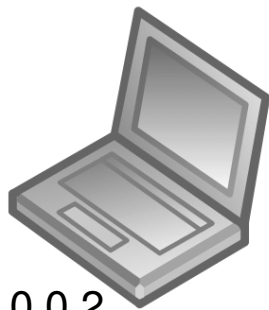
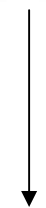


10.0.0.1

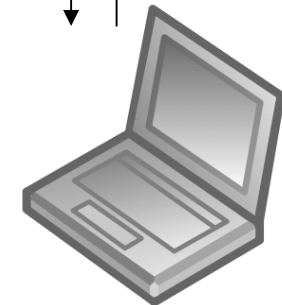
Allocate Response
M: 192.0.2.21 : 25000
R: 192.0.2.22 : 30000

10.0.1.1

Allocate Request



10.0.0.2



10.0.1.2

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

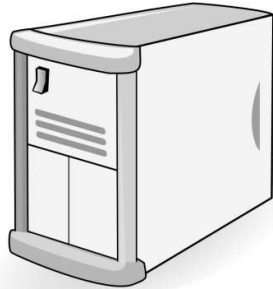


192.0.2.22



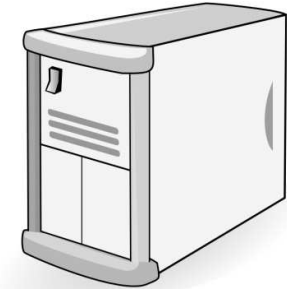
Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000

