



8. Queueing systems

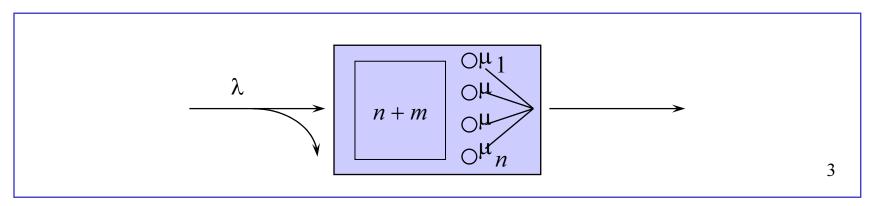
Contents

- Refresher: Simple teletraffic model
- Queueing discipline
- M/M/1 (1 server, ∞ waiting places)
- Application to packet level modelling of data traffic
- M/M/n (*n* servers, ∞ waiting places)

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Simple teletraffic model

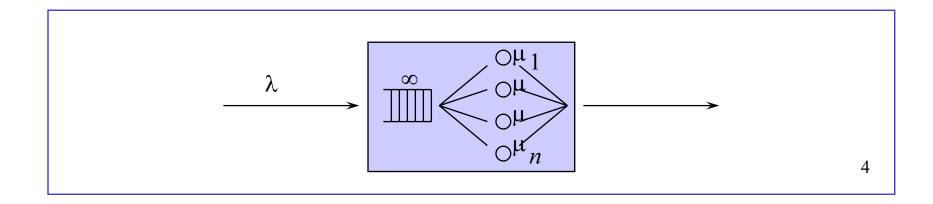
- **Customers arrive** at rate λ (customers per time unit)
 - $1/\lambda$ = average inter-arrival time
- Customers are served by n parallel servers
- When busy, a server serves at rate μ (customers per time unit)
 - $1/\mu$ = average service time of a customer
- There are n + m customer places in the system
 - at least *n* service places and at most *m* waiting places
- It is assumed that blocked customers (arriving in a full system) are lost



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Pure queueing system

- Finite number of servers (n < ∞), n service places, infinite number of waiting places (m = ∞)
 - If all *n* servers are occupied when a customer arrives, it occupies one of the waiting places
 - No customers are lost but some of them have to wait before getting served
- From the customer's point of view, it is interesting to know e.g.
 - what is the probability that it has to wait "too long"?



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Queueing discipline

- Consider a single server (n = 1) queueing system
- Queueing discipline determines the way the server serves the customers
 - It tells
 - whether the customers are served one-by-one or simultaneously
 - Furthermore, if the customers are served one-by-one, it tells
 - in which order they are taken into the service
 - And if the customers are served simultaneously, it tells
 - how the service capacity is shared among them
- Note: In computer systems the corresponding concept is scheduling
- A queueing discipline is called work-conserving if customers are served with full service rate μ whenever the system is non-empty

Work-conserving queueing disciplines

- First In First Out (**FIFO**) = First Come First Served (FCFS)
 - ordinary queueing discipline ("queue")
 - arrival order = service order
 - customers served one-by-one (with full service rate μ)
 - always serve the customer that has been waiting for the longest time
 - default queueing discipline in this lecture
- Last In First Out (LIFO) = Last Come First Served (LCFS)
 - reversed queuing discipline ("stack")
 - customers served one-by-one (with full service rate μ)
 - always serve the customer that has been waiting for the shortest time
- Processor Sharing (**PS**)
 - "fair queueing"
 - customers served simultaneously
 - when *i* customers in the system, each of them served with equal rate μ/i
 - see Lecture 9. Sharing systems

