

End-to-end IP Service Quality and Mobility

- Lecture #2 -

Special Course in Networking Technology

S-38.215

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Agenda

- Agenda
- Goal of this lecture
- Rôles
 - Endpoint
 - Access
- Characteristics of applications
- Mobile applications

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Goal of this lecture

- Definition of the characteristics of applications to be used in mobile environment.
 - “TSpec” this lecture, “RSpec” the next one.
- Impact of multi-access mobile environment
 - Endpoint
 - Access technology
- Analysis of services:
 - Service events that make up the service instance.
 - Service event characteristics.
 - How service is instantiated.

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Rôle of the endpoint

- In multi-access network, the endpoint capabilities vary.
- PC-type endpoint:
 - Lots of processing power, large memory.
 - Large display, stereo speakers.
 - Keyboard & mouse available.
- PDA/cellular phone –type endpoint:
 - More limited processing power & memory.
 - Display smaller, sound capabilities limited.
 - Input devices more constrained.
- **Conclusion:** preferably service instantiation dependent on the type of endpoint.

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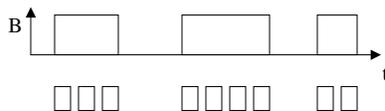
Rôle of access technology

- Access technology has an effect on services:
 - Available bandwidth.
 - RTT, delay variation, packet loss.
 - Service availability.
 - Service instantiation time varies.
 - Service quality support mechanisms vary per access technology.
 - Service quality support mechanisms available.
 - Access network may not be able to provide support for real-time services.
 - There may be no guarantees.
 - Mobility support varies.

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Case VoIP media stream

- VoIP with 8 kHz PCM coding, 8 bit sample resolution
 - Bi-directional
 - “On/off” traffic pattern, talkspurts & silence periods
 - 20 ms frame size => 160 bytes / sample
 - 1 sample / packet
- ⇒ During talkspurts, a 200-byte packet is transmitted every 20 ms (IPv4); during silence periods no packets (VAD).
- ⇒ Requires 80 kbit/s IP layer throughput



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Case HTTP browsing

- HTTP browsing without pipelining.
 - Request / response pattern.
 - HTTP GET requests typically small, HTML page sizes vary.
 - Large non-textual objects such as can be embedded into HTML pages.
 - User wants interactivity: something happens “soon” when a hyperlink is clicked.
 - Non-first time users understand that downloading large content takes time.
 - It is the total downloading time that counts
- ⇒ Small messages best given some statistical capacity in “uplink” direction, downlink can vary more.

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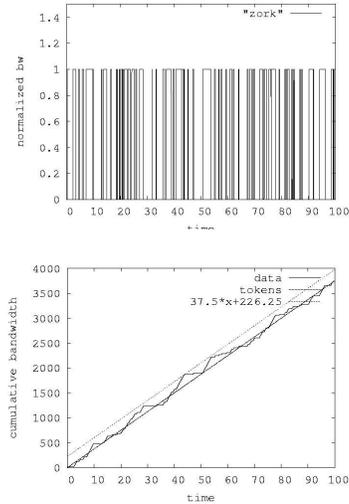
Token bucket parameters

- Bursty VoIP streams can be shaped
 - Token bucket parameters (r , B , M , m) need to match the nominal bit rate and media stream characteristics of the codec (see [example](#)).
 - Buffering delay must be minimized.
- Bursty browsing stream (downlink) has no intrinsic bandwidth requirement.
 - End-to-end delay can be larger ⇒ longer shaping buffering possible.
 - HTTP is run on top of a reliable transport protocol
⇒ throughput is a design parameter.
- TB parameters establish a bandwidth envelope: $\int_0^T d(t) dt \leq rT + B$

