

# Firewalls and intrusion detection systems

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2005-04-11

## Lecture topics

- Firewalls
- Security model with firewalls
- Intrusion detection systems
- Intrusion prevention systems
- How to prevent and detect attacks

## What is a firewall

- Divides network into two (or more) parts with *different security policy*
  - internal network  $\Leftrightarrow$  Internet
  - engineering  $\Leftrightarrow$  accounting: the other network need not be a less secure one than the other one. They just have different security policies or different assets to protect.
  - internal network  $\Leftrightarrow$  public servers  $\Leftrightarrow$  Internet
  - building automation  $\Leftrightarrow$  VoIP  $\Leftrightarrow$  surveillance system
- Enforces security policy
  - allowed traffic
  - prohibited traffic

Refer to IPSec security policy database (SPD): traffic is bypassed, discarded, or bypassed as protected.

- May have additional roles, such as a VPN endpoint

## Firewall types

**Packet-filtering** makes decision based only packet fields

- router ACL (access control list)
- TCP implicit state: for example to disallow incoming connections, firewall will drop any packet that has SYN flag set but no ACK and allows any packet with SYN+ACK.
- difficult with UDP, also some other TCP-based protocols such as FTP in active mode, where server establishes connection to client.

**Stateful** keeps track on connections

- maintains connection state
  - single point of failure
  - has to have some timeout mechanism as the state space is limited. Some attacks may exhaust state space.
    - $\Rightarrow$  random disconnections
- possible to accept related connections: some protocols need an application gateway.

**Application gateway** interpret connection on application level

- checks if application traffic is valid
- protects from a simple port changes like running ssh protocol on port 443 (https).
- may provide a payload inspection to detect malicious payload
- proxy servers
  - call-out
  - in-line (transparent)

**Address-translation** between internal numbering and external addresses

- using NAT provides same security as prohibiting incoming TCP and UDP
- internal topology can be hidden

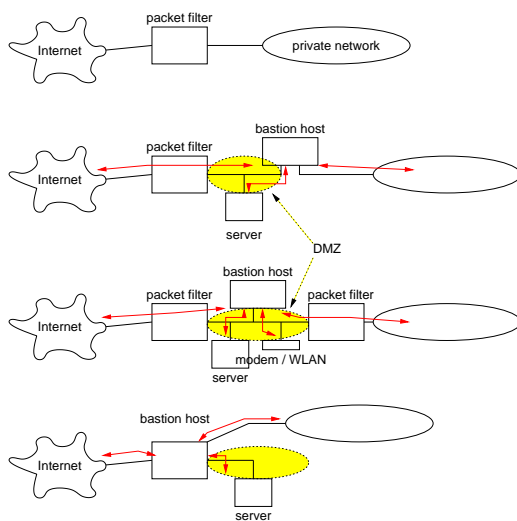
**Host-based** or software firewalls add on application security

- completes application security and access control
- possibly user- and application-level control

**Hybrid** use combination of different types for performance

- check start of connection with application gateway, switch to stateful filtering  
⇒ better performance as the bulk of traffic is handled by the fast path.

## Firewall topologies



## Building firewall rules

- Defining default policy
  - “everything not prohibited is allowed”
    - \* “router” ACL
    - \* enumerate vulnerable services and protect them
  - “everything not allowed is prohibited”
    - \* enumerate need and safe services and allow only those
  - both policies need continuous updating
- There should be only one rule matching for each packet
  - multiple overlapping rules
  - order of rules matters

- performance issues: hardware-based routers/firewalls can handle certain number of rules without significant performance penalty. For software-based firewalls order of rules does matter.
- Possibility to oversight
- High-level specification languages are not a solution

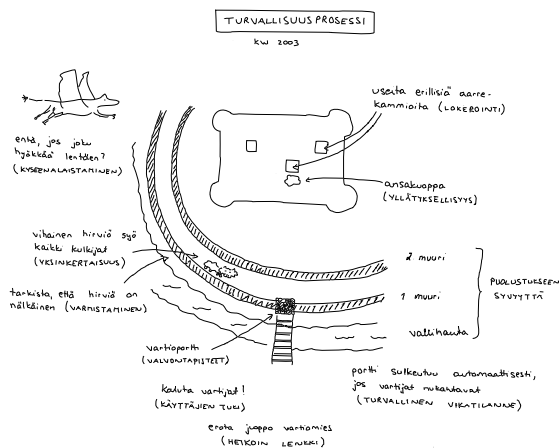
## Deploying multiple firewalls

- Helps to limit the impact of attack
- Protection by diversity
  - on the other hand, multiple systems to update
- Designing rules even more complicated

## What firewall protects and what not

- Protects
  - from known, vulnerable protocols
  - static network configuration
- Does *not* protect for / from
  - executable/active content
  - malicious insider
  - loopholes: modems, WLAN, mobile networks
  - carry-in/out attacks such as notebooks, mass storage, rogue WLAN APs
  - new attacks using applications previously considered safe
  - most DoS attacks
- May result a “hard perimeter, mellow inside”
  - failure to update internal systems
  - selecting insecure protocols and applications

## Security in organisation



## How secure are firewalls

- Common Vulnerabilities and Exposures: 190 matches on “firewall” <http://cve.mitre.org/cve/> (numbers below are from 2005 when there was total 110 entries)
  - Check Point FireWall-1** 34 entries
  - Cisco** 13 entries
  - Juniper** 1 entry
  - Linux** 6
  - Symantec** 17
  - WatchGuard** 11 entries
- More features (VPN, virus checks, QoS protection)
  - ⇒ more code
  - ⇒ more bugs
  - ⇒ more vulnerabilities