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M/M/1 queue

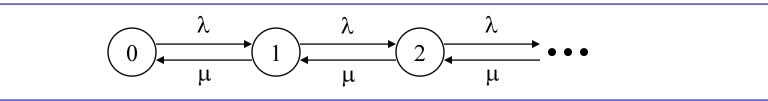
- Consider the following simple teletraffic model:
 - Infinite number of independent customers ($k = \infty$)
 - Interarrival times are IID and exponentially distributed with mean $1/\lambda$
 - so, customers arrive according to a Poisson process with intensity λ
 - One server (n = 1)
 - Service times are IID and exponentially distributed with mean $1/\mu$
 - Infinite number of waiting places $(m = \infty)$
 - Default queueing discipline: FIFO
- Using Kendall's notation, this is an M/M/1 queue
 - more precisely: M/M/1-FIFO queue
- Notation:
 - $\rho = \lambda/\mu = \text{traffic load}$

Related random variables

- X = number of customers in the system at an arbitrary time = queue length in equilibrium
- X* = number of customers in the system at an (typical) arrival time
 = queue length seen by an arriving customer
- W = waiting time of a (typical) customer
- *S* = service time of a (typical) customer
- D = W + S = total time in the system of a (typical) customer = delay

State transition diagram

- Let X(t) denote the number of customers in the system at time t
 - Assume that X(t) = i at some time t, and consider what happens during a short time interval (t, t+h]:
 - with prob. $\lambda h + o(h)$, a new customer arrives (state transition $i \rightarrow i+1$)
 - if i > 0, then, with prob. $\mu h + o(h)$, a customer leaves the system (state transition $i \rightarrow i-1$)
- Process X(t) is clearly a Markov process with state transition diagram



• Note that process X(t) is an irreducible birth-death process with an infinite state space $S = \{0, 1, 2, ...\}$

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Equilibrium distribution (1)

• Local balance equations (LBE):

$$\pi_{i}\lambda = \pi_{i+1}\mu \qquad \text{(LBE)}$$

$$\Rightarrow \pi_{i+1} = \frac{\lambda}{\mu}\pi_{i} = \rho\pi_{i}$$

$$\Rightarrow \pi_{i} = \rho^{i}\pi_{0}, \quad i = 0, 1, 2, \dots$$

• Normalizing condition (N):

$$\sum_{i=0}^{\infty} \pi_{i} = \pi_{0} \sum_{i=0}^{\infty} \rho^{i} = 1$$
 (N)
$$\Rightarrow \pi_{0} = \left(\sum_{i=0}^{\infty} \rho^{i}\right)^{-1} = \left(\frac{1}{1-\rho}\right)^{-1} = 1-\rho, \text{ if } \rho < 1$$

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Equilibrium distribution (2)

 Thus, for a stable system (ρ < 1), the equilibrium distribution exists and is a geometric distribution:

$$\rho < 1 \implies X \sim \text{Geom}(\rho)$$

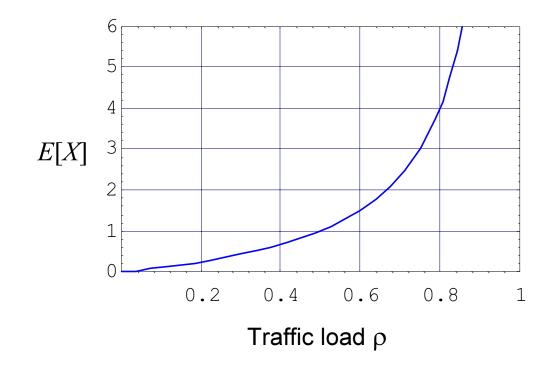
$$P\{X=i\} = \pi_i = (1-\rho)\rho^i, \quad i = 0,1,2,\dots$$

$$E[X] = \frac{\rho}{1-\rho}, \quad D^2[X] = \frac{\rho}{(1-\rho)^2}$$

Remark:

- This result is valid for any work-conserving queueing discipline (FIFO, LIFO, PS, ...)
- This result is **not insensitive** to the service time distribution for **FIFO**
 - even the mean queue length E[X] depends on the distribution
- However, for any symmetric queueing discipline (such as LIFO or PS) the result is, indeed, insensitive to the service time distribution