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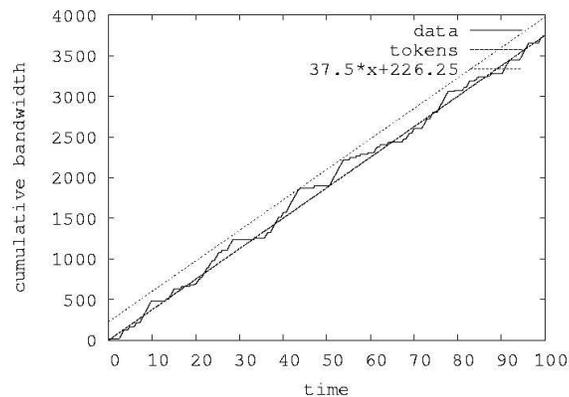
Conditioning example

- VoIP media stream: Poisson modulated ON/OFF process.
- Let us assume
 - Average talkspurt length of 0.6 seconds
 - Average silence period length of 1 seconds
 - Imaginary codec with 100 units/s output rate
- Example trace:
 - $R = 37.5$ units/s
 - $B = 226.25$ units



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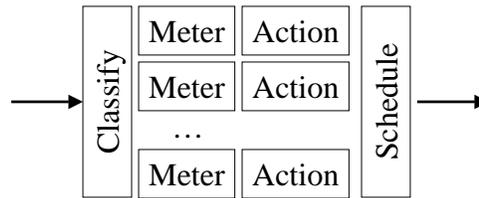
Conditioning example / 2



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Edge conditioning

- Edge conditioning consists of three parts:
 - Classification
 - Metering
 - Action
- Action can be
 - Shaping
 - (Re)marking
 - Dropping
 - Null action
- Metering is based on token bucket parameters
 - Per flow
 - Per traffic aggregate
 - Based on SLA



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Bandwidth utility

- *Bandwidth utility* $U(b)$ tells how much an application can benefit from bandwidth.
- No single measure exists.
- Phenomenological differences between applications:
 - CBR VoIP codec: $U(b)$ almost a step function.
 - G.723.1: two-step function.
 - AMR codec: more steps.
 - Downloading of a large HTTP page: “continuous” function.
- If utility function is known, bandwidths can be allocated by maximizing the total utility of allocated bandwidths.
- In real systems, utility can only be approximated. $U := \max_{\{b_i\}} \sum_i U(b_i)$

[F.Kelly, *European Trans. Telecom. Vol. 8, no.1, 1997, p.33 ff.*]

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Bandwidth utility examples

- Simple utility for elastic applications: linear in bandwidth.
 - Can be made more elaborate by taking into account application type with a per-app linear coefficient of per-app B parameter.
 - Simple-minded application of super-linear utility functions dangerous: utility maximized by non-homogeneous resource allocation.
 - In example, criterion $M > C/B$.
- => **Fairness important.**
- Max-min utility schemes.

$$U(b) = \begin{cases} 1, & b \geq B \\ \frac{b}{B}, & b < B \end{cases}$$

$$U_i(b) = C_i \begin{cases} 1, & b \geq B \\ \left(\frac{b}{B}\right)^\alpha, & b < B \end{cases}$$

[Kao, Zhegura, Utility max-min: an application-oriented..., Infocom'99 p.793 ff.]
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Average throughput

- Some applications are inherently bursty, yet require near-constant throughput.
 - VoIP
- Another example: H.261 video streaming codec
 - Full picture sent periodically => bandwidth peak.
 - Intermediate frame size depends on motion.
- Token bucket regulator must accommodate momentary bandwidth peaks, but typically end-to-end delay allows for shaping to smoothen the bandwidth.
- Contrast with browsing: TB regulator does not need to accommodate all bandwidth peaks; longer delay allowable.

[A. Adas, Traffic models in broadband networks, IEEE Comm. Mag. July 1997, p.82 ff]

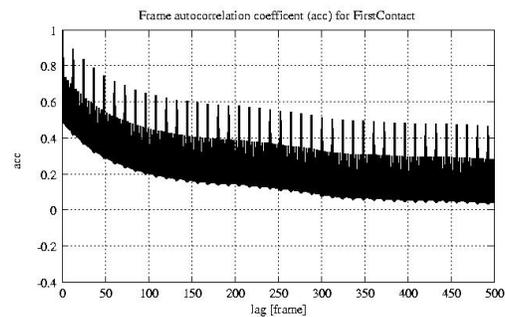
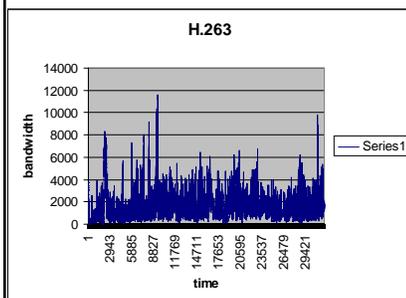
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Estimation of token bucket parameters

