

lect01.ppt

S-38.1145 - Introduction to Teletraffic Theory - Spring 2006

1. Introduction

Contents

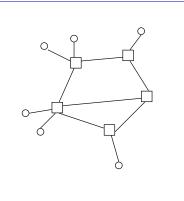
- · Telecommunication networks and switching modes
- Purpose of Teletraffic Theory
- Teletraffic models
- Little's formula

2

1. Introduction

Telecommunication network

- · A simple model of a telecommunication network consists of
 - nodes
 - terminals
 - 0 · network nodes
 - links between nodes
- Access network
 - connects the terminals to the network nodes
- Trunk network
 - connects the network nodes to each other



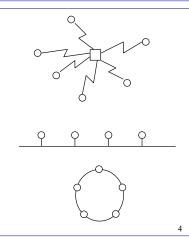
1. Introduction

Shared medium as an access network

- · In the previous model,
 - connections between terminals and network nodes are point-to**point** type (⇒ no resource sharing within the access netw.)
- · In some cases, such as
 - mobile telephone network
 - local area network (LAN) connecting computers

the access network consists of shared medium:

- users have to compete for the resources of this shared medium
- multiple access (MA) techniques are needed



Switching modes

· Circuit switching

- telephone networks
- mobile telephone networks
- optical networks

Packet switching

- data networks
- two possibilities
 - connection oriented: e.g. X.25, Frame Relay
 - · connectionless: e.g. Internet (IP), SS7 (MTP)

· Cell switching

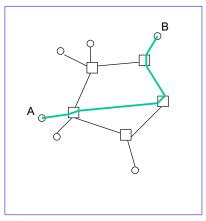
- ATM networks
- connection oriented
- fast packet switching with fixed length packets (cells)

1. Introduction

Circuit switching (1)

Connection oriented:

- connections set up end-to-end before information transfer
- resources reserved for the whole duration of connection
- if resources are not available, the call is blocked and lost
- Information transfer as continuous stream

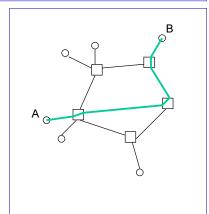


6

1. Introduction

Circuit switching (2)

- Before information transfer
 - Set-up delay
- · During information transfer
 - signal propagation delay
 - no overhead
 - no extra delays
- Example: telephone network

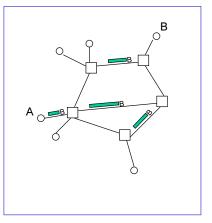


1. Introduction

Connectionless packet switching (1)

Connectionless:

- no connection set-up
- no resource reservation
- no blocking
- Information transfer as discrete packets
 - varying length
 - global address (of the destination)

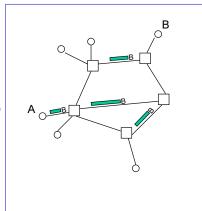


7

5

Connectionless packet switching (2)

- · Before information transfer
 - no delays
- During information transfer
 - overhead (header bytes)
 - packet processing delays
 - queueing delays (since packets compete for joint resources)
 - transmission delays (due to finite capacity links)
 - signal propagation delay
 - packet losses (due to finite buffers)
- Example: Internet (IP-layer)



9

1. Introduction

Contents

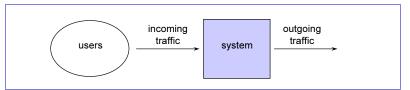
- · Telecommunication networks and switching modes
- Purpose of Teletraffic Theory
- · Teletraffic models
- Little's formula

10

1. Introduction

Traffic point of view

• Telecommunication system from the traffic point of view:

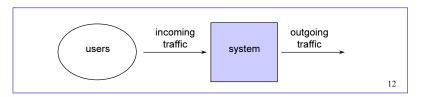


- Ideas:
 - the system serves the incoming traffic
 - the traffic is generated by the users of the system

1. Introduction

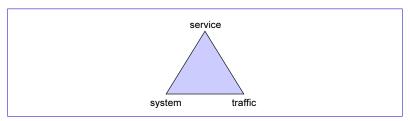
Interesting questions

- Given the system and incoming traffic, what is the quality of service experienced by the user?
- Given the incoming traffic and required quality of service, how should the system be dimensioned?
- Given the system and required quality of service, what is the maximum traffic load?