Mean queue length E[X] vs. traffic load ρ



Mean delay

- Let D denote the total time (delay) in the system of a (typical) customer
 - including both the waiting time W and the service time S: D = W + S
- Little's formula: $E[X] = \lambda \cdot E[D]$. Thus,

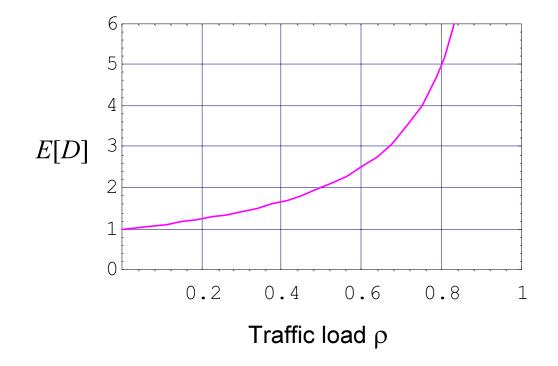
$$E[D] = \frac{E[X]}{\lambda} = \frac{1}{\lambda} \cdot \frac{\rho}{1-\rho} = \frac{1}{\mu} \cdot \frac{1}{1-\rho} = \frac{1}{\mu-\lambda}$$

Remark:

- The mean delay is the same for all work-conserving queueing disciplines (FIFO, LIFO, PS, ...)
- But the variance and other moments are different!

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Mean delay E[D] vs. traffic load ρ



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Mean waiting time

- Let W denote the waiting time of a (typical) customer
- Since W = D S, we have

$$E[W] = E[D] - E[S] = \frac{1}{\mu} \cdot \frac{1}{1-\rho} - \frac{1}{\mu} = \frac{1}{\mu} \cdot \frac{\rho}{1-\rho}$$

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Waiting time distribution (1)

- Let *W* denote the waiting time of a (typical) customer
- Let X* denote the number of customers in the system at the arrival time
- **PASTA**: $P\{X^* = i\} = P\{X = i\} = \pi_i$.
- Assume now, for a while, that $X^* = i$
 - Service times S_2, \ldots, S_i of the waiting customers are IID and $\sim Exp(\mu)$
 - Due to the memoryless property of the exponential distribution, the **remaining** service time S_1^* of the customer in service also follows $Exp(\mu)$ -distribution (and is independent of everything else)
 - Due to the FIFO queueing discipline, $W = S_1^* + S_2 + \ldots + S_i$
 - Construct a Poisson (point) process τ_n by defining $\tau_1 = S_1^*$ and $\tau_n = S_1^* + S_2 + \ldots + S_n$, $n \ge 2$. Now (since $X^* = i$): $W > t \Leftrightarrow \tau_i > t$

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Waiting time distribution (2)

• Since
$$W = 0 \Leftrightarrow X^* = 0$$
, we have
 $P\{W = 0\} = P\{X^* = 0\} = \pi_0 = 1 - \rho$
 $P\{W > t\} = \sum_{i=1}^{\infty} P\{W > t \mid X^* = i\} P\{X^* = i\}$
 $= \sum_{i=1}^{\infty} P\{\tau_i > t\} \pi_i = \sum_{i=1}^{\infty} P\{\tau_i > t\} (1 - \rho) \rho^i$

- Denote by A(t) the Poisson (counter) process corresponding to τ_n
 - It follows that: $\tau_i > t \Leftrightarrow A(t) \le i-1$
 - On the other hand, we know that $A(t) \sim \text{Poisson}(\mu t)$. Thus,

$$P\{\tau_i > t\} = P\{A(t) \le i - 1\} = \sum_{j=0}^{i-1} \frac{(\mu t)^j}{j!} e^{-\mu t}$$
¹⁹

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Waiting time distribution (3)