192.0.2.1



10.0.0.1

192.0.2.21

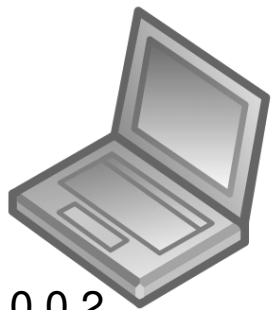


10.0.1.1

Allocate Response

M: 192.0.2.21 : 25000

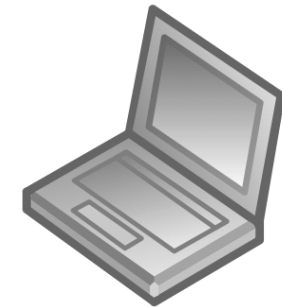
R: 192.0.2.22 : 30000



10.0.0.2

200 OK (answer)

ACK



10.0.1.2

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

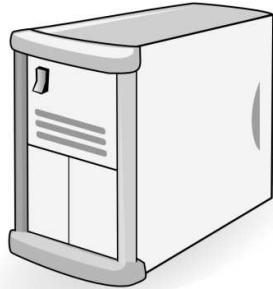


192.0.2.22



Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000

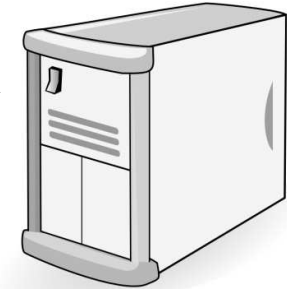
192.0.2.1



Binding Request

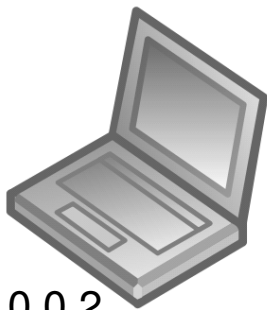
Packet Dropped

192.0.2.21



10.0.0.1

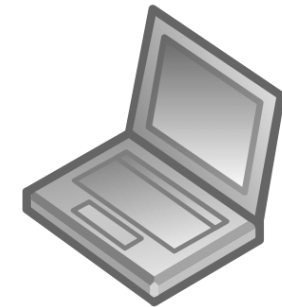
Binding Request



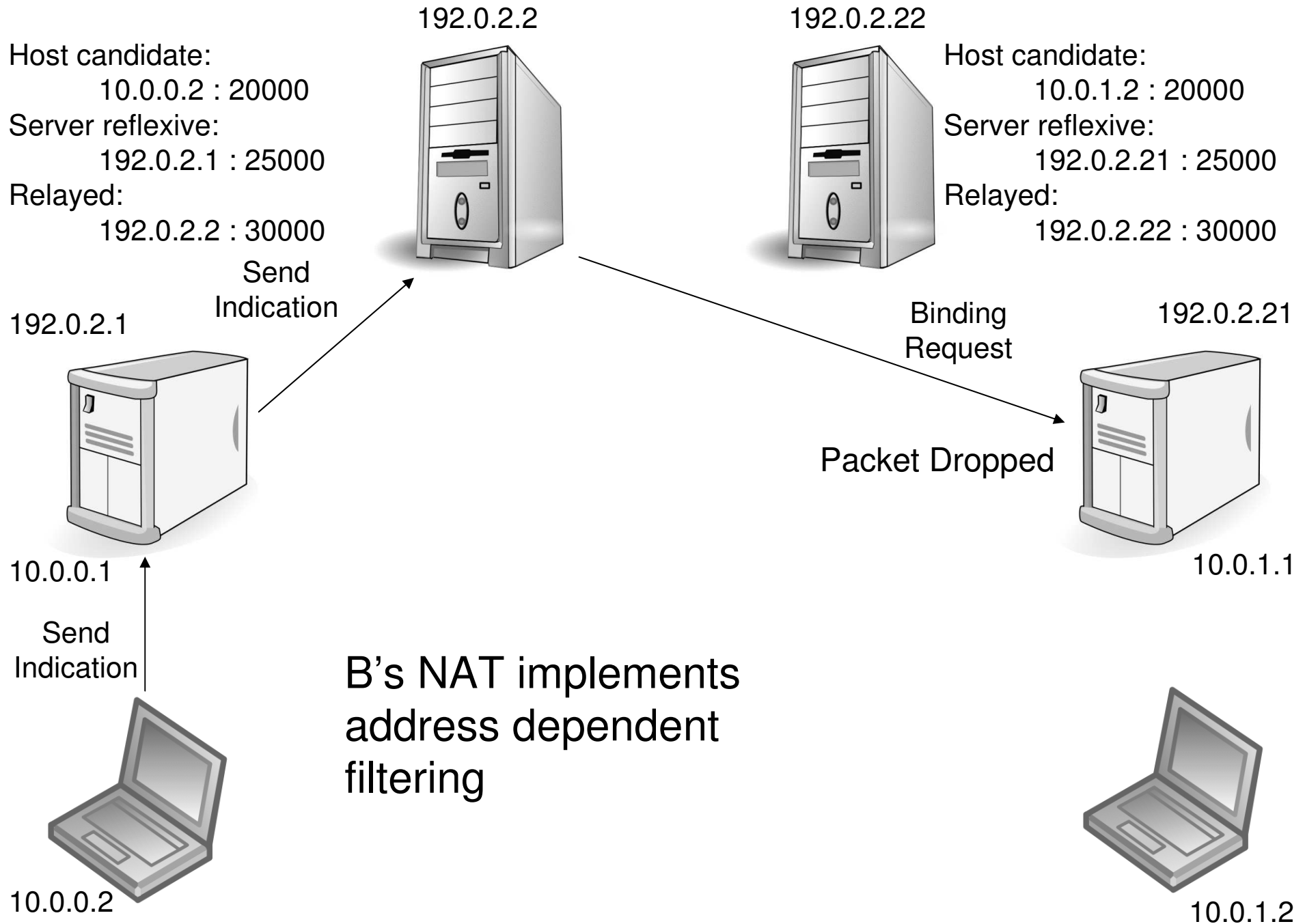
10.0.0.2

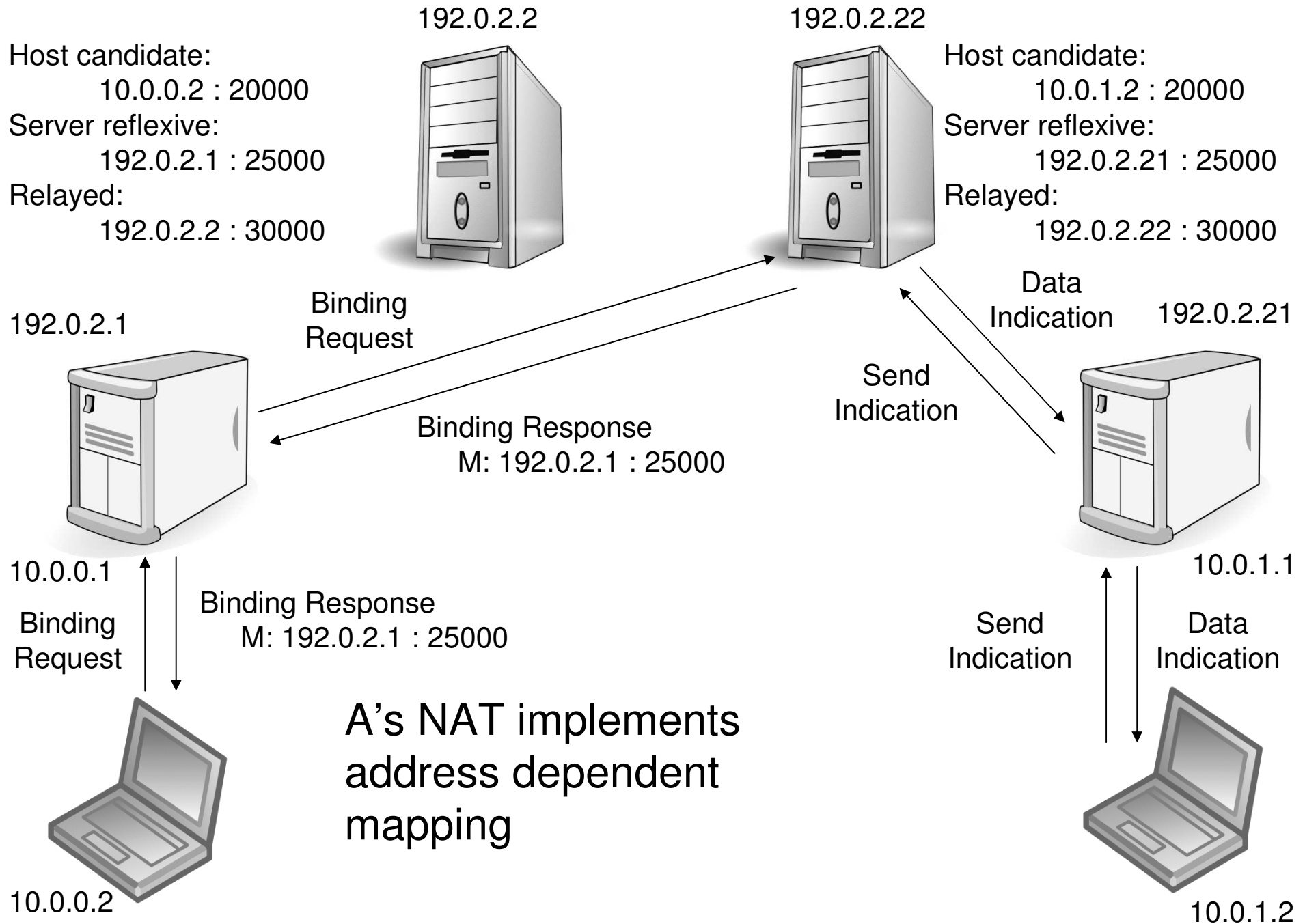
B's NAT implements address dependent filtering

