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Multicast as an ISP service


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Goals of this lecture


- After this lecture you will be able to
 - Give an overall technical view of multicast
 - Addressing, protocols, routing
 - Please note, that Multicast (as a technology) is also lectured at least in 38.2122 and 38.3188
 - List different application and service types that could be deployed with multicast
 - Discuss the problems in bringing Multicast to production use in ISP network
 - Discuss the different aspects of integrating QoS solutions to Multicast




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Packet distribution I


- TCP/IP stack implements four types of packet distribution methods:
 - **Unicast:** From one to another host
 - Normal point-to-point IP traffic
 - The packet is sent only to the receiver
 - **Broadcast:** From one host to all hosts on a network
 - “Is there anybody out there?”
 - The packet is seen by all hosts on the (local) network



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Packet distribution II

- **Anycast:** From one host to one host in a group of hosts
 - To access a service served by multiple hosts (in IPv6)
 - The packet is received by one of the receivers
- **Multicast:** From one to a group of hosts
 - All group members receive the packet
 - Packet duplicated only when necessary



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Many-to-many traffic

- How to send “many-to-many” traffic?
 1. Central servers: Servers duplicate packets to all members of the group (Unicast).
 2. Connected mesh: Every group member has a connection to all or some others (Unicast/Broadcast).
 3. Multicast: Network duplicates the packets as needed.
- First two choices do not scale and waste resources (bandwidth, especially).

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Multicast addressing in IPv4

- Multicast traffic uses class D IP-addresses:
 - Flat hierarchy
 - 224.0.0.0 - 239.255.255.255
 - 224.0.0.0 – 224.0.0.255 reserved for multicast protocols
 - 224.0.1.0 – 224.255.255.255 reserved for various applications
 - 225.0.0.0-238.255.255.255 for free use
 - 239.255.0.0 – 239.255.255.255 for private use (administratively scoped)
 - Not just TTL-bounded, limited vendor support
- See <http://www.iana.org/assignments/multicast-addresses> for complete list.

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Multicast addressing in IPv6

- Address structure

8	flgs(4)	scope(4)	group id (112)
11111111	000t	xyzn	

↑
Transient bit
- The scope field determines the maximum “depth” of the multicast group
 - 1=Node local, 2=Link local, 5=Site local, 8=Organization local, E₁₆=Global scope

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Multicast routing

- Routers form multicast trees
 - The tree leaves are the receivers
 - The trees are maintained with multicast routing protocols and often referred to as spanning trees
 - Trees come in several formats
 - Shortest Path from each source
 - Shortest path from a RP (core of the tree)
 - Unidirectional shared tree
 - Bi-directional shared trees
- Multicast routing protocols include
 - DVMRP, MOSPF, PIM-SM, PIM-DM, CBT, OCBT, HIP and BGMP, to name a few.

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Multicast packet forwarding

- RPF-check in the router:
 - All multicast packets are checked for the source address (not for the destination as in unicast forwarding)
 - The incoming interface must be an interface pointing towards the source address (sender)
 - If this checks, packet is forwarded to appropriate interfaces, otherwise the packet is dropped.

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Multicast group management

- A host must *join* a specific group to receive the traffic in that group but can send to a group without joining.
 - Non-member senders!
- Membership is controlled by the *IGMP* protocol.
 - IGMPv2 is the most commonly used (and an IETF standard)

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IGMPv2

- Backwards compatible with IGMPv1
- 3 message types
 - Membership Query (general or group specific)
 - Membership report (general (v1) or addressed to the group (v2))
 - Leave Group
- IGMPv3 adds support (to Membership Query msgs) for a system to report interest in receiving packets **only** from specific source addresses, or from **all but** specific source addresses, sent to a particular multicast address.

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Hardware for ISP

- De facto architecture
 - IGMPv2, DVMRP, MOSPF, PIM-SM (within autonomous systems) and MSDP or MBGP (for interdomain multicast routing)
 - These protocols are readily usable with most commercial routers
 - However, most of the time these are turned off!
 - Why?

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Why Multicast?