• By combining the previous formulas, we get

$$P\{W > t\} = \sum_{i=1}^{\infty} P\{\tau_i > t\} (1-\rho) \rho^i$$

= $\sum_{i=1}^{\infty} \sum_{j=0}^{i-1} \frac{(\mu t)^j}{j!} e^{-\mu t} (1-\rho) \rho^i$
= $\rho \sum_{j=0}^{\infty} \frac{(\mu t \rho)^j}{j!} e^{-\mu t} (1-\rho) \sum_{i=j+1}^{\infty} \rho^{i-(j+1)}$
= $\rho \sum_{j=0}^{\infty} \frac{(\mu t \rho)^j}{j!} e^{-\mu t} = \rho e^{\mu t \rho} e^{-\mu t} = \rho e^{-\mu (1-\rho)t}$

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Waiting time distribution (4)

 Waiting time W can thus be presented as a product W = JD of two independent random variables J ~ Bernoulli(ρ) and D ~ Exp(μ(1-ρ)):

$$P\{W = 0\} = P\{J = 0\} = 1 - \rho$$

$$P\{W > t\} = P\{J = 1, D > t\} = \rho \cdot e^{-\mu(1-\rho)t}, \quad t > 0$$

$$E[W] = E[J]E[D] = \rho \cdot \frac{1}{\mu(1-\rho)} = \frac{1}{\mu} \cdot \frac{\rho}{1-\rho}$$

$$E[W^{2}] = P\{J = 1\}E[D^{2}] = \rho \cdot \frac{2}{\mu^{2}(1-\rho)^{2}} = \frac{1}{\mu^{2}} \cdot \frac{2\rho}{(1-\rho)^{2}}$$

$$D^{2}[W] = E[W^{2}] - E[W]^{2} = \frac{1}{\mu^{2}} \cdot \frac{\rho(2-\rho)}{(1-\rho)^{2}}$$

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Application to packet level modelling of data traffic

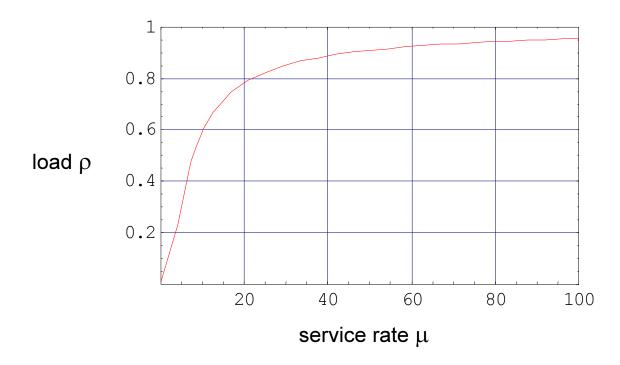
- M/M/1 model may be applied (to some extent) to packet level modelling of data traffic
 - customer = IP packet
 - λ = packet arrival rate (packets per time unit)
 - $1/\mu$ = average packet transmission time (aikayks.)
 - $\rho = \lambda/\mu = \text{traffic load}$
- Quality of service is measured e.g. by the packet delay
 - P_z = probability that a packet has to wait "too long", i.e. longer than a given reference value z

$$P_z = P\{W > z\} = \rho e^{-\mu(1-\rho)z}$$

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Multiplexing gain

- We determine load ρ so that prob. $P_z < 1\%$ for z = 1 (time units)
- Multiplexing gain is described by the traffic load ρ as a function of the service rate μ



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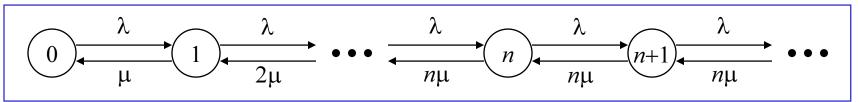
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M/M/n queue