General purpose (1)

- Determine relationships between the following three factors:
 - quality of service
 - traffic load
 - system capacity



13

1. Introduction

General purpose (2)

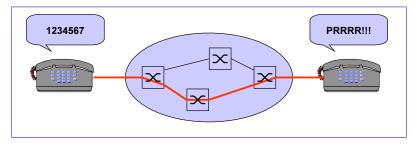
- · System can be
 - a single device (e.g. link between two telephone exchanges, link in an IP network, packet processor in a data network, router's transmission buffer, or statistical multiplexer in an ATM network)
 - the whole network (e.g. telephone or data network) or some part of it
- · Traffic consists of
 - bits, packets, bursts, flows, connections, calls, ...
 - depending on the system and time scale considered
- Quality of service can be described from the point of view of
 - the customer (e.g. call blocking, packet loss, packet delay, or throughput)
 - the system, in which case we use the term **performance** (e.g. processor or link utilization, or maximum network load)

14

1. Introduction

Example

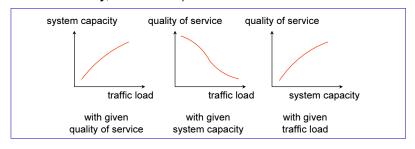
- · Telephone call
 - traffic = telephone calls by everybody
 - system = telephone network
 - quality of service = probability that the phone rings at the destination



1. Introduction

Relationships between the three factors

Qualitatively, the relationships are as follows:



To describe the relationships quantitatively, mathematical models are needed

Teletraffic models

- Teletraffic models are stochastic (= probabilistic)
 - systems themselves are usually deterministic but traffic is typically stochastic
 - "you never know, who calls you and when"
- It follows that the variables in these models are random variables, e.g.
 - number of ongoing calls
 - number of packets in a buffer
- · Random variable is described by its distribution, e.g.
 - probability that there are *n* ongoing calls
 - probability that there are n packets in the buffer
- Stochastic process describes the temporal development of a random variable

Typically,
 the mo

1. Introduction

 the model describes just one part or property of the real system under consideration and even from one point of view

Real system vs. model

- the description is not very accurate but rather approximative
- · Thus,
 - caution is needed when conclusions are drawn

17

1. Introduction

Practical goals

- · Network planning
 - dimensioning
 - optimization
 - performance analysis
- Network management and control
 - efficient operating
 - fault recovery
 - traffic management
 - routing
 - accounting

1. Introduction

Literature

- Teletraffic Theory
 - Teletronikk Vol. 91, Nr. 2/3, Special Issue on "Teletraffic", 1995
 - V. B. Iversen, Teletraffic Engineering Handbook, http://www.tele.dtu.dk/teletraffic/handbook/telehook.pdf
 - J. Roberts, Traffic Theory and the Internet,
 IEEE Communications Magazine, Jan. 2001, pp. 94-99
 http://perso.rd.francetelecom.fr/roberts/Pub/Rob01.pdf
- Queueing Theory
 - L. Kleinrock, Queueing Systems, Vol. I: Theory, Wiley, 1975
 - L. Kleinrock, Queueing Systems, Vol. II: Computer Applications, Wiley, 1976
 - D. Bertsekas and R. Gallager, *Data Networks*, 2nd ed., Prentice-Hall, 1992
 - Myron Hlynka's Queueing Theory Page http://www2.uwindsor.ca/~hlynka/queue.html

19

20

Contents

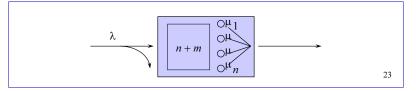
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- · Teletraffic models
- · Little's formula

21

1. Introduction

Simple teletraffic model

- Customers arrive at rate λ (customers per time unit)
 - $-1/\lambda$ = average inter-arrival time
- Customers are served by n parallel servers
- When busy, a server serves at rate μ (customers per time unit)
 - $-1/\mu$ = average service time of a customer
- There are n + m customer places in the system
 - at least *n* service places and at most *m* waiting places
- It is assumed that blocked customers (arriving in a full system) are lost



1. Introduction

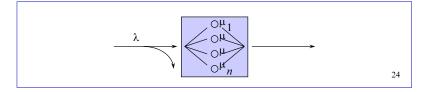
Teletraffic model types

- · Three types of system models:
 - loss systems
 - queueing systems
 - sharing systems
- · Next we will present simple teletraffic models
 - describing a single resource
- These models can be combined to create models for whole telecommunication networks
 - loss networks
 - queueing networks
 - sharing networks

1. Introduction

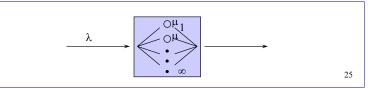
Pure loss system

- Finite number of servers $(n < \infty)$, n service places, no waiting places (m = 0)
 - If the system is full (with all n servers occupied) when a customer arrives, it is not served at all but lost
 - Some customers may be lost
- · From the customer's point of view, it is interesting to know e.g.
 - What is the probability that the system is full when it arrives?