- **Multicast saves bandwidth**
 - Multicast is desirable for sources(servers) and administrators
 - Sources may scale service to large number of users
 - Redundant connections could be joined
 - Receivers do not care whether content is received via unicast or multicast
 - BW/QoS/realtime requirement is the same anyway



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Application drivers


- Unicast killer apps: Email+WWW
- Multicast apps (one-to-many or few-to-few)
 - Audio & Video distribution
 - Push applications (PointCast)
 - Audio & Video conferencing
 - File transfer (updates, file sharing, web caching, distributed databases)



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MBone multicast service


- Virtual network (using the Internet unicast resources) transmitting multicast-packets
 - Tunneling
- Multicast is a method of routing and delivering packets in a many-to-many situation
- MBone is one implementation of multicast method.




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Multicast in MPLS

- Manually with explicit trees
 - Requires extra work, but is an option in small networks
- MPLS networks utilize Multicast Distribution Trees (MDT)
 - Point-to-multipoint and multipoint-to-multipoint LSPs
- MPLS/BGP/VPNs
 - In unicast VPNs the P routers do not maintain per-VPN state (only per-PE)
 - In multicast VPNs the P routers have to maintain some information on active multicast trees (to avoid forwarding unneeded traffic).
- MPLS+Multicast is still very much work in progress







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IP multicast as a service

- Open service –model
 - No restrictions on creating multicast groups, receiving from a group or sending to a group
 - Group membership means that one is reachable to other members of the group
 - No access control
 - Multiple senders may share the same multicast address
 - No pre-emption or address reservation
 - Group size is dynamic and unknown, member status is unknown
- Summary: IP multicast group is not managed (or the management is distributed at best)





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Customer requirements

- ISP is a customer-driven business
- Any service offered must add value to the customer
- For multicast to add value, customers must have
 - Global access to the service
 - Ease and transparency of install and use
 - Controlled group membership
 - Unique addressing (for the session)
 - Reliable transmission





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Management

- Lack of NAT support (work in progress)
- Firewalls do not recognize multicast addresses
 - Solution: tunnels. Consequence: A hole in the wall. Another solution: static routes (and more human management and therefore more expensive service)




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Domain independence

- Traffic sources in other domains are out of the ISP control
 - Traffic control
 - Service level
 - Transit traffic is not preferred
 - Control of address advertisements



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Multicast benefits for the ISP

- Multicast reduces BW required to transport data to multiple recipients
 - Multicast is expensive (aka H-A-R-D) to deployment, installation and management
 - Are BW savings higher than cost of deployment and management?

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Requirements for IP multicast service

- All routers must support multicast (Customers want global access)
- Domain independence
- Additional requirements
 - Group management
 - Authorization for group creation, receiver and sender authorization
 - Distributed multicast address allocation
 - Security
 - Protection against attacks on multicast routes and sessions
 - Support for data integrity mechanisms.
 - Support for network management
 - Billing Multicast

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Multicast and QoS I

- RSVP + IntServ
 - IntServ built with multicast in mind
 - Receiver requests QoS based on Sender description of the service
 - Unidirectional reservations
 - Filters (Shared Explicit, Wildcard, Fixed Filter) illustrate different types of multicast sessions
 - Reservation merging
 - Usual IntServ problems persist
 - Scalability issues
 - Local management (CAC, policy control)

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Multicast and QoS II

- DiffServ
 - There should be no per-group state in the core -> reservation uncertainty
 - Multicast group size not defined in advance and DiffServ not having per-user access control -> unknown resource usage
 - Replication of packets in the network may cause problems in capacity handling
 - Problem of relative QoS
 - PDB in one domain is not necessary the same PDB in other domains -> need for dynamic DSCP allocation -> not a standardized approach



Alternatives for Multicast

- **Current Multicast architecture needs more functionality**
 - Address allocation, Access control, Interdomain management
- **Alternative options**
 - Single-sender service (IGMPv3)
 - Multipeer service
 - Core discovery or advertisement
 - Well-located
 - Replicated
 - Application level multicast
 - Receivers duplicate data and forward it onwards