- Consider the following simple teletraffic model:
 - Infinite number of independent customers ($k = \infty$)
 - Interarrival times are IID and exponentially distributed with mean $1/\lambda$
 - so, customers arrive according to a Poisson process with intensity λ
 - Finite number of servers ($n < \infty$)
 - Service times are IID and exponentially distributed with mean $1/\mu$
 - Infinite number of waiting places $(m = \infty)$
 - Default queueing discipline: FCFS
- Using Kendall's notation, this is an M/M/*n* queue
 - more precisely: M/M/n-FCFS queue
- Notation:
 - $\rho = \lambda/(n\mu) = \text{traffic load}$

State transition diagram

- Let X(t) denote the number of customers in the system at time t
 - Assume that X(t) = i at some time t, and consider what happens during a short time interval (t, t+h]:
 - with prob. $\lambda h + o(h)$, a new customer arrives (state transition $i \rightarrow i+1$)
 - if i > 0, then, with prob. $\min\{i,n\} \cdot \mu h + o(h)$, a customer leaves the system (state transition $i \rightarrow i-1$)
- Process *X*(*t*) is clearly a Markov process with state transition diagram



• Note that process X(t) is an irreducible birth-death process with an infinite state space $S = \{0, 1, 2, ...\}$

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Equilibrium distribution (1)

• Local balance equations (LBE) for i < n:

$$\pi_{i}\lambda = \pi_{i+1}(i+1)\mu \qquad (\text{LBE})$$

$$\Rightarrow \pi_{i+1} = \frac{\lambda}{(i+1)\mu}\pi_{i} = \frac{n\rho}{i+1}\pi_{i}$$

$$\Rightarrow \pi_{i} = \frac{(n\rho)^{i}}{i!}\pi_{0}, \quad i = 0, 1, \dots, n$$

• Local balance equations (LBE) for $i \ge n$:

$$\pi_{i}\lambda = \pi_{i+1}n\mu \qquad (\text{LBE})$$

$$\Rightarrow \pi_{i+1} = \frac{\lambda}{n\mu}\pi_{i} = \rho\pi_{i}$$

$$\Rightarrow \pi_{i} = \rho^{i-n}\pi_{n} = \rho^{i-n}\frac{(n\rho)^{n}}{n!}\pi_{0} = \frac{n^{n}\rho^{i}}{n!}\pi_{0}, \quad i = n, n+1, \dots 28$$

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Equilibrium distribution (2)

• Normalizing condition (N):

$$\sum_{i=0}^{\infty} \pi_{i} = \pi_{0} \left(\sum_{i=0}^{n-1} \frac{(n\rho)^{i}}{i!} + \sum_{i=n}^{\infty} \frac{n^{n}\rho^{i}}{n!} \right) = 1$$
(N)

$$\Rightarrow \pi_{0} = \left(\sum_{i=0}^{n-1} \frac{(n\rho)^{i}}{i!} + \frac{(n\rho)^{n}}{n!} \sum_{i=n}^{\infty} \rho^{i-n} \right)^{-1}$$

$$= \left(\sum_{i=0}^{n-1} \frac{(n\rho)^{i}}{i!} + \frac{(n\rho)^{n}}{n!(1-\rho)} \right)^{-1} = \frac{1}{\alpha + \beta}, \text{ if } \rho < 1$$

Notation: $\alpha = \sum_{i=0}^{n-1} \frac{(n\rho)^{i}}{i!}, \quad \beta = \frac{(n\rho)^{n}}{n!(1-\rho)}$
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Equilibrium distribution (3)

• Thus, for a **stable** system ($\rho < 1$, that is: $\lambda < n\mu$), the equilibrium distribution exists and is as follows:

$$\rho < 1 \implies$$

$$P\{X = i\} = \pi_i = \begin{cases} \frac{(n\rho)^i}{i!} \cdot \frac{1}{\alpha + \beta}, & i = 0, 1, \dots, n\\ \frac{n^n \rho^i}{n!} \cdot \frac{1}{\alpha + \beta}, & i = n, n + 1, \dots \end{cases}$$

$$n = 1: \ \alpha = 1, \ \beta = \frac{\rho}{1-\rho}, \ \pi_0 = \frac{1}{\alpha+\beta} = 1-\rho$$
$$n = 2: \ \alpha = 1+2\rho, \ \beta = \frac{2\rho^2}{1-\rho}, \ \pi_0 = \frac{1}{\alpha+\beta} = \frac{1-\rho}{1+\rho}$$

