

# P2P VoD Systems: Modelling and Performance

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- Peer-to-peer systems
- File sharing: fluid model
- File sharing: steady-state analysis
- File sharing: conclusions
- Video-on-demand: fluid model
- Video-on-demand: steady-state analysis
- Video-on-demand: steady-state synthesis
- Video-on-demand: conclusions

#### **Fundamental principle**

#### Client/Server (CS) paradigm

- Clients download content from servers
- Clear distinction between the two roles
- Service capacity remains the same, while load increases
- When too many clients, transfer times explode
- Offered load bounded by this stability limit (for sure!)
- Peer-to-peer (P2P) systems
  - Peers download pieces of content from other peers/seeds and simultaneously upload downloaded pieces to other peers
  - Blurring of roles: peers not only act as clients (when downloading) but also serve other peers (when uploading)
  - Service capacity scales with the offered load
  - No stability limit (for sure?)

### **Applications**

- P2P used commonly for file sharing (e.g. BitTorrent) and live streaming
- P2P video-ondemand (VoD):
  - Alternative to client-server approaches (YouTube)?
  - Under what conditions?



#### **Quality of Service**

#### • P2P file sharing

- Retrieve the whole file as soon as possible
- Retrieve pieces in any order
- Minimize the file transfer time

#### • P2P streaming

- Retrieve pieces at least at playback rate and in almost sequential order
- Minimize the startup delay (needed to fill the playout buffer)
- P2P video-on-demand
  - Retrieve the whole file
  - Retrieve pieces at least at playback rate and in almost sequential order
  - Minimize the startup delay (needed to fill the playout buffer)

### Why performance modelling?

- Scalability
  - Is the system really scalable?
- Stability
  - If not, where is the stability limit for the load?

#### • Performance

- When stable, is the performance sufficient?

#### **Modelling aspects**

- Dynamic population model
  - describing the evolution of the peer population in the P2P system
- Peer arrival process
  - steady arrival rate, smoothly attenuating arrival rate, or flash crowd?
- Efficiency of resource sharing
  - utilization of a peer's upload capacity
  - effect of the piece/peer selection policy
  - number of parallel connections
- Selfishness / altruism
  - part of peers are free-riders that do not want to share upload capacity
- Download and upload rates
  - homogeneous or heterogeneous peer population?
- Number of permanent seeds
  - correspond to servers in the client-server architecture

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### Model for P2P file sharing

- Life span of a peer consists of two sequential phases:
  - file transfer phase, during which the peers are called leechers
  - sharing phase, during which the peers are called seeds
- Altruistic peers have a longer sharing phase than selfish peers
- Model by Qiu and Srikant (2004):
  - deterministic fluid model (= system of differential equations)
  - describing the system dynamics related to sharing of a single file
  - x(t) = (average) number of leechers at time t
  - y(t) = (average) number of non-permanent seeds at time t



#### Assumptions

- Steady arrival process described by
  - arrival rate  $\lambda$  to transfer phase (arrivals per time unit)
- Efficiency described by
  - upload utilization ratio  $\eta$  (belonging to (0,1])
- Selfishness described by
  - departure rate γ from service phase (departures per time unit)
- Homogeneous peer population with
  - download rate c (file transfers per time unit) and
  - upload rate  $\mu$  (file transfers per time unit)
- No permanent seeds



#### Fluid model

• Switched nonlinear system:

$$\begin{cases} x'(t) = \lambda - \phi(t), \\ y'(t) = \phi(t) - \gamma y(t), \end{cases}$$
(1)

• Aggregate service rate:

$$\phi(t) = \min\{cx(t), \mu(\eta x(t) + y(t))\}.$$
(2)



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#### **Steady-state analysis**