- Straightforward method:
 - Compute token rate
 - Get bucket depth from arrivals process
 - Need to have representative sampling.
 - In IP affected by
 - header overhead
 - MTU size.
- ⇒ Real token bucket parameters used by the network cannot always be fully determined on application layer without referring to link layer.

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Example: H.263 trace (UBR)



[\[http://www.tkn.ee.tu-berlin.de/research/trace/trace.html\]](http://www.tkn.ee.tu-berlin.de/research/trace/trace.html)

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Application characteristics

- Types of service events involved.
 - Characteristics may be different per component.
- Unidirectional / bi-directional.
 - Asymmetry of characteristics
- Type of bandwidth utility: smooth vs. step-like.
 - => Designed bandwidth possible or not.
- Burstiness characteristic: token bucket parameters

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Useful mobile applications

- Mobile environment creates new needs for applications.
- Mobile application types:
 - Telephony / conferencing
 - Instant / near-instant messaging
 - M-mail
 - Fetching of information
 - Killing time

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Mobile applications

- Off-line data transfer (e-mail)
- On-line data transfer
- Browsing
- Messaging
- Presence
- Chat
- Games
- Streaming
- Multimedia-over-IP conferencing

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Data transfer

- Examples: e-mail, downloading of a picture, a piece of music or a game.
- Characteristics:
 - Content size can be large.
 - E-mail: more or less symmetric, downloading: unidirectional/asymmetric.
 - Irregular arrival process for service events.
 - Traffic can be bursty.
 - Applications typically very elastic with respect to utility.
 - TCP-based.
- Conclusion:
 - Bandwidth designed

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Interactive applications

- Examples: browsing, M-commerce.
- Characteristics:
 - Content size varies.
 - Symmetric to asymmetric bandwidth characteristics.
 - Some services have real service instances.
 - Arrival process for service instances random.
 - Arrival process for service events also random, but may be frequent.
 - Request and response temporally close to each other.
 - TCP-based.
- Conclusion: sufficient r & B for most HTTP/WAP requests, downlink throughput can be designed.

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Simple SIP services

- Examples: (instant) messaging, chat/presence.
- Characteristics:
 - Content size varies.
 - Can be unidirectional, symmetric, or asymmetric.
 - Arrival process random for service instances and service events.
 - Some degree of interactivity expected of IM, chat.
 - TCP or UDP based.
- Conclusion:
 - Sufficient throughput for small messages, throughput for large messages may be designed.
 - UDP: shaping rather than dropping.

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Streaming

- Examples: audio or video streaming.
- Characteristics:
 - RTSP signalling:
 - Interactivity.
 - TCP or UDP.
 - Media flow:
 - Periodic PDU stream.
 - Can be bursty or already shaped.
 - UDP, for firewall traversal reasons also TCP.
- Conclusion:
 - Need to provide enough bandwidth, but can apply shaping to media flow.

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Multimedia conferencing

- Examples: audio or video conferencing.
- Characteristics:
 - SIP signalling:
 - Call set-up time requirements.
 - TCP or UDP.
 - Media flow:
 - Periodic PDU stream.
 - Can be bursty (video) or already shaped.
 - UDP, for firewall traversal reasons also TCP.
- Conclusion:
 - Need to provide enough bandwidth for media stream & signalling.

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Summary

- The following characteristics of applications need to be known:
 - Token bucket parameters.
 - Not necessary same on application level and network level.
 - Bandwidth utility.
 - May be application type dependent.
 - Fairness considerations.
- Mobile application categories:
 - Data transfer
 - Interactive applications
 - Simple SIP services
 - Streaming
 - Multimedia conferencing

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