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Probability of waiting

- Let p_W denote the probability that an arriving customer has to wait
- Let *X** denote the number of customers in the system at an arrival time
- An arriving customer has to wait whenever all the servers are occupied at her arrival time. Thus,

$$p_W = P\{X^* \ge n\}$$

• PASTA:
$$P\{X^* = i\} = P\{X = i\} = \pi_i$$
. Thus,

$$p_W = P\{X^* \ge n\} = \sum_{i=n}^{\infty} \pi_i = \sum_{i=n}^{\infty} \pi_0 \cdot \frac{n^n \rho^i}{n!} = \pi_0 \cdot \frac{(n\rho)^n}{n!(1-\rho)} = \frac{\beta}{\alpha+\beta}$$

n=1:
$$p_W = \rho$$

n=2: $p_W = \frac{2\rho^2}{1+\rho}$
₃₁

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Mean number of waiting customers

- Let X_W denote the number of waiting customers in equilibrium
- Then

$$E[X_W] = \sum_{i=n}^{\infty} (i-n)\pi_i = \pi_0 \frac{(n\rho)^n}{n!(1-\rho)} \sum_{i=n}^{\infty} (i-n) \cdot (1-\rho)\rho^{i-n}$$
$$= p_W \cdot \frac{\rho}{1-\rho}$$

$$n = 1: \quad E[X_W] = p_W \cdot \frac{\rho}{1-\rho} = \frac{\rho^2}{1-\rho}$$
$$n = 2: \quad E[X_W] = p_W \cdot \frac{\rho}{1-\rho} = \frac{2\rho^2}{1+\rho} \cdot \frac{\rho}{1-\rho} = \frac{2\rho^3}{1-\rho^2}$$

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Mean waiting time

- Let W denote the waiting time of a (typical) customer
- Little's formula: $E[X_W] = \lambda \cdot E[W]$. Thus,

$$E[W] = \frac{E[X_W]}{\lambda} = \frac{1}{\lambda} \cdot p_W \cdot \frac{\rho}{1-\rho} = \frac{1}{\mu} \cdot \frac{p_W}{n(1-\rho)} = p_W \cdot \frac{1}{n\mu-\lambda}$$

$$n = 1: \quad E[W] = \frac{1}{\mu} \cdot \frac{p_W}{1 - \rho} = \frac{1}{\mu} \cdot \frac{\rho}{1 - \rho}$$
$$n = 2: \quad E[W] = \frac{1}{\mu} \cdot \frac{p_W}{2(1 - \rho)} = \frac{1}{\mu} \cdot \frac{\rho^2}{1 - \rho^2}$$

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Mean delay

- Let D denote the total time (delay) in the system of a (typical) customer
 - including both the waiting time W and the service time S: D = W + S
- Then,

$$E[D] = E[W] + E[S] = \frac{1}{\mu} \cdot \left(\frac{p_W}{n(1-\rho)} + 1\right) = p_W \cdot \frac{1}{n\mu - \lambda} + \frac{1}{\mu}$$

$$n = 1: \quad E[D] = \frac{1}{\mu} \cdot \left(\frac{p_W}{1-\rho} + 1\right) = \frac{1}{\mu} \cdot \left(\frac{\rho}{1-\rho} + 1\right) = \frac{1}{\mu} \cdot \frac{1}{1-\rho}$$
$$n = 2: \quad E[D] = \frac{1}{\mu} \cdot \frac{p_W}{2(1-\rho)} = \frac{1}{\mu} \cdot \left(\frac{\rho^2}{1-\rho^2} + 1\right) = \frac{1}{\mu} \cdot \frac{1}{1-\rho^2}$$