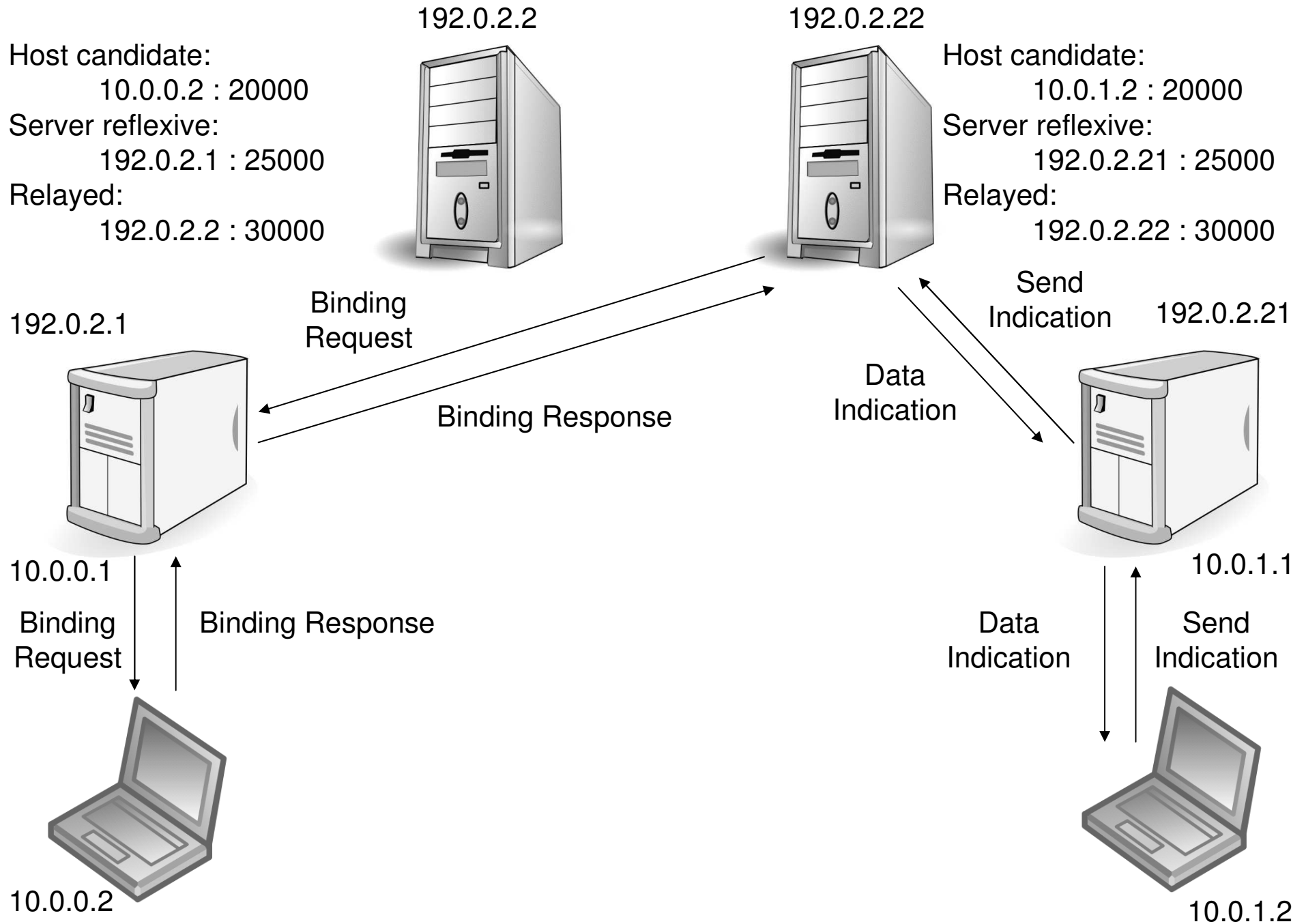
10.0.1.1



10.0.1.2







Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

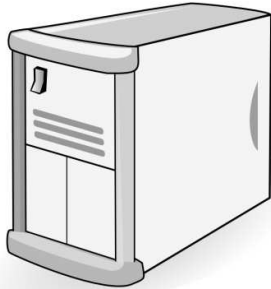


192.0.2.22



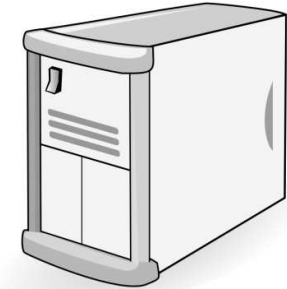
Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000

192.0.2.1

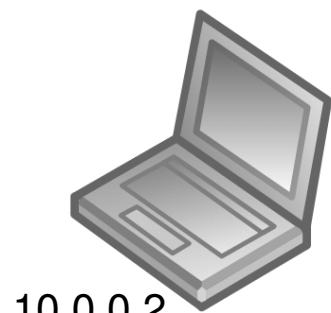


10.0.0.1

192.0.2.21



10.0.1.1

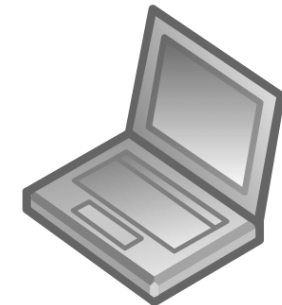


10.0.0.2

INVITE (offer)

200 OK (answer)

ACK



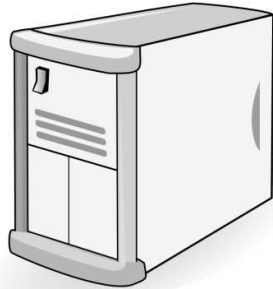
