

- D3/1** Consider telephone traffic carried by a 5-channel link in the telephone network. Use the pure loss system model. New calls arrive according to a Poisson process at rate 2 calls/min, and call holding times are independently and identically distributed with mean 3 min. Determine the following traffic intensities:
- (a) the traffic offered,
  - (b) the traffic carried, and
  - (c) the traffic lost.
- D3/2** Consider the processor of a packet router in a packet switched data network. Traffic consists of data packets to be processed. Use the pure waiting system model with a single server. New packets arrive according to a Poisson process at rate 2 packets/ms, and packet processing times are independently and exponentially distributed with mean 0.4 ms.
- (a) What is the traffic load?
  - (b) What is the probability that an arriving packet will be processed immediately after the arrival (without any waiting)?
  - (c) What is the probability that a packet has to wait longer than 2 ms?
- D3/3** Consider elastic data traffic carried by a 100 Mbps link in a packet switched data network. Use the pure sharing system model with a single server. New flows arrive according to a Poisson process at rate 6 flows/s, and the average size of the files to be transferred is 15 Mb. Determine
- (a) the traffic load,
  - (b) the throughput of a flow, and
  - (c) the average file transfer time.
- D3/4** Consider telephone traffic carried by a link in a packet switched data network. A single call is modelled as a streaming CBR flow with a fixed transmission rate of 64 kbps. The link speed is  $5 \cdot 64$  kbps. Use the infinite system model. New calls arrive according to a Poisson process at rate 2 calls/min, and the average flow duration is 3 minutes. Determine
- (a) the traffic offered,
  - (b) the traffic carried, and
  - (c) the loss ratio.
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**D3/1** The model is a pure loss system of type M/G/n/n with  $n = 5$  (L3/4,10). Since  $\lambda = 2$  calls/min and  $h = 3$  min, the traffic intensity (L3/6) is

$$a = \lambda h = 2 \cdot 3 = 6 \text{ erl}$$

and the call blocking probability (L3/11)

$$B_C = \text{Erl}(n, a) = \text{Erl}(5, 6) = \frac{\frac{6^5}{120}}{1 + 6 + \frac{6^2}{2} + \frac{6^3}{6} + \frac{6^4}{24} + \frac{6^5}{120}} \stackrel{\text{num.}}{=} 0.36$$

(a) The traffic offered (L3/9):

$$a_{\text{offered}} = a = 6 \text{ erl}$$

(b) The traffic carried (L3/9):

$$a_{\text{carried}} = a(1 - B_C) = 6 \cdot (1 - 0.36) = 3.84 \text{ erl}$$

(c) The traffic lost (L3/9):

$$a_{\text{lost}} = aB_C = 6 \cdot 0.36 = 2.16 \text{ erl}$$

**D3/2** The model is a pure single server queueing system of type M/M/1 (L3/18,23) with parameters  $\lambda = 2$  packets/ms and  $1/\mu = 0.4$  ms. Thus, the probability that a packet has to wait longer than  $z$  ms is (L3/24)

$$P_z = \rho e^{-\mu(1-\rho)z},$$

where  $\rho = \lambda/\mu$  is the system load. Note, in particular, that  $P_0 = \rho$ .

(a) The system load (L3/20):

$$\rho = \frac{\lambda}{\mu} = 2 \cdot 0.4 = 0.8$$

(b) The probability that an arriving packet will be processed immediately after the arrival (without any waiting) is (L3/24):

$$P\{\text{"no waiting"}\} = 1 - P_0 = 1 - \rho = 1 - 0.8 = 0.2$$

(c) the probability that a packet has to wait longer than 2 ms is (L3/24):

$$P\{\text{"waiting longer than 2 ms"}\} = P_2 = \rho e^{-\mu(1-\rho)2} = 0.8 \cdot e^{-1.0} \stackrel{\text{num.}}{=} 0.29$$

**D3/3** The model is the pure single server sharing system of type M/G/1-PS (L3/31,36) with parameters  $\lambda = 6$  flows/s,  $S = 15$  Mb, and  $C = 100$  Mbps. The average file transfer time with full rate  $C$  is thus  $1/\mu = S/C = 15/100 = 0.15$  s.

(a) The traffic load (L3/33):

$$\rho = \frac{\lambda}{\mu} = 6 \cdot 0.15 = 0.90$$

(b) The throughput of a flow (L3/37):

$$\theta = C(1 - \rho) = 100 \cdot (1 - 0.9) = 10.0 \text{ Mbps}$$

(c) The average file transfer time (L3/35):

$$D = \frac{S}{\theta} = \frac{15}{10} = 1.5 \text{ s}$$

**D3/4** The model is the infinite system of type M/G/ $\infty$  (L3/44,48) with with the same parameters as in (a),  $\lambda = 2$  calls/min and  $h = 3$  min. The link capacity is  $C = 5 \cdot 64$  kbps, and the bit rate of a single CBR flow  $r = 64$  kbps. Thus,  $n = C/r = 5$ .

(a) The traffic offered (in erlangs)

$$a_{\text{offered}} = a = \lambda h = 2 \cdot 3 = 6 \text{ erl}$$

and as a bit rate (L3/46):

$$R_{\text{offered}} = R = ar = 6 \cdot 64 = 384 \text{ kbps}$$

(b) Let us utilize the loss ratio (determined below in (c)) to find out the traffic lost as a bit rate:

$$R_{\text{loss}} = p_{\text{loss}} R_{\text{offered}} = 0.25 \cdot 384 = 96 \text{ kbps.}$$

The traffic carried (as a bit rate) is thus

$$R_{\text{carried}} = R_{\text{offered}} - R_{\text{loss}} = 384 - 96 = 288 \text{ kbps}$$

(c) The loss ratio (L3/49):

$$p_{\text{loss}} = \text{LR}(n, a) = \text{LR}(5, 6) = \frac{1}{6} \sum_{i=6}^{\infty} (i - 5) \frac{6^i}{i!} e^{-6} \stackrel{\text{num.}}{=} 0.25$$

(The infinite sum is approximated by a finite sum with about 20 terms, which gives a sufficient accuracy.)