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Mean queue length

- Let *X* denote the number of customers in the system (queue length) in equilibrium
- Little's formula: $E[X] = \lambda \cdot E[D]$. Thus,

$$E[X] = \lambda \cdot E[D] = p_W \cdot \frac{\lambda}{n\mu - \lambda} + \frac{\lambda}{\mu} = p_W \cdot \frac{\rho}{1 - \rho} + n\rho$$

$$n = 1: \quad E[X] = p_W \cdot \frac{\rho}{1-\rho} + \rho = \rho \cdot \frac{\rho}{1-\rho} + \rho = \frac{\rho}{1-\rho}$$
$$n = 2: \quad E[X] = p_W \cdot \frac{\rho}{1-\rho} + 2\rho = \frac{2\rho^2}{1+\rho} \cdot \frac{\rho}{1-\rho} + 2\rho = \frac{2\rho}{1-\rho^2}$$

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Waiting time distribution (1)

- Let *W* denote the waiting time of a (typical) customer
- Let X^* denote the number of customers in the system at the arrival time
- The customer has to wait only if $X^* \ge n$. This happens with prob. p_W .
- Under the assumption that $X^* = i \ge n$, the system, however, looks like an ordinary M/M/1 queue with arrival rate λ and service rate $n\mu$.
 - Let W' denote the waiting time of a (typical) customer in this M/M/1 queue
 - Let X^* ' denote the number of customers in the system at the arrival time
- It follows that

$$P\{W = 0\} = 1 - p_W$$

$$P\{W > t\} = P\{X^* \ge n\}P\{W > t \mid X^* \ge n\}$$

$$= p_W \cdot P\{W' > t \mid X^{*'} \ge 1\} = p_W \cdot e^{-n\mu(1-\rho)t}, \quad t > 0$$

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Waiting time distribution (2)

• Waiting time *W* can thus be presented as a product W = JD' of two indep. random variables $J \sim \text{Bernoulli}(p_W)$ and $D' \sim \text{Exp}(n\mu(1-\rho))$:

$$P\{W=0\} = P\{J=0\} = 1 - p_W$$

$$P\{W > t\} = P\{J = 1, D' > t\} = p_W \cdot e^{-n\mu(1-\rho)t}, \quad t > 0$$

$$E[W] = E[J]E[D'] = p_W \cdot \frac{1}{n\mu(1-\rho)} = \frac{1}{\mu} \cdot \frac{p_W}{n(1-\rho)}$$

$$E[W^{2}] = P\{J=1\}E[D'^{2}] = p_{W} \cdot \frac{2}{n^{2}\mu^{2}(1-\rho)^{2}} = \frac{1}{\mu^{2}} \cdot \frac{2p_{W}}{n^{2}(1-\rho)^{2}}$$
$$D^{2}[W] = E[W^{2}] - E[W]^{2} = \frac{1}{\mu^{2}} \cdot \frac{p_{W}(2-p_{W})}{n^{2}(1-\rho)^{2}}$$

Example (1)

- Printer problem
 - Consider the following two different configurations:
 - One rapid printer (IID printing times ~ $Exp(2\mu)$)
 - Two slower parallel printers (IID printing times ~ $Exp(\mu)$)
 - Criterion: minimize mean delay E[D]
 - One rapid printer (M/M/1 model with $\rho = \lambda/(2\mu)$):

$$E[D_1] = \frac{1}{2\mu} \cdot \frac{1}{1-\rho}$$

• Two slower printers (M/M/2 model with $\rho = \lambda/(2\mu)$):

$$E[D_2] = \frac{1}{\mu} \cdot \frac{1}{1-\rho^2} = \frac{1}{2\mu} \cdot \frac{2}{(1-\rho)(1+\rho)} = E[D_1] \cdot \frac{2}{1+\rho} > E[D_1]$$