• Solve the equilibrium of the system by setting x'(t) = y'(t) = 0 in (1):

$$ar{\phi} = \lambda, \ ar{y} = rac{\lambda}{\gamma},$$

- Two cases considered separately:
  - download-constrained system in equilibrium
  - upload-constrained system in equilibrium
- Parameter space divided nicely in two complementary parts each of which has a unique equilibrium solution
  - that are even globally stable by Qiu and Sang (2008)

(3)

### **Download-constrained system**

If

$$c\bar{x} \le \mu(\eta\bar{x} + \bar{y}),\tag{4}$$

then the system is *download-constrained*, and

$$c\bar{x} = \bar{\phi} = \lambda,\tag{5}$$

implying that

$$\begin{cases} \bar{x}_d = \frac{\lambda}{c}, \\ \bar{y}_d = \frac{\lambda}{\gamma}. \end{cases}$$
(6)

The resulting condition for a download-constrained system:

$$\frac{1}{\mu} \le \frac{\eta}{c} + \frac{1}{\gamma}.\tag{7}$$

### **Upload-constrained system**

If

$$c\bar{x} > \mu(\eta\bar{x} + \bar{y}),\tag{8}$$

then the system is *upload-constrained*, and

$$\mu(\eta \bar{x} + \bar{y}) = \bar{\phi} = \lambda, \tag{9}$$

implying that

$$\begin{cases} \bar{x}_u = \frac{\lambda}{\eta} \left( \frac{1}{\mu} - \frac{1}{\gamma} \right), \\ \bar{y}_u = \frac{\lambda}{\gamma}. \end{cases}$$
(10)

The resulting condition for a upload-constrained system:

$$\frac{1}{\mu} > \frac{\eta}{c} + \frac{1}{\gamma}.\tag{11}$$

#### **Deterministic model vs. stochastic simulations**



Figure 1: Experiment 1 : The evolution of the number of seeds as a function of time



Figure 2: Experiment 1 : The evolution of the number of downloaders as a function of time

Source: Qiu and Srikant (2004)

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- Video-on-demand: conclusions

### **Conclusions from the P2P file sharing model**

- Scalability
  - System scalable in the whole parameter space by (6) and (10), in particular for any  $\eta > 0$
- Stability
  - Consequently, system stable for any  $\lambda > 0$
- Performance
  - By Little's formula, the mean file transfer time is

$$T = \frac{\bar{x}}{\lambda} = \max\{\frac{1}{c}, \frac{1}{\eta}\left(\frac{1}{\mu} - \frac{1}{\gamma}\right)\} \le \max\{\frac{1}{c}, \frac{1}{\eta\mu}\} \approx \max\{\frac{1}{c}, \frac{1}{\mu}\}.$$
 (12)

- Thus, no real problems in performance if reasonable download and upload rates with respect to the mean file size
- The last approximation justified for the file sharing application (mainly due to the free retrieving order of pieces)

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#### Model for P2P VoD

- Life span of a peer consists of two overlapping phases:
  - file transfer phase, during which the peers are called leechers
  - watching phase, starting immediately after the initial buffering delay
- Altruistic peers become seeds after the file transfer phase if the watching phase still continues
- Model by Aalto et al. (2010):
  - deterministic fluid model (= system of differential equations)
  - describing the system dynamics related to sharing of a single video file
  - x(t) = (average) number of leechers at time t
  - y(t) = (average) number of non-permanent seeds at time t



#### **Assumptions (1)**

- Steady arrival process described by
  - arrival rate  $\lambda$  (arrivals per time unit)
- Efficiency described by
  - upload utilization ratio  $\eta$  (belonging to (0,1])
- Altruism described by
  - probability  $\zeta$  (for a peer to become a seed)
- Homogeneous peer population with
  - download rate c (file transfers per time unit) and
  - upload rate  $\mu$  (file transfers per time unit)
- Number of permanent seeds = k (belonging to {0,1,2,...})