10.0.1.2

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

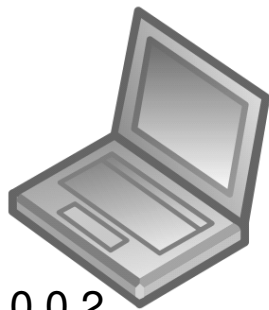
192.0.2.2



192.0.2.1



10.0.0.1

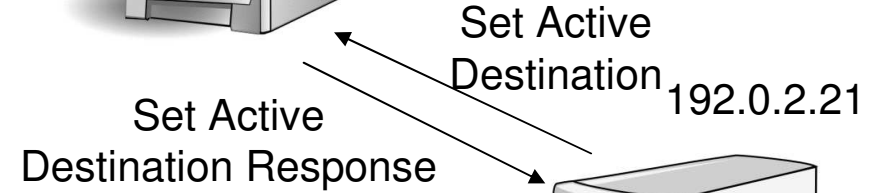


10.0.0.2

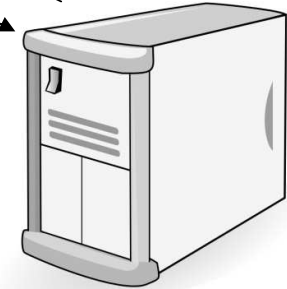
192.0.2.22



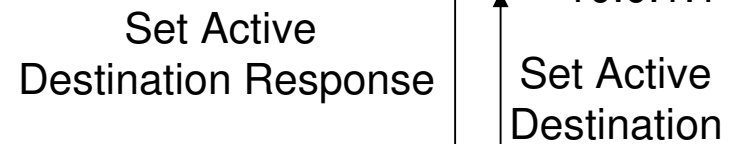
Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000



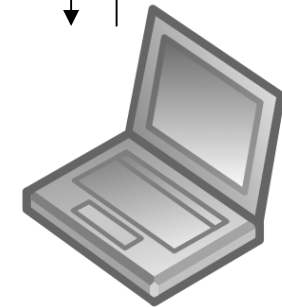
192.0.2.21



10.0.1.1



10.0.1.2



10.0.1.2

ERICSSON 

TAKING YOU FORWARD