## Intrusion Detection Systems

- How to make sure that the firewall is not leaking
  - rule-based
  - anomaly-based
- How to detect internal attacks
- IDS is designed to
  - detect,
  - identify, and
  - report malicious activity
- IDS can be located different places
  - application
  - host
  - network

## Application and host IDS

- An application instrumented to identify abnormal actions
  - high level of abstraction
  - user actions monitored
  - policy violations
  - application log analysis
  - access to encrypted data
  - may not protect from application flaws
- Host instrumented
  - reference monitor
  - actions by a user and an application
  - host log analysis
- Log analysis best done on separate host
  - provides after-the-fact analysis
  - vulnerable to network attacks DoS on log server
  - messages transmitted in clear unless IPSec is deployed

## Network IDS

- Monitors traffic
  - best done with signal splitters operating on physical layer
- Large volume of data
  - low level of abstraction
  - encrypted traffic problematic
- Mostly misuse detection
  - recorded patterns of misuse (signatures)
  - frequent updates (like virus scanners)

```
alert tcp $EXTERNAL_NET any -> $HOME_NET 22
( msg:"EXPLOIT ssh CRC32 overflow /bin/sh";
  flow:to_server,established;
  content: "/bin/sh"; )
```

- Anomaly detection
  - detecting differences to normal
    - \* threshold detection
    - \* statistical profile
    - \* rule-based detection
  - learning system
- Large number of alerts: an example
  - 3700 alerts from corporate network per day
  - 48 should be studied in detail
  - 2 warrant an action

## IDS in large network

- One should monitor every link
  - ⇒ very expensive
- Select important links
  - full census on those
- Do random sampling on other links
  - if one samples every 512th packet and sends it to a central location
    - ⇒ not a big increase in traffic
  - large problems notified immediately

## Honeypots

- A false system similar to production system
  - all access illegal
    - ⇒ any accessing is a potential intruder
- Used as part of IDS
  - a connection results detailed monitoring
- How to keep an attacker from telling the difference from a real system
  - should be not too weak
  - should have “real” data and traffic
  - if a virtual host, should not be visible

## IDS reaction too slow

- IDS identifies attack
  - analysis may not be real-time
  - corrective actions may take time
- Epidemic security problem may require instant actions [5]
- A system can be scanned, attacked, and compromised in a minute or less  
⇒ Need for an automatic security system

## Intrusion Prevention Systems (IPS)

- IDS with an automatic response
- Suffers from a large number of false alerts  
⇒ may result denial of service
- A firewall with automatic ACL update
- Virus scanners are host-based IPS
- Still at early stages
  - does not stop vendors from marketing...

## Traffic traceback

- Problem: where incoming attack traffic originates
- Source IP cannot be trusted
  - sender can put it to any address
  - ingress filtering not deployed universally [1]
- Should not need additional hardware or load on routers
- Scalability problems, few proposals [2, 3, 4]

## Security in Ad-hoc networks

- Ad-hoc networks an interesting topic
  - self-building topology
  - extending network coverage
- Must rely on the other hosts
  - no central authority, block lists
  - no trusted core network
  - routing done by devices
- Public key-based per-packet authentication too heavy
  - modern PC throughput few ten kbit/s, much less for battery-powered device
- How to communicate trustfulness?

## Challenges in All-IP world

- Large number of non-technical users
  - the “--:--” generation
  - rightful ignorance: I want to watch movies — fixing security problems does not match to my idea of relaxing.
- Service provider responsibility
- Multi-vendor environment

## Summary

- Firewall and IDS are good tools
- Must know their limitations
- Future challenges
  - accurate detection of malicious activity
  - security in ubiquitous computing
  - trust in autonomous systems
  - providing security for couch potatoes

## References

- [1] J. Case, R. Mundy, D. Partain, and B. Stewart. Introduction to Version 3 of the Internet-standard Network Management Framework. Request for Comments RFC 2570, Internet Engineering Task Force, April 1999. (Informational) (Obsoleted by RFC3410). URL:<http://www.ietf.org/rfc/rfc2570.txt>.
- [2] Stefan Savage, David Wetherall, Anna Karlin, and Tom Anderson. Practical network support for IP traceback. In *Proceedings of the 2000 ACM SIGCOMM Conference*, August 2000. An early version of the paper appeared as techreport UW-CSE-00-02-01 available at: <http://www.cs.washington.edu/homes/savage/traceback.html>.
- [3] Alex C. Snoeren, Craig Partridge, Luis A. Sanchez, Christine E. Jones, Fabrice Tchakountio, Stephen T. Kent, and W. Timothy Strayer. Hash-Based IP traceback. In Roch Guerin, editor, *Proceedings of the ACM SIGCOMM 2001 Conference (SIGCOMM-01)*, volume 31, 4 of *Computer Communication Review*, pages 3–14, New York, August 27–31 2001. ACM Press.
- [4] Alex C. Snoeren, Craig Partridge, Luis A. Sanchez, Christine E. Jones, Fabrice Tchakountio, Beverly Schwartz, Stephen T. Kent, and W. Timothy Strayer. Single-packet ip traceback. *IEEE/ACM Trans. Netw.*, 10(6):721–734, 2002.
- [5] Stuart Staniford, Vern Paxson, and Nicholas Weaver. How to Own the internet in your spare time. In *Proceedings of the 11th USENIX Security Symposium (Security '02)*. To be appear. URL:<http://www.cs.berkeley.edu/~nweaver/cdc.web/>.