### **Assumptions (2)**

- Startup delay negligible (if video sufficiently long)
  - Thus, the transfer phase and the playback phase start essentially at the same time
- Video watched at (fixed) playback rate
  - Total watching time denoted by z
  - Natural requirement: z > 1/c (since transfer rate always bounded by c)
- Playback quality problems if the transfer phase takes longer than z
  - In this case, the playback phase ends as soon as the transfer is completed
- Selfish peers stay in the system until the end of the transfer phase while altruist peers stay until the end of the playback phase
  - but no longer, which is a worst case scenario



#### Fluid model

• Switched nonlinear system:

$$\begin{cases} x'(t) = \lambda - \phi(t), \\ y'(t) = \zeta \phi(t) - \frac{y(t)}{z - x(t)/\lambda}, \end{cases}$$
(13)

• Aggregate service rate:

$$\phi(t) = \min\{cx(t), \mu(\eta x(t) + y(t) + k)\}.$$
(14)



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#### **Steady-state analysis**

• Solve the equilibrium of the system by setting x'(t) = y'(t) = 0 in (13):

$$\begin{cases} \bar{\phi} = \lambda, \\ \bar{y} = \zeta \lambda (z - \bar{x}/\lambda), \end{cases}$$
(15)

- Two cases considered separately:
  - download-constrained system in equilibrium
  - upload-constrained system in equilibrium
- Multiple solutions found
- Local stability analysis used to rule out some of them

### **Download-constrained system**

If

$$c\bar{x} \le \mu(\eta\bar{x} + \bar{y} + k),\tag{16}$$

then the system is *download-constrained*, and

$$c\bar{x} = \bar{\phi} = \lambda,\tag{17}$$

implying that

$$\begin{cases} \bar{x}_d = \frac{\lambda}{c}, \\ \bar{y}_d = \zeta \lambda \left( z - \frac{1}{c} \right). \end{cases}$$
(18)

The resulting condition for a download-constrained system:

$$\frac{1}{\mu} \le \frac{\eta}{c} + \zeta \lambda \left( z - \frac{1}{c} \right) + \frac{k}{\lambda}.$$
(19)

### **Upload-constrained system (1)**

If

$$c\bar{x} > \mu(\eta\bar{x} + \bar{y} + k), \tag{20}$$

then the system is *upload-constrained*, and

$$\bar{\phi} = \mu(\eta \bar{x} + \bar{y} + k), \tag{21}$$

implying that

$$\begin{cases} \bar{x}_u = \frac{\lambda}{\eta - \zeta} \left( \frac{1}{\mu} - \zeta z - \frac{k}{\lambda} \right), \\ \bar{y}_u = \frac{\zeta \lambda}{\eta - \zeta} \left( -\frac{1}{\mu} + \eta z + \frac{k}{\lambda} \right). \end{cases}$$
(22)

The resulting conditions for a upload-constrained system:  $\eta \neq \zeta$  and

$$\begin{cases} \frac{1}{\mu} > \frac{\eta}{c} + \zeta \left( z - \frac{1}{c} \right) + \frac{k}{\lambda}, & \text{if } \eta > \zeta, \\ \frac{1}{\mu} < \frac{\eta}{c} + \zeta \left( z - \frac{1}{c} \right) + \frac{k}{\lambda}, & \text{if } \eta < \zeta. \end{cases}$$
(23)

#### **Upload-constrained system (2)**

Additionally, for the solution to be meaningful, we require that  $\bar{x}_{u} > 0$  and  $\bar{y}_{u} > 0$ . The former one follows from (23), but the latter one leads to the following additional constraints:

$$\begin{cases} \zeta < \frac{1}{z} \left( \frac{1}{\mu} - \frac{k}{\lambda} \right) < \eta, & \text{if } \eta > \zeta, \\ \eta < \frac{1}{z} \left( \frac{1}{\mu} - \frac{k}{\lambda} \right) < \zeta, & \text{if } \eta < \zeta, \end{cases}$$
(24)

which implies that  $0 < \frac{1}{\mu} - \frac{k}{\lambda} < z$  is a necessary condition for the existence of a non-negative upload constrained solution.