Infinite system

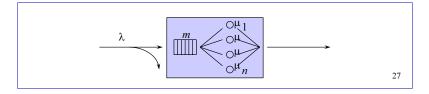
- Infinite number of servers $(n = \infty)$, no waiting places (m = 0)
 - No customers are lost or even have to wait before getting served
- · Sometimes,
 - this hypothetical model can be used to get some approximate results for a real system (with finite system capacity)
- Always,
 - it gives bounds for the performance of a real system (with finite system capacity)
 - it is much easier to analyze than the corresponding finite capacity models



1. Introduction

Lossy queueing system

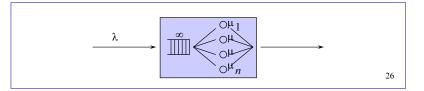
- Finite number of servers (n < ∞), n service places, finite number of waiting places (0 < m < ∞)
 - If all n servers are occupied but there are free waiting places when a customer arrives, it occupies one of the waiting places
 - If all n servers and all m waiting places are occupied when a customer arrives, it is not served at all but lost
 - Some customers are lost and some customers have to wait before getting served



1. Introduction

Pure queueing system

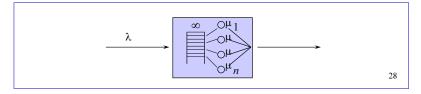
- Finite number of servers (n < ∞), n service places, infinite number of waiting places (m = ∞)
 - If all n servers are occupied when a customer arrives, it occupies one of the waiting places
 - No customers are lost but some of them have to wait before getting served
- From the customer's point of view, it is interesting to know e.g.
 - what is the probability that it has to wait "too long"?



1. Introduction

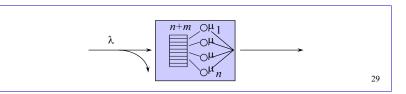
Pure sharing system

- Finite number of servers (n < ∞), infinite number of service places (n + m = ∞), no waiting places
 - If there are at most n customers in the system (x ≤ n), each customer has
 its own server. Otherwise (x > n), the total service rate (nμ) is shared fairly
 among all customers.
 - Thus, the rate at which a customer is served equals $\min\{\mu, n\mu/x\}$
 - No customers are lost, and no one needs to wait before the service.
 - But the delay is the greater, the more there are customers in the system.
 Thus, delay is an interesing measure from the customer's point of view.



Lossy sharing system

- Finite number of servers (n < ∞), finite number of service places (n + m < ∞), no waiting places
 - If there are at most n customers in the system (x ≤ n), each customer has
 its own server. Otherwise (x > n), the total service rate (nμ) is shared fairly
 among all customers.
 - Thus, the rate at which a customer is served equals $\min\{\mu, n\mu/x\}$
 - Some customers are lost, but no one needs to wait before the service.



1. Introduction

Little's formula

- · Consider a system where
 - new customers arrive at rate λ
- · Assume stability:
 - Every now and then, the system is empty
- · Consequence:
 - Customers depart from the system at rate λ
- Let
 - \overline{N} = average number of customers in the system
 - \overline{T} = average time a customer spends in the system = average delay
- · Little's formula:

$$\overline{N} = \lambda \overline{T}$$

1. Introduction

Contents

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30

1. Introduction

Proof (1)

- Let
 - N(t) = the number of customers in the system at time t
 - A(t) = the number of customers arrived in the system by time t
 - B(t) = the number of customers departed from the system by time t
 - T_i = the time customer i spends in the system = its delay
- As $t \to \infty$,

$$\frac{1}{t} \int_{0}^{t} N(s) ds \to \overline{N}, \quad \frac{1}{A(t)} \sum_{i=1}^{A(t)} T_{i} \to \overline{T}, \quad \frac{1}{B(t)} \sum_{i=1}^{B(t)} T_{i} \to \overline{T} \quad (1)$$

• In addition (due to the stability assumption),

$$\frac{1}{t}A(t) \to \lambda, \quad \frac{1}{t}B(t) \to \lambda$$
 (2)

Proof (2)

- We may assume that
 - the system is empty at time t = 0,
 - the customers depart from the system in their arrival order (FIFO)
- Then (see the figure in the following slide)

$$\sum_{i=1}^{B(t)} T_i \le \int_0^t N(s) ds \le \sum_{i=1}^{A(t)} T_i$$

· Thus,

$$\frac{B(t)}{t} \frac{1}{B(t)} \sum_{i=1}^{B(t)} T_i \le \frac{1}{t} \int_0^t N(s) ds \le \frac{A(t)}{t} \frac{1}{A(t)} \sum_{i=1}^{A(t)} T_i$$

• As $t \to \infty$, we have, by (1) and (2),

$$\lambda \overline{T} \le \overline{N} \le \lambda \overline{T}$$

33

Q.E.D.

1. Introduction

Proof (3)

