Trust and threat evaluation

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Lecture topics

- Formal system evaluation
- Threat trees
- Penetration testing
- Malicious logic

Formal evaluation of systems

- Trust based on assurance *evidence*
 - basis for confidence
 - not perfect security
- Helps creating secure systems
- Evaluation methodology features
 - set of requirements for security functionality
 - set of assurance requirements to establish that system meets functional requirements
 - methodology for determining that system meets functional requirements based on analysis of the assurance evidence
 - measure for evaluation (*level of trust*)

What evaluation guarantees?

- The system is perfectly secure NOT!
- Lot of effort put on securing system
 - system evaluation
 - security documentation
 - development methodology
 - \Rightarrow Expensive trill
 - cost of external evaluators
 - cost of own work
- Evaluated system *less likely* to have problems

Evaluation history

- Originated from military and government
- Need for assurance in commerce
 - trustworthy contractors
 - information systems
- Legal requirements, liability
 - cover your back: it was certified

Trusted computer system evaluation criteria (TCSEC)

- Orange Book (1983–1999)
- U.S. Government
- Evaluation classes
 - C1, C2, B1, B2, B3, A1 (D for failed)
 - mainly information confidentiality
 - no requirements for availability
- Functional requirements
 - discretionary access control (DAC)
 - object reuse
 - mandatory access control (MAC) (B1)
 - label requirements (for MAC) (B1)
 - identification and authentication
 - trusted path (B2)
 - audit mechanism
 - architecture
- Operational requirements
 - separation of roles (B2)
 - secure recovery (A1)
 - system integrity validation (A1)
- Assurance requirements
 - configuration management (B2)
 - trusted distribution (A1)
 - system architecture (C1)
 - design specification and verification
 - C1,C2 no requirements
 - B1 informal
 - B2 descriptive top level specification (DTLS, formal)
 - B3 DTLS consistent with security policy
 - ${\bf A1}\,$ formal top level specification (FTLS) with formal methods and mapping to source code
 - testing
 - product documentation
 - internal documentation

TCSEC evaluation classes

- **D: failed** systems that have failed evaluation
- C1: discretionary protection identification and authentication
- C2: controlled access protection used for commercial products
- B1: labelled security protection MAC for set of objects
- B2: structured protection MAC, trusted path, least privilege, covert channel analysis
- B3: security domains reference validation, requirements for development methodology
- A1: verified protection formal methods to evaluate B3 requirements

TCSEC evaluation process

- 1. Application
- 2. Preliminary technical review (PTR)
 - readiness review
- 3. Evaluation
 - (a) design analysis
 - based on documentation
 ⇒ requirements for complete and correct
 - (b) test analysis
 - coverage assessment
 - vendor-supplied tests
 - (c) final review
- Government-sponsored evaluators
- Ratings maintenance program

How good TCSEC is?

- Set baseline for evaluations
- Limited scope
 - only for operating systems
 - U.S. military and government needs (no integrity, availability or other business needs)
- Problems with process
 - additional criteria
 - slow process

FIPS 140 (1994–)

- Evaluation of cryptographic modules
- Based on levels
 - 1. FIPS-approved algorithm; software or hardware
 - 2. physical security, role-based authentication, EAL2
 - 3. enhanced physical security, EAL3
 - 4. detecting and responding to physical access

- Areas of requirements (11 total)
 - cryptographic module specification, parts, interfaces
 - roles, authentication
 - logical model, design
 - physical security (EMI/EMC)
 - operating environment
 - mitigating attacks

level	4	3	2	1
FIPS 140-1	9	66	143	86
FIPS 140-2	-	40	89	88

Common Criteria (1998–)

- International followup to TCSEC, etc.
 - Common Criteria Recognition Agreement (CCRA)
 - ISO 15408
- Three parts
 - CC documents
 - CC evaluation methodology
 - national scheme
- Two types of evaluation
 - **PP: protection profile** implementation independent set of requirements for set of products or systems

ST: security targets evaluation of single product or system

- Requirements
 - security functional
 - * 11 classes with 2-16 families
 - assurance
 - $\ast~10$ classes with 2–8 families

CC Levels build on assurance

- 1. Functionally tested
- 2. Structurally tested
- 3. Methodically tested and checked
- 4. Methodically designed, tested and reviewed
- 5. Semi-formally designed and tested
- 6. Semi-formally verified design and tested
- 7. Formally verified design and tested

Approximate correspondence of TCSEC and CC; number of certified equipment for EALn. TCSEC CC count other