Fig. 1. Left panel: Download-constrained solution with +/+ expressing that  $\bar{x}_d > 0$  and  $\bar{y}_d > 0$  in this area. Right panel: Upload-constrained solution with +/+ [+/-] expressing that  $\bar{x}_u > 0$  and  $\bar{y}_u > 0$  [ $\bar{y}_u < 0$ ] in this area.

#### Summary of the steady-state analysis (2)



Fig. 2. Solution areas, where DL [UL] refers to a positive download [upload] constrained solution. The horizontal bordering line satisfies  $\frac{1}{\mu} = \frac{\eta}{c} + \zeta(z - \frac{1}{c}) + \frac{k}{\lambda}$  and the vertical bordering line satisfies  $\eta = \frac{1}{z}(\frac{1}{\mu} - \frac{k}{\lambda})$ .

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### **Steady-state synthesis (1)**



Fig. 1. Left panel: Solution areas, where DL [UL] refers to a positive download [upload] constrained solution. The horizontal bordering line satisfies  $\frac{1}{\mu} = \frac{\eta}{c} + \zeta(z - \frac{1}{c}) + \frac{k}{\lambda}$  and the vertical bordering line satisfies  $\eta = \frac{1}{z}(\frac{1}{\mu} - \frac{k}{\lambda})$ . Right panel: Steady state synthesis.

### **Steady-state synthesis (2)**

## • If $\eta < \frac{1}{z} \left( \frac{1}{\mu} - \frac{k}{\lambda} \right)$ • If $\eta > \frac{1}{z} \left( \frac{1}{\mu} - \frac{k}{\lambda} \right)$

transfer rate < playback rate, i.e. playback quality problems

• Number of leechers and seeds well estimated by  $(x_0, y_0)$ :

$$\begin{cases} \bar{x}_0 = x_u|_{\zeta=0} = \frac{\lambda}{\eta}(\frac{1}{\mu} - \frac{k}{\lambda}), \\ \bar{y}_0 = y_u|_{\zeta=0} = 0. \end{cases}$$

transfer rate > playback rate, i.e. sufficient playback quality

• If further

$$\frac{1}{\mu} \leq \frac{\eta}{c} + \zeta \lambda \left( z - \frac{1}{c} \right) + \frac{k}{\lambda}.$$

DL constrained system  $(x_d, y_d)$ 

• Otherwise UL constrained system  $(x_u, y_u)$ 

#### **Deterministic model vs. stochastic and BitTorrent simulations**



Fig. 4. Comparison of the fluid model (solid smooth line) against the stochastic model (dashed line) and the BitTorrent simulation (solid jagged line) with  $\zeta = 0.9$  (upper panel) and  $\zeta = 0.3$  (lower panel).

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### **Conclusions from the P2P VoD model**

- Scalability
  - System scalable in the whole parameter space by the synthesis, in particular for any  $\eta > 0$
- Stability
  - Consequently, system stable for any  $\lambda > 0$
- Performance
  - Playback quality problems if the efficiency parameter  $\eta$  is too small
  - On the other hand, performance even "scales" (= good quality for all  $\lambda$ ) if the efficiency parameter  $\eta$  is sufficiently large
  - Transfer rates for DL and UL constrained cases:

$$R_d = c$$
$$1/z < R_u < c$$

#### References

- D. Qiu and R. Srikant, Modeling and performance analysis of BitTorrent like peer-to-peer networks, in ACM SIGCOMM, pp. 367-378, 2004.
- [2] D. Qiu and W. Sang, Global stability of peer-to-peer file sharing systems, Computer Communications, 31, 2, 212-219, 2008.
- [3] S. Aalto, P. Lassila, N. Raatikainen, P. Savolainen, and S. Tarkoma, P2P Video-on-Demand: Steady state and scalability, to appear in *IEEE Globecom*, 2010.