D	-		
-	EAL1	11	
C1	EAL2	62	need for FIPS 140-2 L2 $$
C2	EAL3	43	need for FIPS 140-2 L3 $$
B1	EAL4	63	need for FIPS 140-2 L4 $$
B2	EAL5	1	
B3	EAL6		
A1	EAL7		

ISO 17799

- "A comprehensive set of controls comprising best practises in information security"
 - code of practise
 - specification for information security management system (BS7799-2)
- Generic information security standard
- Control IT system risks
 - product auditing, ...
 - policy must cover all areas
- Plan-Do Check-Act

Threat modelling

- Target: understand and document security threats
- Large number of possible threats
 ⇒ Ad-hoc treat searching incomplete
 ⇒ Must be methodological
- System threat profile described
- Characterisation of system security
- Threat is *not* vulnerability
 - vulnerability is unmitigated threat
 - attack classification important

Collecting information

- How system will be used
- What system depends on (environment)
- What assumptions are made on implementation
- How system interacts to environment
- Basics of internal design decissions

Modelling system

- What are entry points
 - network, services, user interface, files, disk system
- Assets to protect
- Trust levels
 - user groups, unidentified
 - access to assets
- Data flow
 - how data flows from entry to processes

Determining threats

- Threat identification
 - shadow corners of valid-but-malicious data
 - invalid data
- Analyse threat
 - does it result a vulnerability
 - how it can be mitigated

Threat effect classification: STRIDE

Spoofing allows obtaining false identity
Tampering modifies data for goal
Repudiation not providing evidence to point guilty
Information disclosure for unauthorised user
Denial of service for legitimate users
Elevation of privilege for higher trust level

Risk of vulnerability: DREAD

Damage potential if exploited
Reproducibility of vulnerability
Exploitability how easily vulnerability can be exploited
Affected users if exploit is widely available
Discoverability is likehood that vulnerability will be found
Assign weight for each category, calculate average.

Threat tree [1]

- Goal as tree root
- $\bullet\,$ An attack is decomposed to sub-goals

AND all sub-goals must be meet **OR** any of subgoals is sufficient

- Attack costs or pre-requirements can be assigned
 - helps to determine seriousness
- Reuse of attack patterns

Survivability Compromise: Disclosure of ACME proprietary secrets



- tacker over Internet
 - 3. Gain privileged access to Web server
 - 6. Attack ACME intranet using its connections with public telephone network (PTN)
 - OR 1. Monitor communications over PTN for leakage of sensitive information

2. Gain privileged access to machines on intranet connected via Internet

Survivability Compromise: Disclosure of ACME proprietary secrets



ACME web server



How to blackhole traffic



How to establish unauthorised BGP session



How to originate unauthorised prefix attribute into peer route table



Penetration tests

- Experimental evaluation of system security
- Layered models: different threats
 - 1. external attacker without knowledge about system
 - 2. external attacker with access to system
 - 3. internal attacker with access to system
- Flaw hypothesis methodology
 - 1. information gathering
 - 2. flaw hypothesis
 - 3. flaw testing
 - 4. flaw generalisation
 - 5. flaw elimination
- Unsuccessful penetration does not prove system secure
 - cracking contests mostly useless publicity stunts

Malicious logic

Trojan horse user unintentionally executes program

- documented effect (what user expects)
- covert effect (malicious)
- Trojan in compiler [2]

- "free" software add-ons (spyware, adware)
- may replicate itself

Virus inserts itself to file

- may have malicious actions
 - corrupts files
 - destroys equipment
- loss of performance
- several subtypes by infection, implementation method

Worm propagates between systems

- may have an impact on network
- most current malicious logic
 - massmailers
 - chat
 - p2p networks

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Rabbits/bacteria exhaust resources quickly
main(){for(;;)fork();} (DO NOT run code on public systems...)
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Logic bombs event triggers malicious action

• disgruntled employee

Protection against malicious code

- Too coarse access control
 - any program has access to all my data
 - root/administrator omnipotent
- Enforcing principle of least privilege
 - sandboxing
 - capacity model
- Code signing not a solution

Summary

- Formal evaluation supports system development
 - for higher levels must be integral
 - overall quality assurance
- Threat evaluation needed in development
- Threat trees help large system evaluation
 - even if components are certified, network maybe insecure
- Penetration testing practical evaluation
 - needs high level of skill and experience

References

- [1] Bruce Schneier. Secrets and Lies: digital security in a networked world. Wiley Computer Publishing, 2000.
- [2] Ken Thompson. Reflections on trusting trust. Commun. ACM, 27(8):761-763, 1984.